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DB2

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DB2 Update

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Editor

Trevor Eddolls

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DB2 tablespace options

Although DB2 data is accessed at the table level, those skilled in DB2 database design and administration know that the actual data is stored in a structure known as a tablespace. Each tablespace correlates to one or more individual physical VSAM datasets that are used to house the actual DB2 data. When designing DB2 databases, DBAs can choose from three types of tablespaces, each one useful in different circumstances. The three types of tablespace are:

- Simple tablespaces
- Segmented tablespaces
- Partitioned tablespaces.

In general, the predominant tablespace type to use for most applications is the segmented tablespace. Segmented tablespaces provide a good combination of features that mix ease of use and set-up with performance and functionality. Many organizations adhere to standards stating that new DB2 tablespaces should be segmented tablespaces unless a compelling reason exists to choose one of the other tablespace types. You should consider using the other types of DB2 tablespaces in the following cases:

- Use partitioned tablespaces when you wish to encourage parallelism. Although DB2 can and will use parallel access techniques for non-partitioned tablespaces, partitioning data helps DB2 exploit parallelism.
- Consider using partitioned tablespaces when the amount of data to be stored is very large (more than 1 million pages). You will have more control over the placement of data in separate underlying datasets using partitioned tablespaces. This is often a concern with larger DB2 tables.
- Use partitioned tablespaces to reduce utility processing time and decrease contention. It is possible to execute DB2 utilities against single partitions without impacting concurrent access to data in other partitions. Furthermore, the utilities will run faster against

a single partition rather than against the entire tablespace and you will have more control over driving your utility workload. For example, you may not have sufficient time in the batch window to run a REORG on a four million page segmented tablespace, but you might have the time to run a REORG on one partition of that tablespace nightly. With four partitions of one million pages (or perhaps more partitions containing even fewer pages) you may be able to REORG one partition a night.

- Implement partitioned tablespaces to improve data availability. For example, if the data is partitioned by region, the partitions for the Eastern, Southern, and Northern regions can be made available while the Western region partition is being reorganized.
- Use partitioned tablespaces to improve recoverability. Once again, consider the ramifications if the data is partitioned by region. If an error impacts data for the Eastern region only, then only the Eastern partition needs to be recovered. The Southern, Northern, and Western regions can remain on-line, because they are not impacted by the problem in the Eastern region's data.
- Consider partitioned tablespaces to isolate specific data areas in dedicated datasets. If there are specific data 'hot spots' that have higher data modification and/or access activity, you may be able to improve application performance by isolating the 'hot spot' into a single partition that can be tuned for the specific type of application access.
- Use a simple tablespace *only* when you need to mix data from different tables on one page. Simple tablespaces will mix data from each table assigned to the tablespace on each tablespace page. A segmented tablespace will not because each segment in the segmented tablespace is assigned to a single table. If you have two tables that are very frequently joined you might consider loading them into a single simple tablespace, ensuring that each row loaded from the first table is immediately followed by all of the rows from the second table that will be joined to the first table. This can minimize I/O for retrieval. However, DB2 will not maintain this ordering when the data is changed, so this approach is generally useful only for static data.

PARTITIONING CONSIDERATIONS

DB2 can handle up to 254 partitions per tablespace. The actual limit on the number of partitions depends on the DSSIZE of the tablespace. Large tablespaces are those that specify the LARGE parameter or have a DSSIZE greater than 4GB. The LARGE parameter was introduced with V5; DSSIZE with V6. A large tablespace can have from 1 to 254 partitions. Non-LARGE tablespaces are limited to no more than 64 partitions, as are any tablespaces created in a version prior to DB2 V5.

For non-LARGE partitioned tablespaces, the number of partitions impacts the maximum size of the dataset partition as follows:

Number of partitions	Maximum dataset size
1 to 16	4 GB
17 to 32	2 GB
33 to 64	1 GB

Keep these limitations in mind as you design your partitioned tablespaces.

As a general rule-of-thumb try to define tablespace partitions such that no one partition is more than 20 percent larger than the next largest partition. This provides even growth, which eases DASD monitoring and provides approximately even data access requirements and utility processing times across partitions. This is not a hard-and-fast rule though, especially when dealing with 'hot spots'. The 'hot spot' partition may be much smaller than the other partitions, going against the idea of maintaining evenly distributed partitions. This is OK.

Deciding to use a partitioned tablespace is not as simple as merely determining the size of the table. In the early days of DB2, size was the primary consideration for choosing a partitioned tablespace. However, as DB2 has matured and the applications written using DB2 have been modernized, additional considerations will impact your partitioning decisions. Application-level details, such as data contention, performance requirements, degree of parallelism, and the volume of updates to columns in the partitioning index, must factor in the decision to use partitioned tablespaces.

Sometimes designers try to avoid partitioned tablespaces by dividing

a table into multiple tables, each with its own tablespace. This is not wise. Never attempt to avoid a partitioned tablespace by implementing several smaller tablespaces, each containing a subset of the total amount of data. When proceeding in this manner, the designer usually places separate tables into each of the smaller tablespaces. This almost always is a bad design decision because it introduces an uncontrolled and unneeded denormalization. Furthermore, when data that logically belongs in one table is separated into multiple tables, SQL operations to access the data as a logical whole are made needlessly complex. One example of this complexity is the difficulty in enforcing unique keys across multiple tables. Although partitioned tablespaces can introduce additional complexities into your environment, these complexities never outweigh those introduced by mimicking partitioning with several smaller, identical tablespaces. To clarify why this idea is bad, consider these two different ways of implementing a three 'partition' solution:

• The first, recommended, way is to create the table in a single partitioned tablespace with three partitions as follows:

CREATE DB DB_SAMP;

```
CREATE TABLESPACE TS SAMP IN DB SAMP
       ERASE NO NUMPARTS 3
       (PART 1
        USING STOGROUP SG_SAMP1
        PRIQTY 2000 SECQTY 50
        COMPRESS NO.
        PART 2
        USING STOGROUP SG_SAMP2
        PRIQTY 4000 SECQTY 150
        COMPRESS YES.
        PART 3
        USING STOGROUP SG SAMP3
        PRIQTY 1000
        SECOTY 50
        COMPRESS YES)
       LOCKSIZE PAGE BUFFERPOOL BP1 CLOSE NO;
```

```
CREATE TABLE TB_SAMP . . . IN DB_SAMP.TS_SAMP;
```

• The second, ill-advised, way is to create three tablespaces each with its own table as follows:

CREATE DB DB_SAMP;

CREATE	TABLESPACE TS_SAMP1 IN DB_SAMP USING STOGROUP SG_SAMP1 PRIQTY 2000 SECQTY 50 ERASE NO COMPRESS NO LOCKSIZE PAGE BUFFERPOOL BP1 CLOSE NO;
CREATE	TABLESPACE TS_SAMP2 IN DB_SAMP USING STOGROUP SG_SAMP2 PRIQTY 4000 SECQTY 150 ERASE NO COMPRESS YES LOCKSIZE PAGE BUFFERPOOL BP1 CLOSE NO;
CREATE	TABLESPACE TS_SAMP3 IN DB_SAMP USING STOGROUP SG_SAMP3 PRIQTY 1000 SECQTY 50 ERASE NO COMPRESS YES LOCKSIZE PAGE BUFFERPOOL BP1 CLOSE NO;
CREATE	TABLETB_SAMP1INDB_SAMP.TS_SAMP1;TABLETB_SAMP2INDB_SAMP.TS_SAMP2;TABLETB_SAMP3INDB_SAMP.TS_SAMP3;

Now consider how difficult it would be to retrieve data in the second implementation if you did not know which 'partition' the data resides in, or if the data could reside in multiple partitions.

Using the first example a simple SELECT will work:

SELECT *
FROM TB_SAMP
WHERE COL1 = :HOST-VARIABLE;

In the second example, a UNION is required:

```
SELECT *

FROM TB_SAMP1

WHERE COL1 = :HOST-VARIABLE

UNION ALL

SELECT *

FROM TB_SAMP2

WHERE COL1 = :HOST-VARIABLE

UNION ALL

SELECT *

FROM TB_SAMP3

WHERE COL1 = :HOST-VARIABLE;
```

If other tables need to be joined the 'solution' becomes even more

complex. Likewise if data must be updated, inserted, or deleted and you do not know which 'partition' contains the impacted data. The bottom line is: avoid bypassing DB2 partitioning using your own pseudo-partitions.

PARTITIONING PROS AND CONS

Before deciding to partition a tablespace, weigh the pros and cons. Consult the following list of advantages and disadvantages before implementation:

Advantages of a partitioned tablespace:

- Each partition can be placed on a different DASD volume to increase access efficiency.
- Partitioned tablespaces are the only type of tablespace that can hold more than 64GB of data (the maximum size of simple and segmented tablespaces). A partitioned tablespace with extended addressability (EA-enabled) can hold up to 16 terabytes of data. Without being EA-enabled a partitioned tablespace can store up to about 1TB of data.
- Start and stop commands can be issued at the partition level. By stopping only specific partitions, the remaining partitions are available to be accessed thereby promoting higher availability.
- Free space (PCTFREE and FREEPAGE) can be specified at the partition level enabling the DBA to isolate data 'hot spots' to a specific partition and tune accordingly.
- Partitioning can optimize Query I/O, CPU, and Sysplex parallelism by removing disk contention as an issue because partitions can be spread out across multiple devices.
- Tablespace scans on partitioned tablespaces can skip partitions that are excluded, based on the query predicates. Skipping entire partitions can improve overall query performance for tablespace scans because less data needs to be accessed.
- The clustering index used for partitioning can be set up to decrease data contention. For example, if the tablespace will be

partitioned by DEPT, each department (or range of compatible departments) could be placed in separate partitions. Each department is in a discrete physical dataset, thereby reducing inter-departmental contention because of multiple departments coexisting on the same data page. Note that contention remains for data in non-partitioned indexes (although this contention has been significantly reduced in recent versions of DB2).

- DB2 creates a separate compression dictionary for each tablespace partition. Multiple dictionaries tend to cause better overall compression ratios. In addition, it is more likely that the partition-level compression dictionaries can be rebuilt more frequently than non-partitioned dictionaries. Frequent rebuilding of the compression dictionary can lead to a better overall compression ratio.
- The REORG, COPY, and RECOVER utilities can execute on tablespaces at the partition level. If these utilities are set to execute on partitions instead of on the entire tablespace, valuable time can be saved by processing only the partitions that need to be reorganized, copied, or recovered. Partition independence and resource serialization further increase the availability of partitions during utility processing.

Disadvantages of a partitioned tablespace:

- Only one table can be defined in a partitioned tablespace. This is not necessarily a disadvantage because most DBAs follow a one-table-per-tablespace rule.
- The columns of the partitioning index cannot be updated. To change a value in one of these columns, you must delete the row and then reinsert it with the new values.
- The range of key values for which data will be inserted into the table must be known and stable before you create the partitioning index. To define a partition, a range of values must be hard-coded into the partitioning index definition. These ranges will be used to distribute the data throughout the partitions. If you provide a stop-gap partition to catch all the values lower (or higher) than the defined range, monitor that partition to ensure that it does not

grow dramatically or cause performance problems if it is smaller or larger than most other partitions.

• After you define the method of partitioning, the only way to change it is to ALTER the partitioning index to change the LIMITKEY values and reorganize any impacted partitions. Prior to V6 you had to drop and redefine both the partitioning index and tablespace to change LIMITKEY specifications.

In general, partitioned tablespaces are becoming more useful. You might even want to consider using partitioning for most tablespaces (instead of segmented), especially if parallelism is an issue. At least, consider partitioning tablespaces that are accessed in a read-only manner by long-running batch programs. Of course, very small tablespaces are rarely viable candidates for partitioning, even with DB2's advanced I/O, CPU, and Sysplex parallelism features. This is true because the smaller the amount of data to access, the more difficult it is to break it into pieces large enough such that concurrent, parallel processing will be helpful.

When using partitioned tablespaces, try to place each partition of the same partitioned tablespace on separate DASD volumes. Failure to do so can negatively affect the performance of query parallelism performed against those partitions. Disk drive head contention will occur because concurrent access is being performed on separate partitions that co-exist on the same device. Of course, with some of the newer storage devices, such as the ESS Shark hardware from IBM, dataset placement is a non-issue because of the way in which data is physically stored on the device.

SUMMARY

DB2 provides three different types of tablespaces, each of which has its own distinct set of advantages and disadvantages for use depending upon the situation. As a DBA you should understand the implementation details of each type of tablespace and be prepared to choose the right type of tablespace for each situation.

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Interpreting DB2 accounting trace using COBOL

There is no need to mention the wealth of accounting information present in the DB2 SMF accounting record (SMF 101). Many performance problems in an application can be quickly solved once the bottlenecks are identified. Performance monitoring tools help us identify bottlenecks, but what if a site does not have those expensive tools? What if a site wants to move to DB2 but does not want to budget for expensive tools right away. What if a site is planning to get rid of their performance tools to cut down costs? Sometimes it is difficult to justify the investment in performance monitoring tools because many sites using them do not use a lot of the information provided by the tools.

Here is a small utility program that would interpret the SMF 101 record (DB2 accounting record) and produce a report that could help analyse a performance problem or take proactive action. The best part is that the program is written in COBOL and so your program can easily be customized according to the requirements (populating a database for performance analysis, customizing reports etc).

SOFTWARE DEPENDENCIES

This program works well when run against the SMF 101 records from DB2 for OS/390 Version 6 and with the trace destination as SMF. However if your site has some other version of DB2, then you need to check the layouts of DSNDQWA0, DSNDQWHS, DSNDQWAC, DSNDQLAC, and DSNDQXST from the corresponding SDSNMACS library shipped with your version of DB2. If these tally with the layout used by the program (the program uses layouts from the SDSNMACS library shipped with Version 6), it can be used as is, otherwise you need to make modifications to the layout to make it work with your version of DB2.

The program takes SMF101 records as input. It is necessary to have the trace destination as SMF and not GTF because the program processes the SMF101 header information. If the trace destination is GTF then minor changes are required to the length and the header layout definitions.

GETTING STARTED

Before the program can produce any report, DB2 accounting trace (Classes 1, 2, and 3) should be activated with the destination set to SMF (this is the default destination for accounting trace). Additional constraints like PLAN, AUTHID can be applied to restrict the kind of data collected by the trace. The program also makes use of an additional variable, CUTOFF-CPU-TIME, which is currently set to 1.00. The value assigned to this variable is the time in seconds and is used to report only those threads/connections where the 'in DB2 CPU time' is greater than or equal to values assigned to CUTOFF-CPU-TIME.

This helps us to eliminate a lot of output that would otherwise get generated as a result of DB2 Command executions (say like *-display thread*, *-start procedure*, *-stop procedure*, etc) and like SQL statements that may not be of any interest from a monitoring point of view. You can set the value of this variable as per your requirements.

The program also calls an Assembler routine to request storage for the linkage section blocks. These few lines request a GETMAIN and pass the pointer to the calling program. This Assembler code is listed in stream with the JCL to compile and link in *Appendix A*.

If one does not want to use any Assembler code then an alternative way would be to have a COBOL program call our COBOL program so that the caller requests linkage section storage.

GENERATING ACCOUNTING REPORT

This program processes SMF 101 records. An SMF dataset containing SMF records, identified by DDname SMFDATA, is the input to the program. To have better control over the SMF records to be processed, the IBM-supplied utility IFASMFDP can be used. This utility can dump records of interest (SMF Type 101) corresponding to a particular date/time interval. Using this will also be useful if we want to collect SMF data for a few days to analyse, because less resource (both DASD and CPU) would be needed to generate accounting information because of the reduced volume of data to work with. See *Appendix B* for the required JCL.

REPORTS

The job produces two sets of output that can be used to analyse the accounting information. DD cards SYSOUT and BPACCT identify the output destinations.

The first set of reports identified by DD card SYSOUT gives us information about DB2 elapsed and CPU times, DB2 I/O done, I/O wait time within DB2, RDS accounting information like number of DML statements executed by each connection, and DDF accounting information.

The second set of information identified by DD card BPACCT tells us about Buffer Manager accounting. This gives us the breakdown of Get Pages issued by a connection on each of the buffer pools it uses. It also indicates synchronous read/write activity if any.

Sample reports are shown in *Appendix C*. Output is shown here as wrapped because of paper size limitations. The output generated by the program does not wrap over and so it can easily be taken to a spreadsheet for further analysis if required.

Fields on the report include:

- Date date on which the accounting record was written (format YYDDD).
- Time time at which the accounting record was written (format HH:MM:SS).
- Jobname for batch this is job name, for TSO it is the log-on ID, and for CICS it is the user-id specified on the RCT.
- Connect connection type (batch, TSO, utility, CICS).
- Plan application plan name.
- DB2ELAP elapsed time in DB2 (in seconds).
- DB2CPU in DB2 CPU time (in seconds).
- DB2IO total number of get pages requested by connection.
- IO WAIT accumulated I/O elapsed wait time (in seconds).

- Lock Wait accumulated wait time due to latch or local lock contention.
- #Selects number of SELECT statements executed under the connection.
- #Inserts number of INSERT statements executed under the connection.
- #Updates number of UPDATE statements executed under the connection.
- #Deletes number of DELETE statements executed under the connection.
- #Opens number of OPEN cursor statements executed under the connection.
- #Fetches number of FETCH requests executed under the connection.
- #Closes number of CLOSE cursor statements executed under the connection.
- #Prepares number of PREPARE statements executed under the connection.
- Remote Location IP address of the remote location.
- SQL-S SQLs sent from the server.
- SQL-R SQLs received from the client.
- ROWS SENT rows sent to the client.
- BYTES-S number of bytes sent from the server to the client.
- BYTES-R number of bytes received from the client.
- BPID bufferpool identifier (BP0, BP1, etc).
- #GETP number of get pages for the bufferpool.
- SYNC-Read number of synchronous reads.
- SYNC-Write number of synchronous writes.

The elapsed and CPU times shown in the report should not be used for chargebacks because the time interval is from thread connection to thread termination. It does not account for the time spent by an application before the connection or after the termination.

The program reports only three buffer pools, namely BP0, BP1, and BP2, but can be customized to report other/all buffer pools in a subsystem.

APPENDIX A

//***** job	card ****		
//COMPILE EXE	EC PGM=ASMA9Ø,PARM='OBJECT,NODECK'		
//SYSLIB DE	DISP=SHR,DSN=SYS1.MACLIB		
//SYSLIN DE	<pre>DSN=&&LOADSET(GETSTORG),DISP=(NEW,PASS),</pre>		
//	UNIT=SYSDA,SPACE=(800,(50,50,2)),DCB=(BLKSIZE=800)		
//SYSPRINT DE	SYSOUT=*		
//SYSUDUMP DE	SYSOUT=*		
//SYSUT1 DE			
//SYSUT2 DE			
//SYSIN DE			
	TINE TO GET STORAGE		
GETSTORG CSEC			
R1 EQU			
R2 EQU			
R3 EQU			
R4 EQU			
R5 EQU			
R6 EQU R7 EQU			
R8 EQU			
R9 EQU			
R1Ø EQU			
R11 EQU			
R12 EQU			
R13 EQU			
R14 EQU			
R15 EQU			
* PARMS: L	ENGTH OF STORAGE (4-BYTES BINARY - PIC S9(8) COMP)		
	PDATED ADDRESS (4-BYTES BINARY - USAGE POINTER)		
STM F	14,R12,12(R13) SAVE REGS		
LR F	12,R15 LOAD BASE REG		
USING (ETSTORG,R12 ADDRESSABILITY		
LR F	15,R13 OLD SAVE AREA ADDRESS		
	13, SAVEAREA POINT TO NEW SAVE AREA		
	15,4(R13) PUT OLD SAVE ADDR IN IT		
	2,Ø(R1) GET ADDRESS OF STORAGE LENGTH		
ICM F	2,15,Ø(R2) GET ACTUAL STORAGE LENGTH		

L R3,4(R1) GET ADDRESS OF AREA TO UPDATE GETMAIN EU, LV=(2), A=(3) GET STORAGE L R13,4(R13) GET OLD SAVE AREA ADDRESS LM R14,R12,12(R13) RESTORE REGS LA R15,Ø SET RC=Ø LA R15,Ø BR R14 RETURN SAVEAREA DC 18F'Ø' SAVE AREA LTORG END GETSTORG //LKED EXEC PGM=IEWL,PARM='LIST,XREF,LET,RENT,AMODE=24,RMODE=24', 11 COND = (4, LT)//SYSLIN DD DSN=&&LOADSET(GETSTORG),DISP=(OLD,DELETE) //SYSLMOD DD DSN=XXXXX.LOADLIB(GETSTORG),DISP=SHR //SYSPRINT DD SYSOUT=* //SYSUDUMP DD SYSOUT=* //SYSUT1 DD UNIT=SYSDA,SPACE=(1024,(50,50))

APPENDIX B

//***** job card ****
//IFASMFDP EXEC PGM=IFASMFDP
//SYSPRINT DD SYSOUT=*
//SYSUDUMP DD SYSOUT=*
//DUMPIN DD DISP=SHR,DSN=INPUT.SMFDATA
<pre>//DUMPOUT DD DSN=OUTPUT.SMF101.D0131,DISP=(NEW,CATLG),</pre>
<pre>// DCB=(LRECL=32760, RECFM=VBS), SPACE=(CYL, (1,5))</pre>
//SYSIN DD *,DCB=BLKSIZE=80
INDD(DUMPIN,OPTIONS(DUMP))
OUTDD(DUMPOUT,TYPE(101))
DATE(2001031,2001031)
START(Ø8ØØ)
END(1800)
/*
//* SMF1Ø1
//SMF1Ø1 EXEC PGM=SMF1Ø1
//STEPLIB DD DSN=XXXXX.LOADLIB,DISP=SHR
//SYSTSPRT DD SYSOUT=*
//BPACCT DD SYSOUT=*,DCB=LRECL=85
//SMFDATA DD DSN=OUTPUT.SMF101.D0131,DISP=SHR,
// DCB=(LRECL=3276Ø,RECFM=VBS)
//SYSPRINT DD SYSOUT=*
//SYSOUT DD SYSOUT=*,DCB=LRECL=220
//SYSUDUMP DD SYSOUT=*

APPENDIX C

MVS System:OSSF DB2 Sub-System:DB2P Time in Seconds _____ Date Time Jobname Connect Plan DB2 ELAP DB2 CPU _____ Ø1Ø31 19:24:35 \$TSTDL11 UTILITY DSNUTIL 548.73 40.72 Ø1Ø31 20:04:19 \$TST5113 BATCH PITAPLN1 1621.33 244.87 Ø1Ø31 2Ø:Ø5:16 \$TST5113 UTILITY DSNUTIL 32.3Ø 11.6Ø 1.19 Ø1Ø31 20:06:31 \$TST5114 UTILITY DSNUTIL 9.72 Ø1Ø31 20:22:53 SQRW.EXE SERVER DISTSERV 150.52 54.61 _____ DB2I0 IO WAIT Lock Wait#Selects #Inserts #Updates #Deletes _____ 36767 7.24 Ø.ØØ Ø Ø Ø Ø 785845 51.23 Ø.Ø5 16Ø47 1Ø3626 Ø Ø 1.04 Ø.00 Ø Ø 271Ø9 Ø Ø 656 Ø.Ø3 Ø.ØØ 2707 Ø.82 Ø.ØØ Ø Ø Ø Ø Ø Ø Ø Ø #Opens #Fetches #Closes #Prepares Ø Ø Ø Ø 1 Ø 1 14435 Ø Ø Ø Ø Ø Ø Ø Ø 3 3 15 6 Remote Location SQL-S SQL-R ROWS-S BYTES-S BYTES-R - - - - - - - - - - ------- - - - - - - - - - -Ø 15 13 9368 172.10.12.140 9216 Date Time Jobname Connect Plan BPID #GETP SYNC-Read SYNC-Write _____ Ø1Ø31 19:24:35 \$TSTDL11 UTILITY DSNUTIL BPØ 17Ø Ø Ø Ø1Ø31 19:24:35 \$TSTDL11 UTILITY DSNUTIL BP1 Ø 3196 Ø Ø1Ø31 19:24:35 \$TSTDL11 UTILITY DSNUTIL BP2 334Ø1 Ø Ø Ø1Ø31 20:04:19 \$TST5113 BATCH PITAPLN1 BPØ 1575 3 Ø Ø1Ø31 20:04:19 \$TST5113 BATCH PITAPLN1 BP1 98889 2252 Ø Ø1Ø31 2Ø:Ø4:19 \$TST5113 BATCH PITAPLN1 BP2 685381 165 Ø Ø1Ø31 20:05:16 \$TST5113 UTILITY DSNUTIL BPØ 183 Ø Ø Ø Ø1Ø31 20:05:16 \$TST5113 UTILITY DSNUTIL BP1 4889 Ø Ø1Ø31 20:05:16 \$TST5113 UTILITY DSNUTIL BP2 22Ø37 Ø Ø Ø1Ø31 20:06:31 \$TST5114 UTILITY DSNUTIL BPØ 1Ø6 Ø Ø Ø Ø1Ø31 20:06:31 \$TST5114 UTILITY DSNUTIL BP1 55Ø Ø Ø1Ø31 20:22:53 SORW.EXE SERVER DISTSERV BPØ 183 1 Ø 1456 45 Ø1Ø31 2Ø:22:53 SQRW.EXE SERVER DISTSERV BP1 Ø

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SMF101.COB

```
***
        PROGRAM FOR PROCESSING ACCOUNTING TRACE
IDENTIFICATION DIVISION.
Program-Id. SMF101.
Author.
           Pranav Sampat.
ENVIRONMENT DIVISION.
*
INPUT-OUTPUT SECTION.
FILE-CONTROL.
    SELECT SMFDATA ASSIGN TO SMFDATA
    ORGANIZATION SEQUENTIAL.
    SELECT BPACCT ASSIGN TO BPACCT
    ORGANIZATION SEQUENTIAL.
+
DATA DIVISION.
FILE SECTION.
* SMF INPUT: QSAM, RECFM=VBS, BLKSIZE=4096 LRECL=32760
FD SMFDATA
    LABEL RECORDS ARE STANDARD
    RECORDING MODE IS S
     BLOCK CONTAINS 4096 CHARACTERS
    RECORD CONTAINS 18 TO 32756 CHARACTERS.
                     PIC X(18).
Ø1 SMFDATA-MIN-RECORD
Ø1 SMFDATA-MAX-RECORD PIC X(32756).
FD BPACCT
    LABEL RECORDS ARE STANDARD.
Ø1 BPACCT-REC
                      PIC X(8Ø).
WORKING-STORAGE SECTION.
*
Ø1
   WS-BPACCT-REC.
                      PIC 99999.
    Ø5 F-STRDATE
    Ø5 FILLER
                       PIC X VALUE SPACE.
    Ø5 F-STRTIME.
      10 F-STRTIME-HH PIC 99.
                      PIC X VALUE ':'.
      1Ø FILLER
      10 F-STRTIME-MM PIC 99.
      1Ø FILLER
                      PIC X VALUE ':'.
      10 F-STRTIME-SS PIC 99.
    Ø5 FILLER
                      PIC X VALUE SPACE.
    Ø5 F-JOBNAME
                      PIC X(8).
    Ø5 FILLER
                       PIC X VALUE SPACE.
    Ø5 F-CONNECT
                       PIC X(8).
                      PIC X VALUE SPACE.
    Ø5 FILLER
    Ø5 F-PLAN
                       PIC X(8).
    Ø5 FILLER
                      PIC X.
```

Ø5 BPID PIC X(6). Ø5 NUM-GETP PIC ZZZZZZ9. Ø5 FILLER PIC X. PIC ZZZZZZ9. Ø5 NUM-SYNC-READ Ø5 FILLER PIC X. Ø5 NUM-SYNC-WRITE PIC ZZZZZZ9. * Ø1 WS-SMF-RECNO PIC 9(6) COMP VALUE Ø. PIC 9(8) COMP VALUE 32768 SYNC. Ø1 BIN4-32768 Ø1 SMFØØ-PNTR USAGE IS POINTER SYNC. Ø1 SMFØØ-ADDRVAL REDEFINES SMFØØ-PNTR PIC S9(9) COMP. Ø1 SMF1Ø1RPS-PNTR USAGE IS POINTER SYNC. Ø1 SMF1Ø1RPS-ADDRVAL REDEFINES SMF1Ø1RPS-PNTR PIC S9(9) COMP. Ø1 DSNDOWHS-PNTR USAGE IS POINTER SYNC. DSNDQWHS-ADDRVAL REDEFINES DSNDQWHS-PNTR PIC S9(9) COMP. Ø1 SYNC. Ø1 DSNDQWAC-PNTR USAGE IS POINTER DSNDQWAC-ADDRVAL REDEFINES DSNDQWAC-PNTR Ø1 PIC S9(9) COMP. USAGE IS POINTER Ø1 DSNDQXST-PNTR SYNC. Ø1 DSNDQXST-ADDRVAL REDEFINES DSNDQXST-PNTR PIC S9(9) COMP. USAGE IS POINTER Ø1 DSNDQBAC-PNTR SYNC. DSNDQBAC-ADDRVAL REDEFINES DSNDQBAC-PNTR PIC S9(9) COMP. Ø1 USAGE IS POINTER Ø1 DSNDQLAC-PNTR SYNC. DSNDQLAC-ADDRVAL REDEFINES DSNDQLAC-PNTR PIC S9(9) COMP. Ø1 END-OF-FILE-FLAG PIC X VALUE 'N'. Ø1 88 END-OF-FILE VALUE 'Y'. Ø1 BIN2-3Ø PIC 9(4) COMP VALUE 101. Ø1 FILLER REDEFINES BIN2-30. Ø2 FILLER PIC X(1). Ø2 BIN1-1Ø1 PIC X(1). WS-VARS. Ø1 PIC X(8). Ø2 PROGRAM-NAME CUTOFF-CPU-TIME PIC 9V99 VALUE 1.00. Ø2 Ø2 PIC 9(15)V99. TOT-INDB2-ELAP PIC 9(15)V99. Ø2 TOT-DB2-CPU Ø2 TOT-IOWT PIC 9(15)V99. Ø2 TOT-LWT PIC 9(15)V99. Ø2 T0T-I0 PIC 9(15)V999. Ø2 TEMP-HOURS PIC 9(9) COMP. Ø2 TEMP-MIN PIC 9(9) COMP. Ø2 TEMP-SECS PIC 9(9) COMP. Ø2 WS-CNT PIC 9(4) COMP. Ø2 WS-NUM-SELECTS PIC 9(9) COMP. PIC 9(9) Ø2 WS-NUM-INSERTS COMP. Ø2 WS-NUM-UPDATES PIC 9(9) COMP. Ø2 WS-NUM-DELETES PIC 9(9) COMP. Ø2 WS-NUM-OPENS PIC 9(9) COMP. Ø2 WS-NUM-FETCHES PIC 9(9) COMP. PIC 9(9) Ø2 WS-NUM-CLOSES COMP. Ø2 WS-NUM-PREPARES PIC 9(9) COMP. Ø1 WS-HYPENS PIC X(214) VALUE ALL '-'.

Ø1 OUT-REC-HEAD1.	
Ø5 FILLER	PIC X(11) VALUE 'MVS System:'.
Ø5 SYSTEM	PIC X(4).
Ø5 FILLER	PIC X(5) VALUE SPACE.
Ø5 FILLER	PIC X(15) VALUE 'DB2 Sub-System:'.
Ø5 SUBSYS	PIC X(4).
Ø5 FILLER	PIC X(20) VALUE ' Time in Seconds'.0
Ø1 OUT-REC-HEADER.	
Ø5 FILLER	PIC X(5) VALUE 'Date '.
Ø5 FILLER	PIC X(9) VALUE ' Time '.
Ø5 FILLER	PIC X VALUE SPACE.
Ø5 FILLER	PIC X(9) VALUE 'Jobname'.
Ø5 FILLER	PIC X(9) VALUE 'Connect'.
Ø5 FILLER	PIC X(10) VALUE 'Plan '.
Ø5 FILLER	PIC X(10) VALUE 'DB2 ELAP'.
Ø5 FILLER	PIC X(10) VALUE 'DB2 CPU'.
Ø5 FILLER	PIC X(9) VALUE 'DB2IO'.
Ø5 FILLER	PIC X(9) VALUE 'IO WAIT'.
Ø5 FILLER	PIC X(9) VALUE 'Lock Wait'.
Ø5 FILLER	PIC X(9) VALUE '#Selects'.
Ø5 FILLER	PIC X(9) VALUE '#Inserts'.
Ø5 FILLER	PIC X(9) VALUE '#Updates'.
Ø5 FILLER	PIC X(9) VALUE '#Deletes'.
Ø5 FILLER	PIC X(8) VALUE '#Opens'.
Ø5 FILLER	PIC X(9) VALUE '#Fetches'.
Ø5 FILLER	PIC X(8) VALUE '#Closes'.
Ø5 FILLER	PIC X(1Ø) VALUE '#Prepares'.
Ø5 FILLER	PIC X(16) VALUE 'Remote Location'.
Ø5 FILLER	PIC X(7) VALUE 'SQL-S'.
Ø5 FILLER	PIC X(7) VALUE 'SQL-R'.
Ø5 FILLER	PIC X(8) VALUE 'ROWS-S'.
Ø5 FILLER	PIC X(8) VALUE 'BYTES-S'.
Ø5 FILLER	PIC X(8) VALUE 'BYTES-R'.
Ø1 BPACCT-REC-HEADER.	
Ø5 FILLER	PIC X(8) VALUE 'Date '.
Ø5 FILLER	PIC X(8) VALUE 'Time '.
Ø5 FILLER	PIC X(8) VALUE 'Jobname'.
Ø5 FILLER	PIC X(9) VALUE 'Connect'.
Ø5 FILLER	PIC X(9) VALUE 'Plan '.
Ø5 FILLER	PIC X(8) VALUE 'BPID'.
Ø5 FILLER	PIC X(7) VALUE '#GETP'.
Ø5 FILLER	PIC X(1Ø) VALUE 'SYNC-Read'.
Ø5 FILLER	PIC X(1Ø) VALUE 'SYNC-Write'.
Ø1 OUT-REC-DATA.	
Ø5 STRDATE	PIC 99999.
Ø5 FILLER	PIC X VALUE SPACE.
Ø5 STRTIME.	
10 STRTIME-HH	PIC 99.
10 FILLER	PIC X VALUE ':'.
10 STRTIME-MM	PIC 99.
1Ø FILLER	PIC X VALUE ':'.

	1Ø STRTIME-SS	DIC	00
øб			X VALUE SPACE.
			X(8).
			X VALUE SPACE.
			X(8).
			X VALUE SPACE.
			X(8).
			X VALUE SPACE.
			ZZZZZ9.99.
		PIC	
			ZZZZZ9.99.
		PIC	
			ZZZZZZ9.
		PIC	
			ZZZZZ9.99.
		PIC	
05 05	DB2LOCKWAIT FILLER	PIC	¥
			ZZZZZZZ9.
		PIC	
	NUM-INSERTS		
		PIC	
	NUM-UPDATES		
		PIC	
			ZZZZZZZ9.
		PIC	
			ZZZZZZ9.
		PIC	
	NUM-FETCHES		
		PIC	
			ZZZZZZ9.
		PIC	
	NUM-PREPARES		
	FILLER	PIC	
	DDF-REC.		
20		PIC	X(16).
	10 FILLER	PIC	
			ZZZZ9.
		PIC	
	1Ø SQLS-RECEIVED	PIC	ZZZZ9.
	10 FILLER	PIC	
	10 ROWS-SENT		ZZZZZZ9.
	1Ø FILLER	PIC	Х.
			ZZZZZZ9.
		PIC	
	10 BYTES-RECEIVED	PIC	ZZZZZZ9.
			M ID – MVS – */
	SUBSYS SU	JBSYS	S ID – DB2 – */
	STRDATE DA	ATE S	SMF REC WRITTEN */
	STRTIME T	IME S	SMF REC WRITTEN */

*

* * *

* JOBNAME JOB NAME. * FOR BATCH JOB NAME, * FOR TSO. QMF TSO LOGON ID. * FOR CICS USER-ID SPECIFIED ON RCT * CONNECT CONNECTION TYPE * FOR BATCH 'BATCH', * FOR TSO 'TSO', * FOR QMF 'DB2CALL * FOR UTILITY 'UTILITY * FOR CICS CICS-ID * PLAN APPLICATION PLAN NAM * LINKAGE SECTION. ** GENERIC SMF RECORD. Ø1 SMFØØ-RECORD. PIC X(1). Ø2 SMFØØFLG Ø2 SMFØØRTY PIC X(1). PIC 9(9) COMP. Ø2 SMFØØTME Ø2 SMFØØDTE PIC 9(7) COMP-3. Ø2 SMFØØSID PIC X(4). Ø2 SMFØØSSI PIC X(4). Ø2 SMFØØSTP PIC 9(4) COMP. Ø2 FILLER PIC X(32743). Ø1 SMF1Ø1RPS. ** HEADER FOR TYPE 101 RECORD Ø2 SMF1Ø1RHD. Ø3 SMF1Ø1FLG PIC X(1). Ø3 SMF1Ø1RTY PIC X(1). Ø3 SMF1Ø1TME PIC 9(9) COMP. Ø3 SMF1Ø1DTE PIC 9(7) COMP-3. PIC X(4). Ø3 SMF1Ø1SID Ø3 SMF1Ø1SSI PIC X(4). Ø3 FILLER PIC X(6). ** SELF-DEFINING SECTION OF TYPE 101 RECORD Ø2 DSNDQWAØ. Ø3 OWAØ1PSO PIC 9(9) COMP. Ø3 QWAØ1PSL PIC 9(4) COMP. Ø3 QWAØ1PSN PIC 9(4) COMP. Ø3 QWAØ1R10 PIC 9(9) COMP. Ø3 QWAØ1R1L PIC 9(4) COMP. Ø3 QWAØ1R1N PIC 9(4) COMP. Ø3 QWAØ1R20 PIC 9(9) COMP. PIC 9(4) COMP. Ø3 QWAØ1R2L Ø3 QWAØ1R2N PIC 9(4) COMP. Ø3 OWAØ1R30 PIC 9(9) COMP. PIC 9(4) COMP. Ø3 QWAØ1R3L Ø3 QWAØ1R3N PIC 9(4) COMP. Ø3 QWAØ1R40 PIC 9(9) COMP. Ø3 QWAØ1R4L PIC 9(4) COMP.

	 Ø3 QWAØ1R4N Ø3 QWAØ1R50 Ø3 QWAØ1R5L Ø3 QWAØ1R5N Ø3 QWAØ1R60 Ø3 QWAØ1R6L Ø3 QWAØ1R6N Ø3 QWAØ1R70 Ø3 QWAØ1R71 Ø3 QWAØ1R7N Ø3 QWAØ1R80 Ø3 QWAØ1R8L Ø3 QWAØ1R8L Ø3 QWAØ1R8N Ø3 QWAØ1R90 Ø3 QWAØ1R9L Ø3 QWAØ1R9N 	PIC 9(9) COMP. PIC 9(4) COMP. PIC 9(4) COMP. PIC 9(9) COMP. PIC 9(4) COMP. PIC 9(4) COMP. PIC 9(4) COMP. PIC 9(4) COMP. PIC 9(4) COMP. PIC 9(9) COMP.
*	*****	*****
	T DATA SECTION	
	MENTATION STANDA	RD HFADER

*		
Ø1 DSN	DQWHS.	
Ø2	QWHSLEN	PIC 9(4) COMP.
Ø2	QWHSTYP	PIC X(1).
Ø2	QWHSRMID	PIC X(1).
Ø2	QWHSIID	PIC 9(4) COMP.
Ø2 Ø2	QWHSNSDA QWHSSRN	PIC X(1).
ø2 Ø2	QWHSSRN	PIC X(1). PIC 9(9) COMP.
Ø2 Ø2	QWHSSSID	PIC X(4).
Ø2	QWHSSTCK.	110 X(4).
52	Ø3 QWHSSC1	PIC 9(9) COMP.
	Ø3 QWHSSC2	PIC 9(9) COMP.
Ø2	QWHSISEQ	PIC 9(9) COMP.
Ø2	QWHSWSEQ	PIC 9(9) COMP.
Ø2	QWHSMTN	PIC X(4).
Ø2	QWHSLOCN	PIC X(16).
Ø2	QWHSNID	PIC X(8).
Ø2	QWHSLUNM	PIC X(8).
Ø2	QWHSLUUV	PIC X(6).
Ø2	QWHSLUCC	PIC 9(4) COMP.
	TION HEADER	
Ø2 Ø2	QWHCLEN QWHCTYP	PIC 9(4) COMP. PIC X.
Ø2 Ø2	FILLER	PIC X. PIC X.
Ø2 Ø2	QWHCAID	PIC X. PIC X(8).
Ø2	QWHCCV	PIC X(12).
Ø2	QWHCCN	PIC X(8).
ø2	QWHCPLAN	PIC X(8).
ø2	QWHCOPID	PIC X(8).

**** ACCOUNTING DATA Ø1 DSNDQWAC. Ø2 OWACBSC pic x(8). Ø2 QWACESC pic x(8). Ø2 QWACBJST PIC 9(16) COMP. Ø2 OWACEJST PIC 9(16) COMP. Ø2 OWACBSRB PIC 9(16) COMP. Ø2 QWACESRB PIC 9(16) COMP. PIC 9(9) COMP. Ø2 OWACRINV Ø2 QWACNID PIC X(16). Ø2 QWACCOMM PIC 9(9) COMP. Ø2 OWACABRT PIC 9(9) COMP. Ø2 QWCA1. Ø3 QWACASC PIC 9(16) COMP. Ø3 QWACAJST PIC 9(16) COMP. Ø3 QWACASRB PIC 9(16) COMP. Ø3 OWACAWTI PIC 9(16) COMP. Ø3 QWACAWTL PIC 9(16) COMP. Ø3 QWACARNA PIC 9(9) COMP. Ø3 OWACARNE PIC 9(9) COMP. *** RDS STATISTICS BLOCK Ø1 DSNDQXST. Ø2 OXHEAD. PIC XX. Ø3 QXID PIC 9(4) COMP. Ø3 QXLEN Ø3 QXEYE PIC X(4). Ø2 QXSTATS. Ø3 OXSELECT PIC 9(9) COMP. PIC 9(9) COMP. Ø3 OXINSRT Ø3 QXUPDTE PIC 9(9) COMP. Ø3 QXDELET PIC 9(9) COMP. Ø3 QXDESC PIC 9(9) COMP. Ø3 OXPREP PIC 9(9) COMP. Ø3 QXOPEN PIC 9(9) COMP. PIC 9(9) COMP. Ø3 QXCLOSE Ø3 QXCRTAB PIC 9(9) COMP. Ø3 QXCRINX PIC 9(9) COMP. Ø3 QXCTABS PIC 9(9) COMP. Ø3 OXCRSYN PIC 9(9) COMP. PIC 9(9) COMP. Ø3 QXCRDAB Ø3 QXCRSTG PIC 9(9) COMP. Ø3 OXDEFVU PIC 9(9) COMP. PIC 9(9) COMP. Ø3 QXDRPIX Ø3 QXDRPTA PIC 9(9) COMP. Ø3 QXDRPTS PIC 9(9) COMP. Ø3 QXDRPDB PIC 9(9) COMP.

```
Ø3 OXDRPSY
                              PIC 9(9) COMP.
     Ø3 OXDRPST
                              PIC 9(9) COMP.
     Ø3 QXDRPVU
                             PIC 9(9) COMP.
                             PIC 9(9) COMP.
     Ø3 QXALTST
     Ø3 OXFETCH
                              PIC 9(9) COMP.
*** BUFFER MANAGER ACCOUNTING BLOCK
Ø1 DSNDOBAC.
   Ø2 QBACPID PIC 9(9) COMP.
   Ø2 QBACGET PIC 9(9) COMP.
   Ø2 OBACBPX PIC 9(9) COMP.
   Ø2 QBACSWS PIC 9(9) COMP.
   Ø2 QBACSWU PIC 9(9) COMP.
   Ø2 QBACRIO PIC 9(9) COMP.
   Ø2 QBACSEQ PIC 9(9) COMP.
   Ø2 QBACIMW PIC 9(9) COMP.
   Ø2 OBACLPF PIC 9(9) COMP.
   Ø2 QBACDPF PIC 9(9) COMP.
*
*
                        BUFFER POOL ID
         OBACPID
*
                        # OF GET PAGE ISSUED
         OBACGET
                        # OF BUFFER POOL EXPANSIONS required
         QBACBPX
*
                       # OF SETW ISSUED FOR SYSTEM pages
         OBACSWS
*
                        # OF SETW ISSUED FOR UW PAG
         QBACSWU
         QBACRIO
                       # OF SYNCHRONOUS READ I/O
*
         OBACSEO
                       # OF SEQ PREFETCH REQUESTED
                       # OF SYNCHRONOUS WRITE I/O
*
         QBACIMW
                       # OF LIST PREFETCH REQUEST
*
         OBACLPF
*
         QBACDPF
                       # OF DYNAMIC PREFETCH RQST
*** DISTRIBUTED DATA FACILITY ACCOUNTING BLOCK
Ø1 DSNDQLAC.
   Ø2 QLACLOCN PIC X(16).
   Ø2 QLACSQLS PIC 9(9) COMP.
   Ø2 QLACSQLR PIC 9(9) COMP.
   Ø2 QLACROWS PIC 9(9) COMP.
   Ø2 QLACROWR PIC 9(9) COMP.
   Ø2 QLACBYTS PIC 9(9) COMP.
   Ø2 QLACBYTR PIC 9(9) COMP.
   Ø2 QLACCNVS PIC 9(9) COMP.
*
                 - LOCATION
*
     QLACLOCN
*
                - NUMBER OF SOL STATEMENTS SENT
     OLACSOLS
*
               - NUMBER OF SQL STATEMENTS RECEIVED
     QLACSQLR
*
     OLACROWS
                - NUMBER OF ROWS SENT
*
     QLACROWR
                - NUMBER OF ROWS RECEIVED
*
               - NUMBER OF BYTES SENT
     QLACBYTS
```

```
*
     OLACBYTR
               - NUMBER OF BYTES RECEIVED
PROCEDURE DIVISION.
ØØØØ-MAINLINE.
    PERFORM 1000-INITIALIZE
    PERFORM UNTIL END-OF-FILE
      READ SMFDATA INTO SMFØØ-RECORD
        AT END MOVE 'Y' TO END-OF-FILE-FLAG
      END-READ
      IF NOT END-OF-FILE AND
         SMFØØRTY = BIN1-101
           ADD 1 TO WS-SMF-RECNO
           PERFORM 2000-TYPE101-PROCESS
           PERFORM 6000-DISPLAY-ACCOUNTING-INFO
      END-IF
    END-PERFORM.
    CLOSE SMFDATA.
    CLOSE BPACCT.
    GOBACK.
1000-INITIALIZE.
***** Request Storage for the linkage section
                                                *****
CALL 'GETSTORG' USING BIN4-32768 SMFØØ-PNTR.
     SET ADDRESS OF SMFØØ-RECORD TO SMFØØ-PNTR.
     OPEN INPUT SMFDATA.
     OPEN OUTPUT BPACCT.
     WRITE BPACCT-REC
                     FROM BPACCT-REC-HEADER
     WRITE BPACCT-REC
                      FROM WS-HYPENS.
 2000-TYPE101-PROCESS.
** THE FOLLOWING CODE ADDRESSES THE HEADER AND SELF-DEFINING
** SECTIONS
MOVE SMFØØ-ADDRVAL
                          TO SMF1Ø1RPS-ADDRVAL.
     SET ADDRESS OF SMF1Ø1RPS TO SMF1Ø1RPS-PNTR.
***
    4 IS SUBTRACTED AS RDW IS NOT INCLUDED IN OUR RECORD
     COMPUTE DSNDQWHS-ADDRVAL = SMF1Ø1RPS-ADDRVAL + QWAØ1PSO - 4
     SET ADDRESS OF DSNDQWHS TO DSNDQWHS-PNTR.
** THE FOLLOWING CODE ADDRESSES THE ACCOUNTING DATA
     MOVE SMF1Ø1SID
                    TO SYSTEM
     MOVE SMF1Ø1SSI
                    TO SUBSYS
     MOVE SMF1Ø1DTE
                    TO STRDATE , F-STRDATE
     PERFORM 2100-CONVERT-TIME
                   TO STRTIME-HH . F-STRTIME-HH
     MOVE TEMP-HOURS
     MOVE TEMP-MIN
                    TO STRTIME-MM , F-STRTIME-MM
     MOVE TEMP-SECS
                    TO STRTIME-SS , F-STRTIME-SS
     MOVE QWHCCN
                    TO CONNECT , F-CONNECT
     MOVE QWHCCV
                    TO JOBNAME
                                , F-JOBNAME
```

MOVE QWHCPLAN TO PLAN , F-PLAN MOVE Ø TO TOT-INDB2-ELAP , TOT-DB2-CPU . TOT-IOWT . TOT-IWT COMPUTE DSNDQWAC-ADDRVAL = SMF101RPS-ADDRVAL + QWA01R10 - 4SET ADDRESS OF DSNDQWAC TO DSNDQWAC-PNTR PERFORM 3000-COMPUTE-TIMES QWA01R1N TIMES COMPUTE TOT-INDB2-ELAP =(TOT-INDB2-ELAP / 4096000000.0) + 0.005 COMPUTE TOT-DB2-CPU = (TOT-DB2-CPU / 4096000000.0) + 0.005 COMPUTE TOT-IOWT = (TOT-IOWT / 4096000000.0) + 0.005COMPUTE TOT-LWT = (TOT-LWT / 4096000000.0) + 0.005MOVE TOT-INDB2-ELAP TO DB2ELAP MOVE TOT-DB2-CPU TO DB2CPU MOVE TOT-IOWT TO DB2IOWAIT MOVE TOT-LWT TO DB2LOCKWAIT COMPUTE DSNDQXST-ADDRVAL = SMF1Ø1RPS-ADDRVAL + QWAØ1R20 - 4 SET ADDRESS OF DSNDQXST TO DSNDQXST-PNTR MOVE Ø TO WS-NUM-SELECTS WS-NUM-INSERTS WS-NUM-UPDATES WS-NUM-DELETES WS-NUM-OPENS WS-NUM-FETCHES WS-NUM-CLOSES WS-NUM-PREPARES PERFORM 3500-SOL OWA01R2N TIMES. MOVE WS-NUM-SELECTS TO NUM-SELECTS TO NUM-INSERTS MOVE WS-NUM-INSERTS MOVE WS-NUM-UPDATES TO NUM-UPDATES MOVE WS-NUM-DELETES TO NUM-DELETES MOVE WS-NUM-OPENS TO NUM-OPENS MOVE WS-NUM-FETCHES TO NUM-FETCHES MOVE WS-NUM-CLOSES TO NUM-CLOSES MOVE WS-NUM-PREPARES TO NUM-PREPARES MOVE Ø TO TOT-IO. COMPUTE DSNDQBAC-ADDRVAL = SMF1Ø1RPS-ADDRVAL + QWAØ1R30 - 4 SET ADDRESS OF DSNDQBAC TO DSNDQBAC-PNTR PERFORM 4000-COMPUTE-IO QWA01R3N TIMES MOVE TOT-IO TO DB2IO MOVE ALL SPACES TO DDF-REC MOVE '----' TO REMOTE-LOCN COMPUTE DSNDQLAC-ADDRVAL = SMF1Ø1RPS-ADDRVAL + QWAØ1R50 - 4 SET ADDRESS OF DSNDQLAC TO DSNDQLAC-PNTR PERFORM 5000-DDF QWAØ1R5N TIMES. 2000-EXIT. EXIT. 2100-CONVERT-TIME. COMPUTE TEMP-HOURS = SMF1Ø1TME / 36ØØØØ

```
COMPUTE TEMP-MIN
     (SMF1Ø1TME - (TEMP-HOURS * 360000)) / 6000
     COMPUTE TEMP-SECS = (SMF1Ø1TME - (TEMP-HOURS * 36ØØØØ) -
                               (TEMP-MIN * 6000)) / 100.
3000-COMPUTE-TIMES.
    COMPUTE TOT-INDB2-ELAP = TOT-INDB2-ELAP + QWACASC
    COMPUTE TOT-DB2-CPU = TOT-DB2-CPU + OWACAJST + OWACASRB
    COMPUTE TOT-IOWT = TOT-IOWT + QWACAWTI
    COMPUTE TOT-LWT = TOT-LWT + QWACAWTL
    COMPUTE DSNDOWAC-ADDRVAL = DSNDOWAC-ADDRVAL + OWA01R1L
    SET ADDRESS OF DSNDQWAC TO DSNDQWAC-PNTR.
3000-EXIT.
   EXIT.
3500-SOL.
    ADD QXSELECT TO WS-NUM-SELECTS
     ADD QXINSRT
                     TO WS-NUM-INSERTS
     ADD OXUPDTE
                    TO WS-NUM-UPDATES
     ADD QXDELET
                     TO WS-NUM-DELETES
                     TO WS-NUM-OPENS
     ADD QXOPEN
     ADD QXFETCH
                    TO WS-NUM-FETCHES
                     TO WS-NUM-CLOSES
     ADD OXCLOSE
     ADD OXPREP
                     TO WS-NUM-PREPARES
     COMPUTE DSNDQXST-ADDRVAL = DSNDQXST-ADDRVAL + OWAØ1R3L
     SET ADDRESS OF DSNDQXST TO DSNDQXST-PNTR.
3500-EXIT.
    EXIT.
4000-COMPUTE-IO.
     EVALUATE QBACPID
     WHEN Ø
       MOVE 'BPØ'
                            TO BPID
     WHEN 1
       MOVE 'BP1'
                            TO BPID
     WHEN 2
       MOVE 'BP2'
                           TO BPID
     WHEN OTHER
       MOVE 'BPXX'
                            TO BPID
     END-EVALUATE
     MOVE OBACGET
                            TO NUM-GETP
     MOVE QBACRIO
                            TO NUM-SYNC-READ
     MOVE OBACIMW
                            TO NUM-SYNC-WRITE
     COMPUTE TOT-IO = TOT-IO + QBACGET + QBACSWS
     COMPUTE DSNDQBAC-ADDRVAL = DSNDQBAC-ADDRVAL + QWAØ1R3L
     SET ADDRESS OF DSNDQBAC TO DSNDQBAC-PNTR
     IF TOT-DB2-CPU >= CUTOFF-CPU-TIME
        WRITE BPACCT-REC
                               FROM WS-BPACCT-REC
     END-IF.
4000-EXIT.
    EXIT.
```

```
5000-DDF.
    MOVE QLACLOCN TO REMOTE-LOCN
    MOVE OLACSOLS TO SOLS-SENT
    MOVE QLACSQLR TO SQLS-RECEIVED
    MOVE QLACROWS TO ROWS-SENT
    MOVE QLACBYTS TO BYTES-SENT
    MOVE OLACBYTR TO BYTES-RECEIVED.
    COMPUTE DSNDQLAC-ADDRVAL = DSNDQLAC-ADDRVAL + QWAØ1R5L
    SET ADDRESS OF DSNDOLAC TO DSNDOLAC-PNTR.
5000-EXIT.
    EXIT.
6000-DISPLAY-ACCOUNTING-INFO.
    if ws-smf-recno = 1
      DISPLAY OUT-REC-HEAD1
      display ws-hypens
      DISPLAY OUT-REC-HEADER
      display ws-hypens
    end-if.
**** RECORD IS WRITTEN ONLY IF IN DB2 TIME IS MORE THAN 10 MS
IF TOT-DB2-CPU >= CUTOFF-CPU-TIME
       DISPLAY OUT-REC-DATA
    END-IF.
```

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DB2 OLAP Server for OS/390 Version 7.1

IBM has made significant improvements to DB2 OLAP Server, as is shown by the increase in the Version number from 1.1 to 7.1! It is available for other platforms (see any IBM announcement for available hardware) but this article only discusses the mainframe version (OS/ 390 – old-timers will remember that mainframes were declared dead last decade).

OLAP is the acronym for On Line Analytical Processing. Ted Codd coined it in his 1993 paper *Providing OLAP to User-Analysis: An IT Mandate* although its foundation was established in the 1960s with IBM's APL (A Programming Language). OLAP lets users model their business as a multidimensional cube whose typical dimensions are

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geography, product, and time. OLAP comes in the following flavours:

- DOLAP Desktop OLAP that usually employs two-tier architecture with data downloaded from a server for analysis on the client.
- ROLAP Relational OLAP uses a relational database like DB2. Typical architecture is two-tier with a client generating a query and displaying the results.
- MOLAP– Multidimensional OLAP. Typical architecture does analysis on a mid-tier server like OS/390 using the Web for reporting. IBM defines DB2 OLAP Server as a MOLAP.
- HOLAP Hybrid OLAP that attempts to combine the advantages of ROLAP and MOLAP. My view, contrary to IBM, is that DB2 OLAP Server is a HOLAP.

WHY DB2 OLAP SERVER IS A HOLAP

DB2 OLAP Server uses Essbase (Hyperion Solutions Corporation) as its multidimensional engine, but can store the data in DB2 allowing it to perform relational management functions such as back-up and recovery. Any OLAP application can also be assigned to DB2 using a star schema. (*A relational schema for a data warehouse*, *DB2 Update*, Issue 86, December 1999 is my description of star schemas and variants.)

STAR SCHEMA

Star schema is divided into two table types of 'fact' and 'dimension'. A fact table is at the star's 'centre' whereas dimension tables are at its 'points'. Fact tables contain columns to be measured such as sales and units. Dimension tables hold measuring columns such as date and location. A simple star table looks like:

date_dimkey location_dimkey \$\$\$ units

dimkey is a 'foreign key' to a dimension table such as date. DB2 OLAP Server usually creates four fact tables to hold all the data values for a relational cube.

Multidimensional cube software allows users to ask reporting questions (what happened when and where), do planning (what if we did this), and forecast (what next). Software also provides drill-down to details, drill-up to summaries or global views, pivot or slice-and-dice for different perspectives, and crosstabs for displaying data summaries based on their characteristics. DB2 enhanced SQL with new join and GROUP BY features can answer many multidimensional queries without needing multidimensional software. (See my articles *All those joins, DB2 Update,* Issue 97, November 2000, *Have you been cubed?, DB2 Update,* Issue 98, December 2000, and *Have you been cubed? – further thoughts, DB2 Update,* Issue 100, February 2001 for details and examples).

DB2 OLAP Server component Relational Storage Manager (RSM) creates and manages star schema objects (tables, views, indexes, and calculated data) from a metaoutline that you can create using the OLAP Metaoutline tool (a component of DB2 OLAP Integration Server Desktop). The DB2 OLAP Server calculation engine and 100+ functions simplify cube population and query execution, helping to optimize performance.

ANCHOR DIMENSION

DB2 OLAP Server uses an anchor dimension to define the relational cube fact table structure. It should be 'dense' having few nulls ('sparse' is many nulls). Minimizing nulls improves storage efficiency (the DB2 OLAP Server default is the densest dimension). An anchor dimension cannot be deleted or changed after data is loaded without clearing all that data, so be careful choosing it!

STORAGE MANAGERS

DB2 OLAP Server allows you to specify either a multidimensional storage manager or a relational storage manager. Multidimensional is the default but each OLAP application can specify either. Both multidimensional and relational store calculation and report scripts, data load rules, and database outline (defines all database elements) in a file system. Multidimensional stores data in the file system

whereas relational stores it in a DB2 relational database using a star schema.

DESIGNING AN OLAP APPLICATION

The key step is defining your database outline, which contains:

- Dimension and member definitions.
- Dense/sparse dimension tags and attributes.
- Anchor dimension attributes.
- Calculations.
- Shared members.
- Roll-up rules.

Implementation strategies, techniques, and instructions for the above are provided in the *OLAP Database Administrator's Guide, Volumes I & II* (SC27-0788 and SC27-0789).

DB2 OLAP STARTER KIT

The DB2 OLAP Starter Kit is a subset of the DB2 OLAP Server provided free with DB2 UDB V7. It allows you to create a limited OLAP application, which IBM hopes will be so successful that you will acquire the full functionality of DB2 OLAP Server. OLAP applications developed with the Starter Kit can be uploaded to DB2 OLAP Server.

DB2 UDB DATA WAREHOUSE CENTER

DB2 OLAP Server Version 7.1 has been designed to integrate with UDB Data Warehouse Center to manage your OLAP environment. The combined functionality includes:

- Registering and accessing data sources.
- Defining data extraction and transformation steps.
- Populating data warehouse and OLAP applications.

- Managing, interchanging, and reporting metadata.
- Mapping multiple source data into star schema or multidimensional cubes.
- Loading, pre-calculating, and reporting on cubes.

DB2 OLAP Server can use any data warehouse star schema as a data source but cannot use it as its relational cube.

WHY

There is much hype about OLAP. A 1999 study by IDC (*IDC Information Access Tools: 1999 Worldwide Markets and Trends*) forecast a compound growth rate of 43.7% to 2003. This is evidence that many organizations are using or planning to use OLAP seriously. The applications are endless including:

- E-commerce.
- Customer Relationship Marketing (CRM).
- Enterprise Resource Planning (ERP).
- Competitor Analysis.
- Empowering Employees.
- Quality Engineering.

Information Catalog Manager can extract and publish meaningful metadata for users.

I recommend accepting IBM's marketing ploy to use the Starter Kit to develop a test OLAP application. You have little to lose and much to gain!

CONCLUSION

Future articles will discuss various strategies for getting the most out of DB2 OLAP Server and DB2 UDB Data Warehouse Center.

Eric Garrigue Vesely Principal/Analyst Workbench Consulting (Malaysia)

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Sample user-defined functions

This article describes user-defined functions that are provided with DB2. A user-defined function is an extension to the SQL language. A user-defined function is very similar to a stored procedure or a host language subprogram or function. However, a user-defined function is often the better choice for an SQL application because you can invoke a user-defined function in an SQL statement. You can use a user-defined function in an SQL statement in the same way as built-in functions.

The user-defined function runs in the stored procedure's address space that is established in the Workload Manager (WLM) environment. More information about Workload Manager is in red book *DB2 UDB* for OS/390 Version 6 Management Tools Package, SG24-5759-00.

To prepare a user-defined function for execution, you must perform these steps:

- Create the user-defined function with the SQL command CREATE FUNCTION....
- Precompile the user-defined function program and bind the DBRM into a package. You need to do this only if your user-defined function contains SQL statements. You do not need to bind a plan for the user-defined function.
- Compile the user-defined function program with a compiler that supports Language Environment and link-edit the appropriate Language Environment components with the user-defined function. You must also link-edit the user-defined function with RRSAF. The language interface module for RRSAF, DSNRLI, is shipped with the linkage attributes AMODE(31) and RMODE(ANY).
- GRANT EXECUTE authority on the user-defined function or package.

The following is a list of my user-defined functions. I wrote all these function in PL/I:

- AGE returns age in a user-specified format.
- COUNTER a function that increments a variable in the scratchpad each time it is invoked.
- CUMUL a function that generates a cumulative column.
- DATEUDF returns the current date or a user-specified date in a user-specified format.
- **REVERSE** returns the input string in reverse format.
- TIMEUDF returns the current time or a user-specified time in a user-specified format.

AGE

The schema is SYSADM.

The AGE function returns age in two formats. The first one is in 'YYMMDD' format, and the second one is in 'YY YEAR(S) MM MONTH(S) DD DAY(S)' format. The formula used to calculate AGE is:

AGE = CURRENT DATE - DATE(expression)

The expression must be a date.

For example: the current date is '2000-10-26'.

```
SELECT AGE(DATE('1957-Ø8-18')) AGE1,
AGE(DATE('1957-Ø8-18'),'T') AGE2
FROM SYSIBM.SYSDUMMY1
```

Result:

AGE1	AGE2
430208	43 YEARS 2 MONTHS 8 DAYS

COUNTER

The schema is SYSADM.

The COUNTER function increments an integer variable in the scratchpad each time it is invoked. If an integer value is not specified, then the COUNTER function increments by 1.

For example:

```
SELECT COUNTER() "INCREMENT BY 1",
COUNTER(5) "INCREMENT BY 5", NAME
FROM SYSIBM.SYSTABLES
WHERE DBNAME='DSNDBØ6'
AND TSNAME='SYSOBJ'
```

Result:

INCREMENT BY 1	INCREMENT BY 5	NAME
1	5	SYSCONSTDEP
2	10	SYSAUXRELS
3	15	SYSDATATYPES
4	2Ø	SYSROUTINES
5	25	SYSPARMS
6	3Ø	SYSROUTINEAUTH
7	35	SYSTRIGGERS
8	4Ø	SYSSCHEMAAUTH

CUMUL

The schema is SYSADM.

The CUMUL function returns a cumulative value for each row. The argument must be smallint, integer, decimal, or float.

For example:

SELECT ID, CUMUL(ID) "CUMUL ID", SALARY, CUMUL(SALARY) "CUMUL SALARY" FROM Q.STAFF WHERE JOB='MGR'

Result:

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ΙD	CUMUL ID	SALARY	CUMUL SALARY
1Ø	1Ø	18357.50	18357.500
3Ø	4Ø	17506.75	35864.250

5Ø	9Ø	20659.80	56524.050
100	19Ø	18352.80	74876.850
14Ø	33Ø	21150.00	96026.850
16Ø	49Ø	22959.20	118986.050
21Ø	7ØØ	20010.00	138996.050
24Ø	94Ø	19260.25	158256.300
26Ø	1200	21234.00	179490.300
27Ø	147Ø	18555.50	198045.800
29Ø	176Ø	19818.00	217863.800

DATEUDF

????, ' \'??????
????,'M'??????
????,'N'??????
????, ' U'??????
????,'E'?????
????,'S'?????
????,'L'?????
????,'J'?????
????,'D'?????

The schema is SYSADM.

The DATEUDF function returns the date in one of the following formats:

- Weekday returns the English name for the day of the week, in mixed case.
- Month returns the full English name of the month, in mixed case.
- Normal returns the date in the default format 'YYYY-MM-DD'.
- Usa returns the date in the format 'MM-DD-YYYY'.
- European returns the date in format 'DD.MM.YYYY'.
- Short returns the date in format 'DD Mon YYYY'.
- Long returns the date in format 'DD Month YYYY'.
- Julian returns the date in format 'YYYDDD'.
- **D**ays returns the number of days, including the current day, in the format 'DDD' (no leading zeros).

The argument must be a date.

For example: the current date is '2000-10-27'.

```
SELECT DATEUDF(DATE('1957-Ø8-18'),'W') AS DAY,
DATEUDF(DATE('1957-Ø8-18'),'M') AS MONTH,
DATEUDF(DATE('1957-Ø8-18'),'S') AS "SHORT DAY",
DATEUDF(DATE('1957-Ø8-18'),'L') AS "LONG DAY",
DATEUDF(DATE('1957-Ø8-18'),'L') AS JULIAN,
DATEUDF(DATE('1957-Ø8-18'),'E') AS EUROPEAN,
DATEUDF(DATE('1957-Ø8-18'),'E') AS USA,
DATEUDF(DATE('1957-Ø8-18'),'U') AS USA,
DATEUDF(DATE('1957-Ø8-18'),'D') AS DAYS,
DATEUDF(DATE('1957-Ø8-18'),'N') AS DEFAULT
FROM SYSIBM.SYSDUMMY1;
```

Result:

DAY	MONTH	SHORT DAY	LONG DAY	JULIAN
Sunday	August	18 Aug 1957	18 August 1957	1957230
and:				
EUROPEAN	USA	DAY	S DEFAULT	
18.08.19	 57 Ø8/1	8/1957 230	1957-Ø8-18	

REVERSE

The schema is SYSADM.

The REVERSE function returns an input string in reverse format. The argument must be a character string.

For example:

```
SELECT REVERSE('UDF SAMPLE') AS REVERSE FROM SYSIBM.SYSDUMMY1
```

Result:

```
REVERSE
-----
ELPMAS FDU
```

TIMEUDF

The schema is SYSADM.

The TIMEUDF function returns the time in one of the following formats:

- Normal returns the time in the default format 'hh:mm:ss'.
- Minutes returns the number of minutes since midnight in the format 'mmm' (no leading zeros).
- Seconds returns the number of seconds since midnight in the format 'sssss' (no leading zeros).
- **D** returns the time in the format 'hhmmss'.
- Civil returns the time in the Civil format 'hh:mm am/pm'.

For example:

```
SELECT TIMEUDF(TIME(CURRENT TIME),'N') AS DEFAULT,
TIMEUDF(TIME(CURRENT TIME),'M') AS MINUTES,
TIMEUDF(TIME(CURRENT TIME),'S') AS SECONDS,
TIMEUDF(TIME(CURRENT TIME),'D') AS LOCAL,
TIMEUDF(TIME(CURRENT TIME),'C') AS CIVIL
FROM SYSIBM.SYSDUMMY1 WITH UR
```

Result:

DEFAULT	MINUTES	SECONDS	LOCAL	CIVIL
14:45:22	885	53122	144522	2:45pm

THE UDF (USER-DEFINED FUNCTION) OBJECTS

The first step is to define my sample UDFs to DB2.

AGE

CREATE FUNCTION SYSADM.AGE (DATE) RETURNS VARCHAR(10) SPECIFIC AGE EXTERNAL NAME 'AGE' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL DETERMINISTIC READS SQL DATA DBINFO FENCED COLLID AGE WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT NO SCRATCHPAD NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2: CREATE FUNCTION SYSADM.AGE (DATE , VARCHAR(1)) RETURNS VARCHAR(3Ø) SPECIFIC AGED EXTERNAL NAME 'AGED' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL DETERMINISTIC READS SQL DATA DBINFO FENCED COLLID AGED WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT NO SCRATCHPAD NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2: COUNTER CREATE FUNCTION SYSADM.COUNTER RETURNS INTEGER SPECIFIC COUNTER EXTERNAL NAME 'COUNTER' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL NOT DETERMINISTIC NO SQL NO DBINFO FENCED NO COLLID WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT SCRATCHPAD 100 NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2: CREATE FUNCTION SYSADM.COUNTER (INTEGER) RETURNS INTEGER SPECIFIC COUNTERV EXTERNAL NAME 'COUNTERS' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SOL NOT DETERMINISTIC NO SQL NO DBINFO FENCED NO COLLID WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT SCRATCHPAD 100 NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2; CUMUL CREATE FUNCTION SYSADM.CUMUL (INTEGER) RETURNS INTEGER SPECIFIC CUMULI EXTERNAL NAME 'CUMULI' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL

NOT DETERMINISTIC NO SQL NO DBINFO FENCED NO COLLID WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT SCRATCHPAD 100 NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2; CREATE FUNCTION SYSADM.CUMUL (DECIMAL(15.3)) RETURNS DECIMAL(15,3) SPECIFIC CUMULD EXTERNAL NAME 'CUMULD' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL NOT DETERMINISTIC NO SQL NO DBINFO FENCED NO COLLID WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT SCRATCHPAD 100 NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2: CREATE FUNCTION SYSADM.CUMUL (DOUBLE) RETURNS DOUBLE SPECIFIC CUMULF EXTERNAL NAME 'CUMULF' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL NOT DETERMINISTIC NO SQL NO DBINFO FENCED

NO COLLID WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT SCRATCHPAD 100 NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2: DATEUDF CREATE FUNCTION SYSADM.DATEUDF (DATE , VARCHAR(1)) RETURNS VARCHAR(2Ø) SPECIFIC DATEUDF EXTERNAL NAME 'DATEUDF' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SOL DETERMINISTIC READS SQL DATA DBINFO FENCED COLLID DATEUDF WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT NO SCRATCHPAD NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2; CREATE FUNCTION SYSADM.DATEUDF (DATE) RETURNS VARCHAR(1Ø) SPECIFIC DATEUDFD EXTERNAL NAME 'DATEUDFD' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL DETERMINISTIC READS SOL DATA DBINFO FENCED COLLID DATEUDFD

WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT NO SCRATCHPAD NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2; REVERSE CREATE FUNCTION SYSADM.REVERSE (VARCHAR(4Ø46)) RETURNS VARCHAR(4Ø46) SPECIFIC REVERSE EXTERNAL NAME 'REVERSE' -- MVS load module LANGUAGE PLI PARAMETER STYLE DB2SQL DETERMINISTIC NO SOL DBINFO FENCED NO COLLID WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN EXTERNAL ACTION RETURNS NULL ON NULL INPUT SCRATCHPAD 100 NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2; TIMEUDF CREATE FUNCTION SYSADM.TIMEUDF (TIME , VARCHAR(1)) RETURNS VARCHAR(1Ø) SPECIFIC TIMEUDF -- MVS load module EXTERNAL NAME 'TIMEUDF' LANGUAGE PLI PARAMETER STYLE DB2SOL DETERMINISTIC READS SQL DATA DBINFO

FENCED COLLID TIMEUDF WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT NO SCRATCHPAD NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2: CREATE FUNCTION SYSADM.TIMEUDF (TIME) RETURNS VARCHAR(1Ø) SPECIFIC TIMEUDFD EXTERNAL NAME 'TIMEUDFD' -- MVS load module LANGUAGE PLT PARAMETER STYLE DB2SQL DETERMINISTIC READS SOL DATA DBINFO FENCED COLLID TIMEUDFD WLM ENVIRONMENT DSNNWLM1 -- WLM application STAY RESIDENT YES PROGRAM TYPE MAIN NO EXTERNAL ACTION RETURNS NULL ON NULL INPUT NO SCRATCHPAD NO FINAL CALL DISALLOW PARALLEL ASUTIME NO LIMIT SECURITY DB2;

LOAD MODULES

AGE – PL/I source code

/* */ UDF : AGE /* INPUT : UDF_PARM1 CHAR(1Ø) */ /* OUTPUT: UDF RESULT VARCHAR(1Ø) */ CHAR(1Ø); /* INPUT PARAMETER CHAR(1Ø) VAR; /* RESULT PARAMETER DCL UDF_PARM1 */ DCL UDF RESULT */ BIN FIXED(15); /* INDICATOR FOR INPUT PARM */ BIN FIXED(15); /* INDICATOR FOR RESULT */ DCL UDF_IND1 DCL UDF INDR */ /* SCRATCHPAD DCL 1 UDF SCRATCHPAD. 3 UDF_SPAD_LEN BIN FIXED(31), 3 UDF_SPAD_TEXT CHAR(100); EXEC SOL INCLUDE UDFINFO: /* DBINFO */ DCL (ADDR, LENGTH, SUBSTR, NULL) BUILTIN: EXEC SQL INCLUDE SQLCA; /* RETURNS AGE IN FORMAT 'YYMMDD' -> CURRENT DATE - DATE */ /* AGE('1957-Ø8-18') -> '430202' CURRENT DATE ='2000-10-20' */ EXEC SQL SELECT CHAR(INTEGER(CURRENT DATE - DATE(:UDF PARM1))) INTO :UDF RESULT FROM SYSIBM.SYSDUMMY1 WITH UR;

END AGE;

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AGED – PL/I source code

```
* PROCESS SYSTEM(MVS);
AGET: PROC(UDF_PARM1, UDF_PARM2, UDF_RESULT,
             UDF_IND1, UDF_INDR,
             UDF_SQLSTATE, UDF_NAME, UDF_SPEC_NAME,
             UDF DIAG MSG. UDF SCRATCHPAD.
             UDF_CALL_TYPE, UDF_DBINFO)
        OPTIONS(MAIN NOEXECOPS REENTRANT):
 UDF : AGET
/*
                                                              */
/*
      INPUT : UDF PARM1 CHAR(1Ø)
                                                              */
/*
           : UDF_PARM2
                                                              */
                         VARCHAR(1)
      OUTPUT: UDF_RESULT VARCHAR(3Ø)
 /*
                                                              */
CHAR(1Ø); /* INPUT PARAMETER
CHAR(1) VAR; /* INPUT PARAMETER
CHAR(3Ø) VAR; /* RESULT PARAMETER
 DCL UDF_PARM1
                                                              */
 DCL UDF_PARM2
                                                              */
 DCL UDF_RESULT
                                                              */
 DCLUDF_IND1BINFIXED(15);/*INDICATORFORINPUTPARMDCLUDF_INDRBINFIXED(15);/*INDICATORFORRESULT
                                  /* INDICATOR FOR INPUT PARM */
                                                              */
 DCL 1 UDF_SCRATCHPAD,
                                   /* SCRATCHPAD
                                                              */
```

3 UDF_SPAD_LEN BIN FIXED(31), 3 UDF_SPAD_TEXT CHAR(100); EXEC SOL INCLUDE UDFINFO: /* DBINFO */ DCL (ADDR, LENGTH, SUBSTR, NULL) BUILTIN: EXEC SQL INCLUDE SQLCA: /* RETURNS AGE IN FORMAT 'YY YEAR(S) MM MONTH(S) DD DAY(S)' */ /* AGE((DATE('2000-08-18'),'T') -> 43 YEARS 2 MONTHS 2 DAYS */ EXEC SOL SELECT CASE WHEN(SMALLINT(YY)) > 1 THEN YY||' YEARS ' ELSE YY || ' YEAR ' END CONCAT CASE WHEN(SMALLINT(MM)) > 1THEN STRIP(CHAR(SMALLINT(MM)))||' MONTHS ' ELSE STRIP(CHAR(SMALLINT(MM)))||' MONTH ' END CONCAT CASE WHEN(SMALLINT(DD)) > 1THEN STRIP(CHAR(SMALLINT(DD)))||' DAYS' ELSE STRIP(CHAR(SMALLINT(DD)))||' DAY ' END INTO :UDF_RESULT FROM(SELECT SUBSTR(AGE,1,LENGTH(AGE)-4) YY, SUBSTR(AGE, 3, 2) MM, SUBSTR(AGE,5,2) DD FROM (SELECT STRIP(CHAR(INTEGER(CURRENT DATE-DATE(:UDF_PARM1)))) AGE FROM SYSIBM.SYSDUMMY1) X) Y;

END AGET;

Editor's note: the remaining UDFs will be published next month.

Bernard Zver (Slovenia)	© Xephon 2001

Developer Solutions has announced new versions of HiT Software's HiT OLEDB and ODBC DB2 middleware server products. Enhancements include monitoring of remote Windows server middleware activity including connection pooling.

The HiT Server Manager allows a remote developer or administrator to see OLE DB and ODBC DB2 SQL connections by DB2 server, calling application and User ID, start time of connections, status of connection, traffic volume over connection, and memory utilization. DB2 server connections can be pooled and reused to optimise performance. The HiT Server Manager can define the time-out and dropping of these pooled connections.

HiT OLEDB Server and HiT ODBC Server versions are available for DB2 UDB on its supported platforms; there is also an AS/400-specific version.

For further information contact:

Developer Solutions, Bradninch Court, Castle Street, Exeter EX4 3PL, UK. Tel: (01392) 262626.

URL: http://www.developer-solutions.co.uk.

* * *

Compuware has launched Version 7.0 of its XPEDITER/TSO automated testing tool for DB2 Stored Procedures. The debugging option intercepts S/390 DB2 stored procedures that originate from any type of application.

Testing is done in the actual DB2/Workload Manager address space, with no need for simulation of how the stored procedure will run. Developers can trap and step through Stored Procedures written in COBOL, Assembler, PL/I, and C.

Also, they can suspend and resume program execution at any point in the program, display and change program data variables, alter program logic, and set statement execution counts. The product works with XPEDITER/Code Coverage and File-AID DB2.

For further information contact: Compuware, 31440 Northwestern Highway, Farmington Hills, MI 48334-2564, USA. Tel: (248) 737 7300. URL: http://www.compuware.com/

products/xpediter/tso/db2.htm.

SteelEye has announced that LifeKeeper for Windows NT and Linux now works with DB2 Universal Database to provide reliability features for eServer xSeries Intelbased servers. The company worked with IBM to integrate the latest version of LifeKeeper Next Generation Enterprise Reliability platform for DB2 UDB running in Windows NT and Linux.

LifeKeeper monitors system, storage, and application health, maintains client connectivity, and provides continuous data access, regardless of where the clients reside. It's designed to prevent system or application failures.

For further information contact: SteelEye Technology, 2660 Marine Way, Suite 200, Mountain View, CA 94043, USA. Tel: (877) 319 0108. URL: http://www.steeleye.com.



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