



166

CICS

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In this issue

- 3 A shutdown assist and task purge utility
 - 10 Interpreting temporary storage behaviour
 - 19 CEMT Logger – an alternative design – part 2
 - 28 Displaying CPU usage by TCB – revisited
 - 29 Accessing CICS control blocks in COBOL
 - 42 Listing the TMONCICS control file RLC definitions
 - 48 CICS news
-

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Robert Burgess

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A shutdown assist and task purge utility

INTRODUCTION

This article presents a program to purge CICS tasks. It functions either as a shutdown assist program or as a utility to cancel tasks by transaction-id and/or user-id. To use as a shutdown assist, you simply add the program to the PLTSD (for CICS Releases 4.1 and below) or define an associated transaction (eg KILL) and specify it as SDTRAN (for Transaction Server releases). To use in utility mode, you must invoke the program as transaction KILL from a console or terminal and supply the appropriate terminal input.

Here are some console examples:

- ‘F PRODCICS,KILL TRAN=UPDT’ cancels tasks with the transaction-id UPDT.
- ‘F PRODCICS,KILL USER=MRHAPPY’ cancels tasks with the user-id MRHAPPY.
- ‘F PRODCICS,KILL TRAN=UPDT,USER=MRHAPPY’ cancels tasks meeting both criteria.

Operators and systems programmers favour this utility over CEMT to cancel tasks because they are freed from the burden of remembering task numbers and because they can cancel more than one task at a time. The ability to cancel multiple tasks having the same transaction-id is especially useful during deadlock situations.

KILLTASK

```
*ASM XOPTS(CICS,FE,SP)
*
* PROGRAM: KILLTASK
*
* PURPOSE: Purge CICS tasks
*
*           PRINT ON,NOGEN
*
*           DFHREGS
```

```

*
* Register Usage:
*
* R2 -> task count
* R3 code base
* R4 -> task list
* R5 -> tran list corresponding
* R6 -> EIB
* R8 -> TIOA start 3
* R9 -> TIOA offset 3 optional
* RB -> TIOA end 3
* RD -> dynamic storage
*

```

```

DFHEISTG DSECT
CVDA      DS      F
RESP      DS      F
TARGET    DS      PL4
FACILITY  DS      CL4
FACTYPE   DS      F
TASKCNT   DS      F
LOGTEXT   DS      ØH
LOGPART1  DS      CL32
LOGTASK   DS      CLØ7
LOGPART2  DS      CLØ7
LOGTRAN   DS      CLØ4
LOGPART3  DS      CLØ7
LOGTERM   DS      CLØ4
LOGPART4  DS      CLØ7
LOGUSER   DS      CLØ8
LOGPART5  DS      CL14
LOGTYPE   DS      CLØ5
LOGTEXTL  EQU    *-LOGTEXT
LOGL      DS      H
TIOALEN   DS      H
TRANPARAM DS      CL4
USERPARAM DS      CL8
HOWSTART  DS      CL2
KILLFLAG  DS      C
PARAMFLAG DS      C

```

```

*
KILLTASK DFHEIENT CODEREG=(3),DATAREG=(13),EIBREG=(6)
KILLTASK AMODE 31
KILLTASK RMODE ANY

```

```

*
EXEC CICS HANDLE CONDITION ERROR(DUMP)
*
BAL R7,LIST           get list of tasks
L   R2,TASKCNT
LTR R2,R2
BZ  RETURN

```

```

XC      PARMFLAG,PARMFLAG          initialize to no input parms
CLC     EIBTRNID,=C'KILL'
BNE     INITIAL
*
EXEC    CICS ASSIGN STARTCODE(HOWSTART)
*
CLC     HOWSTART,=C'TD'             possible terminal input?
BNE     INITIAL                     n - skip receive
*
EXEC    CICS RECEIVE SET(R8)        +
      LENGTH(TIOALEN)
*
LH      R11,TIOALEN
C       R11,=F'10'                  minimum meaningful tioalen?
BL      INITIAL                     n - skip parse
LA      R10,1                       set parsing increment
AR      R11,R8                      point to last byte tioa
S       R11,=F'6'                   ... less 6
LA      R8,4(R8)                   begin parse after transid
LA      R9,4                         initialize tioa offset
GETPARAM DS      0H
TM      PARMFLAG,X'03'              bothparms loaded?
BC      3,INITIAL                   y - we're done parsing
TM      PARMFLAG,X'01'              tran parm already loaded?
BC      3,GETPARAM1                 y - skip parse for TRAN=
CLC     0(5,R8),=C'TRAN='
BNE     GETPARAM1
LA      R8,5(R8)                    skip over TRAN=
LA      R9,5(R9)                    bump tioa offset accordingly
LH      R12,TIOALEN
SR      R12,R9
BCTR   R12,R0
C       R12,=F'3'
BNH    LOADTRAN
LA      R12,3
LOADTRAN LA      R7,TRANPARAM
EX      R12,TRANMVC
TR      TRANPARAM,XLTAB             blank any valid delimiter
LA      R12,3
PADTRAN CLI     0(R7),X'40'
BNE     PADTRAN1
MVI    1(R7),X'40'
PADTRAN1 LA     R7,1(R7)
BCT    R12,PADTRAN
OI     PARMFLAG,X'01'              signal tran parm input
B      GETPARAM1
*
TRANMVC MVC     0(1,R7),0(R8)
*
GETPARAM1 TM    PARMFLAG,X'02'      user parm already loaded?

```

	BC	3,GETPARM2	y - skip parse for USER=
	CLC	Ø(5,R8),=C'USER='	
	BNE	GETPARM2	
	LA	R8,5(R8)	skip over USER=
	LA	R9,5(R9)	bump tioa offset accordingly
	LH	R12,TIOALEN	
	SR	R12,R9	
	BCTR	R12,RØ	
	C	R12,=F'7'	
	BNH	LOADUSER	
	LA	R12,7	
LOADUSER	LA	R7,USERPARM	
	EX	R12,USERMVC	
	TR	USERPARM,XLTAB	blank any valid delimiter
	LA	R12,7	
PADUSER	CLI	Ø(R7),X'4Ø'	
	BNE	PADUSER1	
	MVI	1(R7),X'4Ø'	
PADUSER1	LA	R7,1(R7)	
	BCT	R12,PADUSER	
	OI	PARMFLAG,X'Ø2'	signal user parm input
GETPARM2	LA	R9,1(R9)	increment tioa offset
	BXLE	R8,R1Ø,GETPARM	
	B	INITIAL	end parse
	*		
USERMVC	MVC	Ø(1,R7),Ø(R8)	
	*		
INITIAL	DS	ØH	
	XC	KILLFLAG,KILLFLAG	initialize kill flag
	MVC	LOGPART1,MSGPART1	set up CSMT log msg
	MVC	LOGPART2,MSGPART2	
	MVC	LOGPART3,MSGPART3	
	MVC	LOGPART4,MSGPART4	
	MVC	LOGPART5,MSGPART5	
	LA	R7,LOGTEXTL	get length for TD write
	STH	R7,LOGL	
	MVC	CVDA,DFHVALUE(PURGE)	set up task purge
	MVC	LOGTYPE,PURGE	
	BAL	R7,KILL	do first level cancels
	CLI	KILLFLAG,X'FF'	were there any?
	BNE	RETURN	n - then we're done
	EXEC	CICS DELAY FOR SECONDS(3)	
	BAL	R7,LIST	see what tasks remain
	L	R2,TASKCNT	
	LTR	R2,R2	
	BZ	RETURN	
	MVC	CVDA,DFHVALUE(FORCEPURGE)	we're done kidding around
	MVC	LOGTYPE,FORCE	
	BAL	R7,KILL	do second level cancels
	B	RETURN	

```

*
LIST      DS      0H
*
          EXEC    CICS INQUIRE TASK LIST          +
          LISTSIZE(TASKCNT)                        +
          SET(R4)                                   +
          SETTRANSID(R5)
*
          BR      R7
*
KILL      DS      0H
          ZAP     TARGET,0(4,R4)                   get taskno from list
          CP      TARGET,EIBTASKN                  no suicides
          BE      KILLNEXT
*
          EXEC    CICS INQUIRE                      +
          TASK(TARGET)                             +
          FACILITY(FACILITY)                       +
          FACILITYTYPE(FACTYPE)                   +
          USERID(LOGUSER)
*
          CLI     PARMFLAG,X'00'
          BE      KILL00
          CLI     PARMFLAG,X'01'
          BE      KILL01
          CLI     PARMFLAG,X'02'
          BE      KILL02
          CLI     PARMFLAG,X'03'
          BE      KILL03
          B       RETURN                            illogic - shouldn't get here
KILL00    LA     R8,EXCLTAB                          check exclude table
          LA     R10,4
          LR     R11,R8
          LA     R11,EXCLLEN(R11)
          BCTR   R11,R0
KILL00X   CLC   0(4,R8),0(R5)                        task's tran in exclude list?
          BE     KILLNEXT                            y - then bypass cancel
          BXLE  R8,R10,KILL00X
          B     KILLTSK
KILL01   CLC   TRANPARM,0(R5)
          BE     KILLTSK
          B     KILLNEXT
KILL02   CLC   USERPARM,LOGUSER
          BE     KILLTSK
          B     KILLNEXT
KILL03   CLC   TRANPARM,0(R5)
          BNE   KILLNEXT
          CLC   USERPARM,LOGUSER
          BE     KILLTSK
          B     KILLNEXT

```

```

KILLTSK MVI KILLFLAG,X'FF' show that we fired
*
EXEC CICS SET
TASK(TARGET)
PURGETYPE(CVDA)
RESP(RESP)
*
CLC RESP,DFHRESP(NORMAL)
BNE KILLNEXT
MVC LOGTERM,=4X'40' clear termid
CLC FACTYPE,DFHVALUE(TERM) is facility a terminal?
BNE KILLLOG n - don't log facility
MVC LOGTERM,FACILITY
*
KILLLOG UNPK LOGTASK,TARGET format log msg
OI LOGTASK+6,X'F0'
MVC LOGTRAN,0(R5)
*
EXEC CICS WRITEQ TD
QUEUE('CSMT')
FROM(LOGTEXT)
LENGTH(LOGL)
*
KILLNEXT LA R4,4(R4) bump to next task in list
LA R5,4(R5) next tran in list
BCT R2,KILL repeat kill for each task
BR R7
*
DUMP DS 0H
EXEC CICS DUMP TASK DUMPCODE(EIBTRNID)
B RETURN
*
RETURN DS 0H
EXEC CICS RETURN
*
MSGPART1 DC CL32'KILLTASK issued cancel for task '
MSGPART2 DC CL07', tran '
MSGPART3 DC CL07', term '
MSGPART4 DC CL07', user '
MSGPART5 DC CL14', purgetype = '
*
PURGE DC CL05'PURGE'
FORCE DC CL05'FORCE'
*
EXCLTAB DS 0H Tran Exclude Table
DC CL4'AAON' Abend-Aid
DC CL4'DSNC' DB2
DC CL4'OMEG' Omegamon
EXCLLEN EQU *-EXCLTAB
*

```



```

XLTAB   DS      0H
00004905
DC      XL16'400102030405060708090A0B0C0D0E0F'  00 = null
DC      XL16'101112131415161718191A1B1C1D1E1F'
DC      XL16'202122232425262728292A2B2C2D2E2F'
DC      XL16'303132333435363738393A3B3C3D3E3F'
DC      XL16'404142434445464748494A404C4D4E4F'  4B = .
DC      XL16'505152535455565758595A5B5C5D5E5F'  5E = ;
DC      XL16'606162636465666768696A406C6D6E6F'  6B = ,
DC      XL16'7071727374757677787940707C7D7E7F'  7A = :
DC      XL16'8081828384858687888980808C8D8E8F'
DC      XL16'9091929394959697989990909C9D9E9F'
DC      XL16'A0A1A2A3A4A5A6A7A8A9A0A0ACADAEAF'
DC      XL16'B0B1B2B3B4B5B6B7B8B9B0B0BCBDBEBF'
DC      XL16'C0C1C2C3C4C5C6C7C8C9C0C0CCCDCECF'
DC      XL16'D0D1D2D3D4D5D6D7D8D9D0D0DCDDDEDF'
DC      XL16'E0E1E2E3E4E5E6E7E8E9E0E0ECEDEEEF'
DC      XL16'F0F1F2F3F4F5F6F7F8F9F0F0FCFDFEFF'
*
      LTORG
      END    KILLTASK

```

PROGRAM NOTES

There are a few features of KILLTASK that you may want to know about before using it at your site. First-level cancels are issued in the form of a task PURGE; only if tasks remain after a three-second interval are second-level cancels issued in the form of a task FORCE.

Any transactions you want to make ineligible for cancellation should be placed in the transaction exclude table. Our table excludes the DB2 connection, Abend-Aid/FX, and Omegamon because these products have their own entries in the PLTSD, and it is generally a good idea to exclude non-application transactions. Logging to transient data queue CSMT is performed for any tasks purged.

Finally, a good deal of freedom is allowed with respect to utility mode terminal input. Parameters 'USER=' and 'TRAN=' may be entered with either one first. Spacing and use of delimiters is free-form.

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CICS Systems Programmer (USA)

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Interpreting temporary storage behaviour

INTRODUCTION

The temporary storage component of CICS has been rewritten and restructured into a CICS Domain in CICS Transaction Server Release 1. The old DFHTSP code of CICS/ESA Release 4.1.0 and below has been replaced with a new suite of programs, written using object-oriented techniques. This enhancement has brought with it the benefits of structural design, stability, and maintainability that the Domain model provides to CICS.

With the rewrite of temporary storage for CICS Transaction Server, a number of queries against the restructured code have been reported to the CICS change team. This article addresses these observations, and explains their background and how they should be interpreted.

CICS TRANSACTION SERVER

This article makes reference to CICS Transaction Server for OS/390 Releases 1 and 2. CICS Transaction Server is a member of the OS/390 family of MVS-based software servers. IBM has integrated CICS with a set of supporting software, offering a single product in their place.

The CICS component of CICS Transaction Server Release 1 has a Release number of 0510. The CICS component of CICS Transaction Server Release 2 has a Release number of 0520. There is no separately orderable product such as 'CICS/ESA 5.1.0', however – it is the CICS component of CICS Transaction Server Release 1. IBM has recently shipped CICS Transaction Server Release 3. The CICS component of CICS Transaction Server Release 3 has a Release number of 0530.

BACKGROUND TO CICS TEMPORARY STORAGE MANAGEMENT

The basic concepts of temporary storage control under CICS are the same in CICS Transaction Server as they were for CICS/ESA Release 4.1.0 and prior releases. An API is provided for CICS applications to

write to, read from, and delete record data held on temporary storage queues. The queue name is a unique identifier on a given CICS system that represents the collection of records held upon the queue.

Using the EXEC CICS WRITEQ, READQ, and DELETEQ API commands, applications can store, retrieve, and delete temporary storage queues. The data on each queue can be held on one of several media. Main temporary storage resides in-core within the CICS address space. Auxiliary temporary storage is maintained in a number of Control Intervals residing on a CICS-maintained VSAM ESDS, DFHTEMP. Finally, shared temporary storage (introduced in CICS Transaction Server Release 1) allows queue data to be stored within a Coupling Facility.

Auxiliary temporary storage data is read in and written out of the CICS address space at the Control Interval level; that is, CICS will read and write entire Control Intervals to and from DFHTEMP. Within each Control Interval, record data for different auxiliary temporary storage queues may well be held contiguously. CICS maintains index information for retrieval of given records on each queue.

To optimize performance, an EXEC CICS DELETEQ of a temporary storage queue will not cause CICS to read in and rearrange every Control Interval that contains a record on the queue being deleted. Such a physical deletion of data on every DELETEQ would be expensive – a queue could have up to 32,767 records upon it, scattered across many Control Intervals on DFHTEMP.

To read in each Control Interval via I/O and then compress them to remove the redundant record data with Move Character Long (MVCL) instructions would increase path length and response time within temporary storage processing to an unacceptable level. Instead, CICS removes the references to the deleted queue from within its internal control blocks. It also updates the temporary storage byte map – this is a control block with a byte for every Control Interval on DFHTEMP.

The byte map records the amount of Free Space remaining within each Control Interval, and is used when CICS selects a target Control Interval for storing a new record, which will be added to an auxiliary temporary storage queue. By updating the bytes for the Control

Intervals containing records for a deleted queue, to reflect the fact that this space is now available once more, CICS logically deletes the data.

When a Control Interval is later selected to store a new auxiliary temporary storage record for a WRITEQ request, CICS will select a Control Interval that contains enough Free Space to accommodate the new record, by means of the byte map. CICS will then read the Control Interval into a temporary storage buffer in-core within the address space. It then has to compress the buffer to squeeze out all the redundant record space. It does this by determining the records within the Control Interval that are still required (ie still have existing queues) and moving them along into one contiguous block, starting at the beginning of the Control Interval.

This series of MVCL instructions results in the records for deleted queues being overwritten by this compression process, and the Free Space in the Control Interval being repositioned into a contiguous area at the end of the Control Interval. This Free Space can then be used to accommodate the new record being added.

As an aside, CICS uses this compression process as a good point to validate the consistency of the temporary storage control block data and Control Interval contents. If the Free Space in the buffer is not in agreement with the amount of Free Space, as calculated from the records found within the Control Interval and the temporary storage control blocks, corruption of some kind has occurred since this Control Interval was previously compressed. If such data integrity loss is detected, CICS issues a DFHTS1310 abend.

CONTROL INTERVAL SELECTION AND BUFFER COMPRESSION

A user migrated his CICS environment from CICS/ESA Release 4.1.0 to CICS Transaction Server. The workload and temporary storage definitions were kept constant between the two versions, and yet various puzzling differences in the temporary storage statistics data were observed.

Figure 1 shows some of the statistics fields produced by the temporary storage component of CICS, firstly from a CICS/ESA Release 4.1.0 region and then from a CICS Transaction Server region. As can be

	<i>CICS/ESA</i>	<i>CICS Transaction Server</i>
Total number of transactions	334,840	307,689
Put/Putq auxiliary storage	770,619	701,072
Get/Getq auxiliary storage	644,097	624,229
Peak temporary storage	5,780	5,619
Times queues created	286,558	272,469
Control Interval size	12,288	12,288
Longest auxiliary temp	8,048	8,024
Number of Control Intervals	24,000	23,999
Peak Control Intervals in use	18,006	21,793
Number of temp storage	396,505	20,137

Figure 1: Comparison of statistics

seen, the workloads and temporary storage usage on both systems were very similar.

The fields that surprised the user were the ratio of ‘Number of Control Intervals available’ to the ‘Peak Control Intervals in use’. In CICS/ESA Release 4.1.0 this ratio was 75%, in CICS Transaction Server it was over 90%. Also, the ‘Number of temp storage compressions’ in CICS Transaction Server (that is, Control Interval compressions within temporary storage buffers) was only 5% of the value it used to be in CICS/ESA Release 4.1.0.

Taking the ratio of peak Control Intervals in use to total number of Control Intervals available first, it needs to be understood how CICS/ESA Release 4.1.0 managed the selection of Control Intervals to satisfy EXEC CICS WRITEQ requests. In particular, the notion of the ‘75 percent rule’ must be explained.

EXEC CICS WRITEQ requests on a cold-started CICS/ESA Release

4.1.0 system are allocated into Control Intervals from the beginning of the DFHTEMP dataset. The first Control Interval is used for storing data for each WRITEQ command that is issued. Successive data records are written into the Control Interval contiguously from the start. This process continues until a WRITEQ request is made with data too large to fit the remaining Free Space within the Control Interval. At this point, CICS/ESA Release 4.1.0 switches to use the second Control Interval, and record data is added from the start of this. In such a way, successive Control Intervals are used to store temporary storage record data.

This process continues until 75% of the Control Intervals on DFHTEMP have been written to. EXEC CICS WRITEQ requests after this point are directed back to the start of the dataset once more. CICS uses the Free Space data from the byte map to select a Control Interval with enough Free Space to accommodate the requests. It looks back to the start of the dataset because, the theory is, old queues written earlier to CICS may well have been deleted by now. As explained above, such deleted queue data will remain in Control Intervals but no longer be required. If the byte map shows a Control Interval has sufficient Free Space for the new record being written, CICS will read it into a buffer and compress it to move all the required records to the start of the Control Interval.

The reason CICS/ESA Release 4.1.0 retained 25% of Control Intervals was to provide space for special header records. These are generated internally by CICS when handling records that are larger than the Control Interval size ('spanned records'). Special header records require an empty Control Interval when being written. If DFHTEMP contained fragmented data in each Control Interval, a Special header record could not be stored. This is why CICS/ESA Release 4.1.0 tried to maintain a percentage of free Control Intervals for use by large items such as Special header records.

With CICS Transaction Server, temporary management still tries to maintain a percentage of empty Control Intervals for use in storing large records. However, the algorithm for providing this has been changed to only implement the ruling when DFHTEMP has been extended and cannot be enlarged any more. In other words, the total primary and secondary capacity of the dataset has been reached. The

reasoning behind this is that it is more efficient to allocate an empty Control Interval when writing a record to temporary storage than it is to read in a Control Interval via an I/O and then to compress it to generate a contiguous Free Space for use by new requests. Therefore, while it is still possible to enlarge the dataset and provide further empty Control Intervals for use by temporary storage, this action is preferable to the alternative of returning to the start of DFHTEMP and searching for an existing Control Interval with sufficient space to handle new WRITEQ requests.

With this understood, the second confusing statistic becomes understandable. Since new empty Control Intervals are being selected in preference to old ones being reused, the rate of reuse of existing Control Intervals is reduced. Therefore, the number of temporary storage buffer compressions is reduced.

TEMPORARY STORAGE PERFORMANCE CONSIDERATIONS

On migrating to CICS Transaction Server from CICS/ESA Release 4.1.0, a user noticed that the CPU costs and response times from their transactions were increasing during the run of CICS. The transactions were intensive users of auxiliary temporary storage. The system definitions for the numbers of buffers and strings, Control Interval sizes, and DFHTEMP attributes were all unchanged.

Analysis of the problem showed that the trace option TS=3 was being specified. With the restructure of temporary storage into a Domain, it now has its own trace component that can be set independently of the other traceable components of CICS. In CICS/ESA Release 4.1.0, temporary storage trace calls were made from within the Application Domain (AP).

Specifying TS=3 instructs CICS to perform consistency checking of its control blocks for any corruption that may have occurred. In effect, this is the same activity that could be carried out by a DFHTRAP in previous releases. In fact, such a DFHTRAP is available from IBM to perform consistency checking in CICS/ESA Release 4.1.0 – it is designed to capture the moment when a corruption occurs that could otherwise lead to a DFHTS1310 abend eventually occurring.

Having such validation work being carried out by CICS has a noticeable effect upon the pathlength of a temporary storage request, and the CPU consumption needed to achieve it. Because the consistency checking takes longer to complete, more auxiliary temporary storage data exists in the system – this explained the gradual degradation in performance and increase in CPU costs during a run of CICS.

The solution was to correct the trace option for the temporary storage trace component. TS=3 should only be set in either a test environment, or when a production system is known to be experiencing temporary storage corruption and a resolution to the problem requires such a pragmatic approach.

QIDERR VERSUS INVREQ

Prior to CICS Transaction Server Release 1, there was a restriction that prevented a temporary storage queue being created that had a null name (ie binary zeros, or X'0000000000000000').

However, it was regarded as valid for applications to attempt to read or delete such a queue – in these circumstances, CICS/ESA Release 4.1.0 would return a qiderr response to the EXEC CICS READQ or DELETEQ API request. Qiderr is a ‘soft’ response that indicates to the application that a given queue does not exist. Applications should cater for the possibility of qiderr being returned on a READQ or DELETEQ request since it is quite likely that a queue may not exist (or did once exist but has since been deleted) when such a request is made.

With CICS Transaction Server, the decision was made to tighten up the API regarding this response. Because it is not possible to create a temporary storage queue that has a null name, applications should have no valid reason for attempting to issue EXEC CICS READQ or DELETEQ requests against such a queue.

Any such requests against a queue indicate that application logic is incorrect – most probably because the variable field used to specify the queue name value was not set up properly and so left set to its initial value (typically one of binary zeros). Therefore, a stronger API response than qiderr was deemed appropriate to be returned to the

application. CICS Transaction Server therefore returns invreq in such circumstances.

It was felt appropriate to differentiate between API calls against temporary storage queues that did not currently exist, but which could validly exist at some point under CICS (ie qiderr), as opposed to requests against queues that could never have existed (ie invreq).

SHARED TEMPORARY STORAGE CONSIDERATIONS

Shared temporary storage (or temporary storage data sharing) provides multiple MVS regions in a parallel sysplex with access to CICS temporary storage queue data. Shared temporary storage queues are held within pools; each such pool corresponds to a list structure within a coupling facility.

Shared temporary storage queues are non-recoverable. However, since they are stored within the coupling facility, they are normally preserved across a CICS restart or even an MVS re-IPL.

CICS systems that use shared temporary storage gain access to a pool of queues via a temporary storage data sharing server for a given pool. All such access is achieved via cross-memory calls to the server for the pool.

There are many benefits to using shared temporary storage for queue management. These include the performance improvement of accessing queue data from a coupling facility compared to function shipping the request to another CICS system (a Queue Owning Region or QOR).

Users considering implementing shared temporary storage can refer to further information on the CICS system definition, resource definition, and security implications from the *CICS System Definition Guide*, *CICS Resource Definition Guide*, and *CICS RACF Security Guide*.

CHANGES VISIBLE FROM THE CICS DUMP FORMATTER

The restructuring of temporary storage into a CICS Domain, and the implementation of object-oriented programming techniques, have led to a number of changes visible from the CICS dump formatter.

TS=2 formats the various temporary storage control blocks as before. In CICS Transaction Server, however, this now starts with the Domain and class anchors. The TSA is the Domain anchor block, and each class within the object-oriented code has its own anchor block too. These are followed by the DTN (Digital Tree Nodes) control blocks. These represent the tree structure used to maintain temporary storage queue nodes on a CICS system.

Each queue is then broken down into its component control blocks. In CICS Transaction Server, a queue is represented by a TSQ control block. Each record on a queue has a TSI control block to describe the item. For main queues, a TSI addresses a TSM control block, and for each auxiliary queue there is a TSX.

Additional control blocks (such as those specific to auxiliary and shared queue management) are also displayed if appropriate to the CICS system that generated the dump.

The various temporary storage control blocks visible when formatting a CICS Transaction Server system dump using the 'TS' verb exit are documented in the *CICS Problem Determination Guide*.

SUMMARY AND CONCLUSIONS

I hope that this article has helped explain the background to the CICS temporary storage Domain in CICS Transaction Server, and also given an indication of some of the variations that may be encountered when comparing temporary storage activity and statistics between CICS/ESA Release 4.1.0 and CICS Transaction Server.

Editor's note: readers wishing to discuss the material in this article can contact the author via e-mail at andy_wright@uk.ibm.com. CICS is a registered trademark of International Business Machines Corporation.

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CEMT Logger – an alternative design – part 2

This month we conclude the redesigned system to record CEMT output to the CSMT transient data queue.

```
*****
P-PARSE-CEMT-OUTPUT.
* this routine handles up to 100 lines of output.
*****
MOVE L-WORKAREA-TIOALEN TO W-WORK-LENGTH.
MOVE L-WORKAREA-TIOA-DATA(1:W-WORK-LENGTH) TO W-TIOA-DATA.
MOVE +0 TO W-UNSTRING-COUNT.
UNSTRING W-TIOA-DATA
  DELIMITED BY C-SBA-CHAR
  INTO W-CEMT-OUTPUT-LINE(1)
    W-CEMT-OUTPUT-LINE(2)
    W-CEMT-OUTPUT-LINE(3)
    W-CEMT-OUTPUT-LINE(4)
    W-CEMT-OUTPUT-LINE(5)
    W-CEMT-OUTPUT-LINE(6)
    W-CEMT-OUTPUT-LINE(7)
    W-CEMT-OUTPUT-LINE(8)
    W-CEMT-OUTPUT-LINE(9)
    W-CEMT-OUTPUT-LINE(11)
    W-CEMT-OUTPUT-LINE(12)
    W-CEMT-OUTPUT-LINE(13)
    W-CEMT-OUTPUT-LINE(14)
    W-CEMT-OUTPUT-LINE(15)
    W-CEMT-OUTPUT-LINE(16)
    W-CEMT-OUTPUT-LINE(17)
    W-CEMT-OUTPUT-LINE(18)
    W-CEMT-OUTPUT-LINE(19)
    W-CEMT-OUTPUT-LINE(20)
    W-CEMT-OUTPUT-LINE(21)
    W-CEMT-OUTPUT-LINE(22)
    W-CEMT-OUTPUT-LINE(23)
    W-CEMT-OUTPUT-LINE(24)
    W-CEMT-OUTPUT-LINE(25)
    W-CEMT-OUTPUT-LINE(26)
    W-CEMT-OUTPUT-LINE(27)
    W-CEMT-OUTPUT-LINE(28)
    W-CEMT-OUTPUT-LINE(29)
    W-CEMT-OUTPUT-LINE(30)
    W-CEMT-OUTPUT-LINE(31)
    W-CEMT-OUTPUT-LINE(32)
    W-CEMT-OUTPUT-LINE(33)
    W-CEMT-OUTPUT-LINE(34)
    W-CEMT-OUTPUT-LINE(35)
    W-CEMT-OUTPUT-LINE(36)
```

W-CEMT-OUTPUT-LINE(37)
W-CEMT-OUTPUT-LINE(38)
W-CEMT-OUTPUT-LINE(39)
W-CEMT-OUTPUT-LINE(40)
W-CEMT-OUTPUT-LINE(41)
W-CEMT-OUTPUT-LINE(42)
W-CEMT-OUTPUT-LINE(43)
W-CEMT-OUTPUT-LINE(44)
W-CEMT-OUTPUT-LINE(45)
W-CEMT-OUTPUT-LINE(46)
W-CEMT-OUTPUT-LINE(47)
W-CEMT-OUTPUT-LINE(48)
W-CEMT-OUTPUT-LINE(49)
W-CEMT-OUTPUT-LINE(50)
W-CEMT-OUTPUT-LINE(51)
W-CEMT-OUTPUT-LINE(52)
W-CEMT-OUTPUT-LINE(53)
W-CEMT-OUTPUT-LINE(54)
W-CEMT-OUTPUT-LINE(55)
W-CEMT-OUTPUT-LINE(56)
W-CEMT-OUTPUT-LINE(57)
W-CEMT-OUTPUT-LINE(58)
W-CEMT-OUTPUT-LINE(59)
W-CEMT-OUTPUT-LINE(60)
W-CEMT-OUTPUT-LINE(61)
W-CEMT-OUTPUT-LINE(62)
W-CEMT-OUTPUT-LINE(63)
W-CEMT-OUTPUT-LINE(64)
W-CEMT-OUTPUT-LINE(65)
W-CEMT-OUTPUT-LINE(66)
W-CEMT-OUTPUT-LINE(67)
W-CEMT-OUTPUT-LINE(68)
W-CEMT-OUTPUT-LINE(69)
W-CEMT-OUTPUT-LINE(70)
W-CEMT-OUTPUT-LINE(71)
W-CEMT-OUTPUT-LINE(72)
W-CEMT-OUTPUT-LINE(73)
W-CEMT-OUTPUT-LINE(74)
W-CEMT-OUTPUT-LINE(75)
W-CEMT-OUTPUT-LINE(76)
W-CEMT-OUTPUT-LINE(77)
W-CEMT-OUTPUT-LINE(78)
W-CEMT-OUTPUT-LINE(79)
W-CEMT-OUTPUT-LINE(80)
W-CEMT-OUTPUT-LINE(81)
W-CEMT-OUTPUT-LINE(82)
W-CEMT-OUTPUT-LINE(83)
W-CEMT-OUTPUT-LINE(84)
W-CEMT-OUTPUT-LINE(85)
W-CEMT-OUTPUT-LINE(86)
W-CEMT-OUTPUT-LINE(87)
W-CEMT-OUTPUT-LINE(88)

```

W-CEMT-OUTPUT-LINE(89)
W-CEMT-OUTPUT-LINE(90)
W-CEMT-OUTPUT-LINE(91)
W-CEMT-OUTPUT-LINE(92)
W-CEMT-OUTPUT-LINE(93)
W-CEMT-OUTPUT-LINE(94)
W-CEMT-OUTPUT-LINE(95)
W-CEMT-OUTPUT-LINE(96)
W-CEMT-OUTPUT-LINE(97)
W-CEMT-OUTPUT-LINE(98)
W-CEMT-OUTPUT-LINE(99)
W-CEMT-OUTPUT-LINE(100)
TALLYING IN W-UNSTRING-COUNT
ON OVERFLOW
MOVE 'CEMT OUTPUT TRUNCATED' TO W-MSG-TEXT
MOVE '10' TO W-MSG-NO
PERFORM P-HANDLE-ERROR
END-UNSTRING.
PERFORM P-WRITE-TIOA
VARYING K FROM +1 BY +1 UNTIL K > W-UNSTRING-COUNT.

```

P-WRITE-TIOA.

```

MOVE W-CEMT-OUTPUT-DATA(K) TO W-TDQ-BUFFER.
MOVE '**' TO W-TDQ-BUFFER(1:2).
SUBTRACT +1 FROM LENGTH OF W-TDQ-BUFFER
GIVING W-WORK-LENGTH.
PERFORM VARYING J FROM +3 BY +1
UNTIL J > W-WORK-LENGTH
EVALUATE TRUE
  WHEN W-TDQ-BUFFER(J:1) = C-INSERT-CURSOR
    MOVE SPACE TO W-TDQ-BUFFER(J:1)
  WHEN W-TDQ-BUFFER(J:1) = X'00'
    MOVE SPACE TO W-TDQ-BUFFER(J:1)
  WHEN W-TDQ-BUFFER(J:1) = C-START-FIELD-CHAR
    MOVE SPACES TO W-TDQ-BUFFER(J:2)
    ADD +1 TO J
  WHEN W-TDQ-BUFFER(J:1) = LOW-VALUE
    MOVE SPACE TO W-TDQ-BUFFER(J:1)
  WHEN OTHER
    CONTINUE
END-EVALUATE
END-PERFORM.
EVALUATE TRUE

```

***** DON'T SHOW ANYTHING AFTER THE 'STATUS:' LINE

```

WHEN W-TDQ-BUFFER(3:7) = 'STATUS:'
  PERFORM P-WRITE-TDQ
  ADD W-UNSTRING-COUNT TO K
WHEN W-TDQ-BUFFER(3:77) = SPACES
  CONTINUE
WHEN W-TDQ-BUFFER(3:77) = LOW-VALUES

```

```
CONTINUE
WHEN OTHER
PERFORM P-WRITE-TDQ
END-EVALUATE.
```

```
*****
P-INQUIRE-REQID.
```

```
*****
MOVE C-REQID TO W-REQID.
```

```
EXEC CICS INQUIRE
REQID(W-REQID)
NOHANDLE
END-EXEC.
```

```
*****
P-CANCEL-REQID.
```

```
*****
PERFORM WITH TEST AFTER UNTIL EIBRESP NOT = DFHRESP(NORMAL)
EXEC CICS CANCEL
REQID(C-REQID)
NOHANDLE
END-EXEC
END-PERFORM.
```

```
*****
P-START-AGAIN.
```

```
*****
```

```
**** FIRST CANCEL ANY OUTSTANDING REQIDS WITH THE SAME NAME.
```

```
**** THIS PREVENTS INADVERTENT 'SPAWNING' PROBLEMS
```

```
PERFORM P-CANCEL-REQID.
EXEC CICS START
TRANSID(C-MONITOR-TRANS-ID)
AFTER
MINUTES(W-INTERVAL-BIN-MINS)
REQID(C-REQID)
RESP(W-RESP)
RESP2(W-RESP2)
END-EXEC.
```

```
IF W-RESP NOT = DFHRESP(NORMAL) THEN
MOVE 'START TRANSID ERROR' TO W-MSG-TEXT
MOVE '11' TO W-MSG-NO
PERFORM P-HANDLE-ERROR
GO TO 0000-CICS-RETURN
END-IF.
```

```
*****
P-HANDLE-ERROR.
```

```
*****
```

```
IF W-RESP NOT = DFHRESP(NORMAL) THEN
MOVE W-RESP TO W-RESP-PIC
MOVE W-RESP2 TO W-RESP2-PIC
```

```

        STRING
            W-MSG-TEXT                                DELIMITED BY ' '
            ',RESP=' W-RESP-PIC                       DELIMITED BY SIZE
            ',RESP2=' W-RESP2-PIC                    DELIMITED BY SIZE
            INTO W-MSG-TEXT
        END-STRING
    END-IF.
    PERFORM P-WRITE-MSG.

```

```

*****
P-WRITE-TDQ.
*****

```

```

        EXEC CICS WRITEQ TD QUEUE(C-MSG-QUEUE)
            FROM(W-TDQ-BUFFER)
            LENGTH(LENGTH OF W-TDQ-BUFFER)
            NOHANDLE
        END-EXEC.
        MOVE SPACES TO W-TDQ-BUFFER.

```

```

*****
P-WRITE-MSG.
*****

```

```

        PERFORM P-GET-TIMESTAMP.
        MOVE W-TIME TO W-MSG-TIME.
        MOVE W-MSG TO W-TDQ-BUFFER.
        PERFORM P-WRITE-TDQ.
        MOVE SPACES TO W-MSG-TEXT.

```

```

*****
P-GET-TIMESTAMP.
*****

```

```

        EXEC CICS ASKTIME
            ABSTIME(W-ABSTIME)
            NOHANDLE
        END-EXEC.
        EXEC CICS FORMATTIME
            ABSTIME(W-ABSTIME)
            DDMMYYYY(W-DDMMYYYY)
            DATESEP('/')
            TIME(W-TIME)
            TIMESEP(':')
            NOHANDLE
        END-EXEC.

```

```

*****

```

ZZZZCOUT

```

*****
*
*   MODULE NAME = ZZZZCOUT
*
*****

```

```

* DESCRIPTIVE NAME = CICS/ESA XZCOUT GLOBAL USER EXIT PROGRAM      *
*                                                                 *
* FUNCTION - CAPTURE CEMT OUTPUT FOR EVENTUAL WRITING TO CSMT     *
* LOG BY ZZZCEMT (Q.V.).                                          *
*                                                                 *
* GWA LENGTH = 91 BYTES                                          *
*                                                                 *
* STATUS = 4.1.Ø                                                *
*                                                                 *
* ERROR CONDITION CODES ARE SAVED IN THE GWA.                    *
* PROGRAM ZZZCEMT CAN CHECK THIS CODE AND TAKE APPROPRIATE      *
* NOTIFICATION ACTION.                                           *
*                                                                 *
*-----*
*-----*

```

EJECT

```

*
* USER EXIT INTERFACE FOR EXIT POINT XZCOUT
      DFHUEXIT TYPE=EP, ID=(XZCOUT)
      DFHUEXIT TYPE=XPIENV                USED FOR XPI CALL
*
* TERMINAL INPUT/OUTPUT AREA (USED BY CEMT)
      COPY      DFHTIOA
*
* TERMINAL CONTROL TABLE
      PRINT     OFF,NOGEN
      COPY      DFHTCTTE
      PRINT     ON,GEN
*
* XPI DSECT FOR INQUIRE TRANSACTION CALL
      COPY      DFHXMIQY
*
* GWA AREA
GWA      DSECT
GWA_DWORD      DS  D          WORK
GWA_WORKPTR1   DS  F          -> WORKAREA 1
GWA_WORKPTR2   DS  F          -> WORKAREA 2
GWA_OFFSET     DS  F          OFFSET IN WORKAREA
GWA_TIOALEN    DS  F          TIOA LENGTH
GWA_TERMID     DS  CL4        TERMINAL-ID
GWA_TRANID     DS  CL4        TRANSACTION-ID
GWA_USERID     DS  CL8        USER-ID
GWAERRNO       DS  XL1        ERROR NO.
ALL_OK         EQU X'ØØ'      NO WORRIES
ERR_NO_WORKAREA EQU X'Ø1'      NO VALID WORKAREA
ERR_XPI_CALL   EQU X'Ø2'      XPI CALL FAILED (INQ TRANS)
ERR_BUFFER_FULL EQU X'Ø3'      TIOA WON'T FIT IN BUFFER
ERR_TCTTE_NULL EQU X'Ø4'      TCTTE INVALID
GWAMSG         DS  CL49        FREE FOR ERROR MESSAGE
GWA_MAX_LEN    EQU 65536      SIZE OF EACH WORKAREA
GWA_STATUS     DS  XL1        WORKAREA STATUS

```



```

GWA_INITIAL    EQU X'00'          INITIAL STATE - NOT READY
GWA_WORK1      EQU C'1'          WORKAREA1 ACTIVE
GWA_WORK2      EQU C'2'          WORKAREA2 ACTIVE
*
GWALEN    EQU    *-GWA          LENGTH OF GWA
*
* LENGTH OF HEADER FIELDS FOR SAVED CEMT DATA
* USERID + TERMID + TIOA LENGTH (PACKED)
HDRLEN    EQU    L'GWA_USERID+L'GWA_TERMID+3
*
*****
* REGISTER EQUATES
*****
*R1        EQU    1             UEP BASE ON ENTRY, WORK
UEPBAR     EQU    2             -> USER EXIT PARMS
GWABAR     EQU    3             -> GLOBAL WORK AREA
*R4        EQU    4             WORK
*R5        EQU    5             WORK
*R6        EQU    6             WORK
*R7        EQU    7             WORK
WORKPTR    EQU    8             WORKAREA PTR
TIOABAR    EQU    9             -> TIOA
TCTTEAR    EQU    10           -> TCTTE
BASEREG    EQU    11           PROGRAM BASE
*R12       EQU    12           WORK, UEPXSTOR BASE
*R13       EQU    13           SAVEAREA POINTER
*R14       EQU    14           WORK
LINKREG    EQU    14           SUBROUTINE LINKAGE
*R15       EQU    15           ENTRY ADDRESS, RETURN CODE
EJECT
*
*****
ZZZCOUT    AMODE 31
ZZZCOUT    RMODE ANY
ZZZCOUT    CSECT
          SAVE (14,12)          SAVE REGISTERS
          LR    BASEREG,R15
          USING ZZZCOUT,BASEREG SET UP PROGRAM BASE REGISTER
          B     START           GO AROUND EYECATCHER
          DC    C'*ZZZCOUT*'
          DC    C'*)&SYSDATE*'
          DC    C'*)&SYSTIME*'
          DC    C'*)&VERSION 1.0*'
*
START      DS    0H
          LR    UEPBAR,R1
          USING DFHUEPAR,UEPBAR -> USER EXIT PARAMETER LIST
*
* CHECK GWA
          L     R1,UEPGAL        -> GWA LENGTH
          CLC   0(2,R1),=AL2(GWALEN) GWA LENGTH OK ?

```

```

BNE RETURN NO, GET OUT ->
ICM GWABAR,B'1111',UEPGAA GET GWA ADDRESS
BZ RETURN IF BAD, GET OUT ->
USING GWA,GWABAR -> GWA
*
* CHECK THAT THE WORKAREA POINTER IS VALID
CLI GWA_STATUS,GWA_INITIAL OPEN FOR BUSINESS ?
BE RETURN NO, EXIT ->
CLI GWA_STATUS,GWA_WORK1 WORKAREA 1 ?
BE WORKAREA1 YES ->
CLI GWA_STATUS,GWA_WORK2 WORKAREA 2 ?
BE WORKAREA2 YES ->
B NO_WORKAREA ELSE PROBLEM ->
*
WORKAREA1 DS 0H
ICM WORKPTR,B'1111',GWA_WORKPTR1 GET WORKAREA 1 ADDR
BZ NO_WORKAREA BAD ADDR ??? ->
B FIND_TERMINAL ELSE CONTINUE ->
*
WORKAREA2 DS 0H
ICM WORKPTR,B'1111',GWA_WORKPTR2 GET WORKAREA 2 ADDR
BZ NO_WORKAREA BAD ADDR ??? ->
*
FIND_TERMINAL DS 0H
ICM TCTTEAR,B'1111',UEPTCTTE -> TCTTE
BZ NO_TCTTE EXIT IF NONE ->
MVC GWA_TERMID,TCTTETI SAVE TERMID
*
* FIND THE USER-ID USING XPI CALL
L R12,UEPXSTOR -> XPI STORAGE
LA R12,4(R12) LEAVE 1ST 4 BYTES
USING DFHXMIQ_ARG,R12
L R13,UEPSTACK -> KERNEL STACK
DFHXMIQX CALL,CLEAR, XPI CALL INQ TRANSACTN +
IN, INPUT PARMS +
FUNCTION(INQUIRE_TRANSACTION), +
OUT, OUTPUT PARMS +
TRANSACTION_ID(GWA_TRANID), +
USERID(GWA_USERID), +
RESPONSE(*), +
REASON(*)
CLI XMIQ_RESPONSE,XMIQ_OK XPI CALL OK ?
BNE XMIQ_ERROR NO ->
*
* IF TRANSACTION IS NOT 'CEMT', EXIT
CLI GWA_TRANID,C'C' C TRANSID ?
BNE RETURN NO ->
CLC GWA_TRANID+1(3),=CL3'EMT' 'CEMT' ?
BNE RETURN NO ->
*
* COPY THE TIOA DATA TO THE ACTIVE WORKAREA. 1ST 8 BYTES IS USER-ID,
* THEN TIOA DATA LENGTH (2 BYTES), THEN TIOA DATA.

```

```

ICM  TIOABAR,B'1111',UEPTIOA      -> TIOA
BZ   RETURN                        GET OUT IF NO TIOA
XR   R7,R7
ICM  R7,B'0011',TIOATDL           GET TIOA DATA LENGTH
BZ   RETURN                        GET OUT IF NO DATA
ST   R7,GWA_TIOALEN              SAVE TIOA LENGTH
L    R1,GWA_OFFSET               GET OFFSET
AR   R1,R7                       OFFSET + TIOA LENGTH
LA   R1,HDRLEN(R1)               + USERID + LEN
C    R1,=AL4(GWA_MAX_LEN)        > BUFFER SIZE ?
BH   BUFFER_FULL                 YES ->
*
LR   R5,R7                       SET LENGTH
LA   R6,TIOADBA                  -> TIOA DATA
LR   R4,WORKPTR                  -> BUFFER
A    R4,GWA_OFFSET               -> WRITE AREA
MVC  0(8,R4),GWA_USERID          SAVE USER-ID
MVC  8(4,R4),GWA_TERMID         SAVE TERM-ID
L    R1,GWA_TIOALEN             GET TIOA LENGTH
XC   GWA_DWORD,GWA_DWORD        CLEAR WORK FIELD
CVD  R1,GWA_DWORD               CONVERT TO DECIMAL
MVC  8+4(3,R4),GWA_DWORD+8-3    SAVE LAST 3 BYTES
LA   R4,HDRLEN(R4)             BUMP PTR
MVCL R4,R6                      COPY TIOA
*
* RECALCULATE NEW OFFSET
L    R4,GWA_OFFSET              GET OFFSET
LA   R4,HDRLEN(R4)             ADD LEN + USER-ID ETC
A    R4,GWA_TIOALEN            + TIOA LEN = OFFSET
ST   R4,GWA_OFFSET             SAVE IT
*
B    RETURN                     AND EXIT ->
EJECT
*
*****
* RETURN TO CICS
*****
RETURN DS  0H
L      R13,UEPEPSA              RESTORE R13
LA     R15,UERCNORM             SET RC = OK
RETURN (14,12),RC=(15)
EJECT
*
*****
* ERROR HANDLING
*****
NO_TCTTE DS  0H                 TCTTE ADDRESS ZERO
MVI    GWAERRNO,ERR_TCTTE_NULL  FLAG ERROR
MVC    GWAMSG,=CL(L'GWAMSG)'NO TCTTE '
B      RETURN                   AND EXIT ->
*
BUFFER_FULL DS  0H              TIOA WON'T FIT

```

```

MVI   GWAERRNO,ERR_BUFFER_FULL           FLAG ERROR
MVC   GWAMSG,=CL(L'GWAMSG)'BUFFER FULL'
B     RETURN                               AND EXIT ->
*
XMIQ_ERROR DS 0H                           XPI CALL FAILED
MVI   GWAERRNO,ERR_XPI_CALL              FLAG ERROR
MVC   GWAMSG,=CL(L'GWAMSG)'XPI CALL FAILED'
B     RETURN                               AND EXIT ->
*
NO_WORKAREA DS 0H
MVI   GWAERRNO,ERR_NO_WORKAREA           NO VALID WORKAREA
MVC   GWAMSG,=CL(L'GWAMSG)'INVALID WORKAREA'
B     RETURN                               AND EXIT ->
EJECT
*
*****
CONSTNTS DS 0F
LTORG
END   ZZZZCOUT

```

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Displaying CPU usage by TCB – revisited

The code in the article *Displaying CPU usage by TCB*, published in *CICS Update*, June 1999 and July 1999, Issues 163 and 164, contained some errors.

The following amendments should be noted:

- In Issue 163, page 26, after the ‘Important: pictures should not be modified’ statement in the IPPCDTCB program, the line:

```
01 IPPCGTCB PIC X(08) VALUE 'IPPCDTCB'
```

should read:

```
01 IPPCGTCB PIC X(08) VALUE 'IPPCGTCB'
```

- In Issue 164, pages 39 and 40, the same program has two ‘PERFORM ACCESS-CALTAB’ lines. These should be ‘PERFORM ACCESS-GTCB’.

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Accessing CICS control blocks in COBOL

Many recently trained programmers and other IT industry workers often treat COBOL with contempt. However, on the mainframe the majority of existing applications are written in COBOL. More significant is the fact that new applications continue to be developed in COBOL – so the language must have some good points!

Recently, I have found myself at various client sites where I needed to examine some of the data in CICS control blocks. The main problem of being a consultant is that each environment you encounter is different – you cannot rely on the availability of an Independent Software Vendor (ISV) tool. However, almost always the installation has COBOL. So I decided to write a program I could take with me to sites to use for the purpose of examining internal CICS structures in COBOL that I knew would usually be available wherever I went. I found this to be much easier than might be thought.

THE SOURCE

Since the run-time environment available at different places is also different, I decided that it would be best to take the program in source code form. The application is extremely simple, consisting of a BMS mapset, a program, and three resource definitions. Theoretically, the program and mapset definitions are not required if program auto-install is being used (PGAIPGM=ACTIVE in the SIT), but it should be noted that the program must run in CICS key since it references areas of CICS key storage which are fetch-protected.

I am currently expanding the program and mapset in order to display more information, but the code included with this article could easily be modified and extended for your own purposes. Note that the code is written to be able to run in any CICS/ESA environment from Version 4.1 through to the latest available release, Version 5.3 (the CICS component of Transaction Server 1.3). It has been tested on all of those releases with the exception of 5.1, although I do not anticipate any problems in that version.

THE RESOURCE DEFINITIONS

The names used for these resources can be anything that conforms to your naming conventions. The illustrated definitions only show the relevant parameters; the normal CICS defaults can be assumed for all other values. All definitions must, of course, be placed in an appropriate group and installed.

Transaction ADDR

TRANSaction ==> addr
DEscription ==> control block display
PROGram ==> addrdisp
TASKDATAloc ==> Any
SPurge ==> Yes
TPUrge ==> Yes

Program ADDRDISP

PROGram ==> addrdisp
Description ==> control block display
DATAlocation ==> Any
EXECKey ==> CICS

Mapset ADDRMAP

Mapset ==> ADDRMAP
Description ==> Control Block Maps

ADDRDISP

IDENTIFICATION DIVISION.

PROGRAM-ID. ADDRDISP.

*

* Address manipulation – MVS & CICS control blocks

*

* This program will work in all versions of CICS/ESA from

* V4.1 through 5.3 (CICS Transaction Server 1.3). It is

* constructed in such a manner as to allow it to be

* adapted to later releases without great difficulty.

* All release-specific data names and values have a prefix

* of 'Vvr-' where 'v' is the Version number and 'm' the

* Release level, eg 'V41-'. Note that the Modification

* level is ignored when interrogating CICS as to its RELEASE

* via the INQUIRE SYSTEM command; this is based on the

* assumption that no general fix will change the basic

* number of Domains nor the number of gates allowed per

* Domain. However, newer releases may change some of the
 * control structures referenced by the program, so careful
 * review of the LINKAGE SECTION would need to be performed
 * as part of any adaptation to a specific release.

*
 ENVIRONMENT DIVISION.
 CONFIGURATION SECTION.
 SPECIAL-NAMES.

*
 * The following is used for determining which characters
 * are "printable" when displaying data areas.
 *

```

CLASS PUNCTUATION IS '$' '.' '<' '(' '+' '|' '&' '!' '£'
                      '*' ')' ';' '¬' '-' '/' '≥' ', ' %'
                      '_' '>' '?' '`' ':' '#' '@' '""'
                      '{' '}' '\'.
  
```

DATA DIVISION.
 WORKING-STORAGE SECTION.
 01 FILLER.

*
 * NB, If the number of Domains ever exceeds 40, then the
 * ADDRMAP, the following value, and the definition of
 * KCB-VECTOR in the LINKAGE SECTION below must change.
 *

```

03 MAX-IN-MAP          PIC 9(04) COMP VALUE 40.
03 MAX-NO-OF-DOMAINS  PIC 9(04) COMP VALUE 26.
88 V41-NO-DOMAINS     VALUE 26.
88 V51-NO-DOMAINS     VALUE 29.
88 V52-NO-DOMAINS     VALUE 29.
88 V53-NO-DOMAINS     VALUE 36.
03 WHAT-VERSION.
05 WV                 PIC X(03).
88 V41                 VALUE '041'.
88 V51                 VALUE '051'.
88 V52                 VALUE '052'.
88 V53                 VALUE '053'.
05 FILLER             PIC X(01).
  
```

*
 * The following is used to format an entry on the main
 * Domain information display panel.

```

01 OUT-LINE           VALUE SPACES.
03 DD-ID             PIC X(02).
03 FILLER            PIC X(01).
03 DD-IX            PIC 9(02).
03 FILLER            PIC X(01).
03 DD-ADDR          PIC X(08).
03 FILLER            PIC X(01).
  
```

*
 * The following is used to format a line of data on

* the detail display panel.

*

```
Ø1 OUT-ADDR          VALUE SPACES.
   Ø3 OA-ADDR        PIC X(Ø8).
   Ø3 FILLER         PIC X(Ø3).
   Ø3 OA-OFF        PIC X(Ø6).
   Ø3 FILLER         PIC X(Ø3).
   Ø3 OA-CORE       OCCURS 4.
       Ø5 OA-DATA    PIC X(Ø8).
       Ø5 FILLER     PIC X(Ø1).
   Ø3 FILLER         PIC X(Ø2).
   Ø3 OA-EBCDIC     PIC X(16).
   Ø3 FILLER         PIC X(Ø5).
```

*

* These are the various work areas required by the logic.

*

```
Ø1 FILLER.
   Ø3 LINE-LIMIT    PIC 9(Ø4) BINARY.
   Ø3 SEG-LIMIT     PIC 9(Ø4) BINARY.
   Ø3 BYTE-LIMIT    PIC 9(Ø4) BINARY.
   Ø3 WORK-LIMIT    PIC 9(Ø4) BINARY.
   Ø3 DATA-LIMIT   PIC 9(Ø4) BINARY.
   Ø3 LENGTH-3      PIC 9(Ø4) BINARY VALUE 3.
   Ø3 LENGTH-4      PIC 9(Ø4) BINARY VALUE 4.
   Ø3 WORK-OFF      PIC 9(Ø8) BINARY.
   Ø3 WORK-OFF-X REDEFINES WORK-OFF.
       Ø5 FILLER    PIC X(Ø1).
       Ø5 WO-LOW.
           Ø7 WOL-BYTE PIC X(Ø1) OCCURS 3.
   Ø3 WORK-LTH      PIC 9(ØØ4) BINARY.
   Ø3 WORK-PTR      PIC 9(ØØ8) BINARY.
   Ø3 WORK-PTR-X REDEFINES WORK-PTR.
       Ø7 WP-BYTE    PIC X(Ø1) OCCURS 4.
   Ø3 ADDR-PTR REDEFINES WORK-PTR
       POINTER.
   Ø3 VECTOR-IX     PIC 9(Ø4) COMP.
   Ø3 IX            PIC 9(Ø4) COMP.
   Ø3 IW            PIC 9(Ø4) COMP.
   Ø3 MSGNO         PIC 9(ØØ4) BINARY.
   Ø3 MSGS.
       Ø5 FILLER    PIC X(6Ø) VALUE
           'Only <PF3>, <PF7> & <PF8> are valid.'.
       Ø5 FILLER    PIC X(6Ø) VALUE
           'At last page.'.
       Ø5 FILLER    PIC X(6Ø) VALUE
           'At first page.'.
   Ø3 MSG REDEFINES MSGS PIC X(6Ø)
       OCCURS 3.
   Ø3 FOUND-IND     PIC X(Ø1) VALUE 'N'.
       88 FOUND-ONE VALUE 'Y'.
```



```

      Ø3  END-MSG                PIC X(26) VALUE
          '      Processing Terminated.'
*
* As per usual, this program requires a BMS mapset.
*
COPY ADDRMAP.
COPY DFHBMSCA.
COPY DFHAID.
*
LINKAGE SECTION.
*
* Since two different maps are used for the two
* different types of data displayed, some info needs
* to be saved between the pseudo-conversational tasks.
* Also when data is being displayed, other data is
* required in order to control the paging.
*
Ø1  DFHCOMMAREA.
      Ø3  LAST-DISPLAYED        PIC X.
          88  MAIN-SHOWN        VALUE 'M'.
          88  DETAIL-SHOWN      VALUE 'D'.
      Ø3  DOING                 PIC X.
          88  INPUT-MAIN        VALUE 'M'.
          88  INPUT-DETAIL      VALUE 'D'.
      Ø3  DOMAIN                PIC X(ØØ2).
      Ø3  LAST-CURSOR           PIC 9(ØØ4) BINARY.
      Ø3  START-WORK            PIC 9(ØØ8) BINARY.
      Ø3  START-ADDR REDEFINES START-WORK
          POINTER.
      Ø3  TOTAL-LTH             PIC 9(ØØ4) BINARY.
      Ø3  CURRENT-PAGE          PIC 9(ØØ4) BINARY.
      Ø3  TOTAL-PAGES           PIC 9(ØØ4) BINARY.
*
* The following fields describe various MVS (OS/39Ø) and
* CICS/ESA (CICS/TS) control structures. These are release-
* dependent, so need to be reviewed as part of any
* adaptation to any newer release of either.
*
* The PSA is a description of the MVS (OS/39Ø) Prefixed Storage
* Area (low storage). This area is most unlikely to change since
* it is used by all CICS programs to locate DFHEIP.
*
Ø1  PSA.
      Ø3  FILLER                 PIC X(54Ø).
      Ø3  PSA-TCB-PTR            POINTER.
      Ø3  FILLER                 PIC X(ØØ4).
      Ø3  PSA-ASCB-PTR          POINTER.
*
* The ASCB is a description of the MVS (OS/39Ø)
* Address Space Control Block.

```

```

*
  01 ASCB.
    03 FILLER          PIC X(336).
    03 ASCB-ASSB-PTR  POINTER.
*
* The ASSB is a description of the MVS (OS/390)
* Address Space Secondary Block.
*
  01 ASSB.
    03 FILLER          PIC X(168).
    03 ASSB-JSAB-PTR  POINTER.
*
* The JSAB is a description of the MVS (OS/390)
* Job Scheduler Address space Block.
*
  01 JSAB.
    03 FILLER          PIC X(020).
    03 JSAB-JOBID     PIC X(008).
    03 JSAB-JOBNAME   PIC X(008).
*
* The TCB is a description of the MVS (OS/390) Task
* Control Block. This area is most unlikely to change
* since it is used by all CICS programs to locate DFHEIP.
*
  01 TCB.
    03 FILLER          PIC X(208).
    03 TCB-TCBEXT-PTR POINTER.
*
* The TCBEXT is a description of the MVS (OS/390) Task
* Control Block EXTension. This area is most unlikely
* to change since it is used by all CICS programs to
* locate DFHEIP.
*
  01 TCBEXT.
    03 FILLER          PIC X(020).
    03 TCBEXT-AFCB-PTR POINTER.
*
* The AFCB is a description of the CICS Authorized Function
* Control Block. The structure actually consists of a
* prefix, a vector list (a set of addresses), and a trailer.
* There is no direct pointer to the trailer but it can be
* found by adding the lengths of the prefix and the vector
* list to the address of the AFCB itself. The structure
* of this area is most unlikely to change since it is used
* by all CICS programs to locate DFHEIP.
*
  01 AFCB.
    03 FILLER          PIC X(006).
    03 AFCB-VLIST-LTH PIC 9(004) BINARY.
    03 FILLER          PIC X(008).

```

*
 * The AFT is a description of the CICS Authorized Function
 * Trailer. This area is most unlikely to change.
 *

```

Ø1 AFT.
  Ø3 FILLER                PIC X(ØØ4).
  Ø3 AFT-AFCS-PTR          POINTER.
  
```

*
 * The AFCS is a description of the CICS Authorized Function
 * Common Structure. This area is most unlikely to change.
 *

```

Ø1 AFCS.
  Ø3 FILLER                PIC X(ØØ8).
  Ø3 AFCS-KCB-PTR          POINTER.
  
```

*
 * The KCB is a description of the CICS/ESA Kernel Anchor Block.
 * Although this structure is essentially the same in the releases
 * on which this program was developed, note that the size of the
 * array of vectors (addresses) to the Domain table entries is
 * release-dependent. Also note that the offset (FILLER) to the
 * array did change between Version 3 and Version 4.
 *

```

Ø1 KCB.
  Ø3 FILLER                PIC X(376).
  Ø3 KCB-ERROR-VECTOR      POINTER.
  Ø3 KCB-VECTOR            POINTER
                           OCCURS 26 TO 4Ø
                           DEPENDING ON MAX-NO-OF-DOMAINS.
  
```

*
 * The DOMAIN-TABLE is a description of the CICS/ESA Kernel
 * Domain table entry essential to the linkage architecture.
 * The part of this area described is most unlikely to change.
 *

```

Ø1 DOMAIN-TABLE.
  Ø3 FILLER                PIC X(ØØ3).
  Ø3 DT-ID                 PIC X(ØØ2).
  Ø3 FILLER                PIC X(ØØ3).
  Ø3 DT-IX                 PIC 9(ØØ8) BINARY.
  Ø3 FILLER                PIC X(ØØ4).
  Ø3 DT-ANCHOR-PTR         POINTER.
  Ø3 DT-ANCHOR-PTR-X REDEFINES DT-ANCHOR-PTR
                           PIC 9(ØØ8) BINARY.
  
```

*
 * The ANCHOR is a generic description of 256 bytes of any
 * Domain's anchor block. The standard used by Domains is
 * for the length of the control block to be placed in the
 * first halfword of the block itself. There are, however
 * a couple of exceptions to this rule.
 *

* First the AP (Application) Domain's anchor block has an

* historical structure based on the original design of CICS
 * in the late 1960s. This area is more commonly known as the
 * Common System Area (CSA). For CICS/ESA this has been adapted
 * so that it now contains a length in the second halfword of
 * the CSA. However the length does not actually reflect the
 * length as described by the data areas manual. This program
 * uses the length found in the CSA.

*
 * Second the RX (Recovery Services - introduced in 5.3)
 * Domain's anchor block does not follow the standard
 * convention. There is apparently no length contained within
 * the control block itself, so the program assumes a size of
 * 4096 (4K) bytes.

*
 Ø1 ANCHOR.
 Ø3 A-BLOCK.
 Ø5 A-LENGTH PIC 9(ØØ4) BINARY.
 Ø5 A-AP-LTH PIC 9(ØØ4) BINARY.
 Ø5 FILLER PIC X(252).
 Ø3 FILLER REDEFINES A-BLOCK.
 Ø5 A-LINE OCCURS 16.
 Ø7 A-WORD OCCURS 4.
 Ø9 A-BYTE PIC X(Ø1)
 OCCURS 4.

PROCEDURE DIVISION.

*
 * First set the release-dependent values.
 *

```
EXEC CICS INQUIRE SYSTEM
      RELEASE(WHAT-VERSION)
      NOHANDLE
END-EXEC
EVALUATE TRUE
  WHEN V41
    SET V41-NO-DOMAINS TO TRUE
  WHEN V51
    SET V51-NO-DOMAINS TO TRUE
  WHEN V52
    SET V52-NO-DOMAINS TO TRUE
  WHEN V53
    SET V53-NO-DOMAINS TO TRUE
END-EVALUATE
```

*
 * The PSA is the start of virtual storage.
 *

```
SET ADDRESS OF PSA TO NULL
IF EIBCALEN = Ø
```

*
 * Initially display the Main Domain Information.
 *

```

EXEC CICS GETMAIN
      FLENGTH(LENGTH OF DFHCOMMAREA)
      SET (ADDRESS OF DFHCOMMAREA)
END-EXEC
MOVE LOW-VALUES TO DFHCOMMAREA
MOVE 1          TO LAST-CURSOR
PERFORM SEND-MAIN
ELSE
*
* Remember what we last did.
*
MOVE LAST-DISPLAYED TO DOING
*
* Only certain Attention Identifiers (keys) are
* acceptable. So we take appropriate action.
*
EVALUATE EIBAID
      WHEN DFHENTER
*
* <Enter> is only allowed on the main display
* to specify which anchor block is to be shown.
*
IF MAIN-SHOWN
EXEC CICS RECEIVE
      MAP('MAINMAP')
      MAPSET('ADDRMAP')
      NOHANDLE
END-EXEC
*
* We need to get BMS to tell us where the
* cursor was left so we can display the
* corresponding anchor block.
*
EVALUATE EIBRESP
      WHEN DFHRESP(NORMAL)
*
* Find which was requested.
*
PERFORM VARYING IX FROM 1 BY 1
UNTIL IX > MAX-NO-OF-DOMAINS
OR FOUND-ONE
      IF ANCHLNF(IX) = DFHBMCUR
      MOVE IX TO VECTOR-IX
      SET FOUND-ONE TO TRUE
      END-IF
END-PERFORM
IF FOUND-ONE
PERFORM SEND-DETAIL
ELSE
PERFORM SEND-CONTROL

```

```

        END-IF
        WHEN DFHRESP(MAPFAIL)
            PERFORM SEND-CONTROL
        WHEN OTHER
*
*
*           This should never happen, but...
*
*           EXEC CICS ABEND
*                   ABCODE('OOPS')
*                   NOHANDLE
*           END-EXEC
        END-EVALUATE
    ELSE
        MOVE 1 TO MSGNO
        PERFORM SEND-CONTROL
    END-IF
WHEN DFHPF3
*
*
*       <PF3> is used for the standard exit function.
*
*       IF MAIN-SHOWN
*           EXEC CICS SEND
*                   FROM(END-MSG)
*                   ERASE
*                   NOHANDLE
*           END-EXEC
*           EXEC CICS RETURN
*           END-EXEC
*       ELSE
*           PERFORM SEND-MAIN
*       END-IF
WHEN DFHPF8
*
*
*       <PF8> is used to scroll forward if detail data
*       is displayed and there is more to show.
*
*       IF MAIN-SHOWN
*           PERFORM SEND-CONTROL
*       END-IF
*       IF CURRENT-PAGE = TOTAL-PAGES
*           MOVE 2 TO MSGNO
*           PERFORM SEND-CONTROL
*       END-IF
*       ADD 1          TO CURRENT-PAGE
*       MOVE LOW-VALUES TO DETLMAPO
*       MOVE SPACES    TO DMSGO
*       PERFORM BUILD-IT
WHEN DFHPF7
*
*
*       <PF7> is used to scroll backward if detail data
*       is displayed and we have previously scrolled

```

```

*           forward.
*
          IF MAIN-SHOWN
            PERFORM SEND-CONTROL
          END-IF
          IF CURRENT-PAGE = 1
            MOVE 3 TO MSGNO
            PERFORM SEND-CONTROL
          END-IF
          SUBTRACT 1 FROM CURRENT-PAGE
          MOVE LOW-VALUES TO DETLMAPO
          MOVE SPACES      TO DMSGO
          PERFORM BUILD-IT
        WHEN DFHCLEAR
*
*           <Clear> simply causes a screen refresh.
*
          IF MAIN-SHOWN
            PERFORM SEND-MAIN
          ELSE
            MOVE LAST-CURSOR TO VECTOR-IX
            PERFORM BUILD-IT
          END-IF
        WHEN OTHER
*
*           Any other key simply generates an error message.
*
          MOVE 1 TO MSGNO
          PERFORM SEND-CONTROL
        END-EVALUATE
      END-IF
    EXEC CICS RETURN
          TRANSID(EIBTRNID)
          COMMAREA(DFHCOMMAREA)
    END-EXEC
  .
SEND-MAIN.
*
* Indicate what we have done.
*
      SET MAIN-SHOWN TO TRUE
*
* Clear the output map area and set where we want the cursor.
*
      MOVE LOW-VALUES      TO MAINMAPO
      MOVE -1              TO ANCHLNL(LAST-CURSOR)
*
* Place the Job ID, Job Name, and CICS release in the output map.
*
      PERFORM ADDRESS-JSAB
      MOVE JSAB-JOBID     TO MJOBIDO

```

```

MOVE JSAB-JOBNAME TO MJOBMNO
MOVE WHAT-VERSION TO MRELO
*
* Fill in the Domain data.
*
PERFORM ADDRESS-KCB
PERFORM VARYING VECTOR-IX FROM 1 BY 1
UNTIL VECTOR-IX > MAX-NO-OF-DOMAINS
    SET ADDRESS OF DOMAIN-TABLE TO KCB-VECTOR(VECTOR-IX)
    IF DT-ID ALPHABETIC
        MOVE DT-ID TO DD-ID
        MOVE DT-IX TO DD-IX
        SET ADDR-PTR TO DT-ANCHOR-PTR
        CALL 'HEXMANIP' USING DT-ANCHOR-PTR,
            LENGTH-4,
            DD-ADDR
        MOVE OUT-LINE TO ANCHLNO(VECTOR-IX)
        IF WORK-PTR = ZERO
            MOVE DFHBMASK TO ANCHLNA(VECTOR-IX)
        END-IF
    ELSE
*
* NB The display fields are unprotected in order to allow
* the user to use the <Tab> key to position the cursor.
* So if the table entry is unused, we protect it.
*
        MOVE DFHBMASK TO ANCHLNA(VECTOR-IX)
    END-IF
END-PERFORM
*
* In order to allow the map to be used across all of the
* releases (and possibly for those in the future as well),
* the excess positions on the map are protected.
*
    IF MAX-NO-OF-DOMAINS < MAX-IN-MAP
        PERFORM VARYING VECTOR-IX FROM MAX-NO-OF-DOMAINS
            BY 1
        UNTIL VECTOR-IX > MAX-IN-MAP
            MOVE DFHBMASK TO ANCHLNA(VECTOR-IX)
        END-PERFORM
    END-IF
EXEC CICS SEND
    MAP('MAINMAP')
    MAPSET('ADDRMAP')
    ERASE
    CURSOR
    NOHANDLE
END-EXEC
.
SEND-DETAIL.
*
```



```

* Clear the output map area.
*
      MOVE LOW-VALUES      TO DETLMAPO
      MOVE SPACES          TO DMSGO
*
* Address the requested anchor block.
*
      PERFORM ADDRESS-KCB
      SET ADDRESS OF DOMAIN-TABLE TO KCB-VECTOR(VECTOR-IX)
      IF (DT-ID ALPHABETIC)
      AND (DT-ANCHOR-PTR-X NOT = ZERO)
*
*       For valid entries, display the requested data.
*
      MOVE DT-ID            TO DOMAIN,
                          DDOMIDO
      MOVE VECTOR-IX       TO LAST-CURSOR
*
*       Indicate what we have done.
*
      SET DETAIL-SHOWN     TO TRUE
      SET ADDRESS OF ANCHOR TO DT-ANCHOR-PTR
      SET START-ADDR TO DT-ANCHOR-PTR
      MOVE 1               TO CURRENT-PAGE
      EVALUATE DT-ID
*
*       Allow for the exceptions to the length convention.
*
      WHEN 'AP'
          MOVE A-AP-LTH TO TOTAL-LTH
      WHEN 'RX'
          MOVE 4096     TO TOTAL-LTH
      WHEN OTHER
          MOVE A-LENGTH TO TOTAL-LTH
      END-EVALUATE
      COMPUTE TOTAL-PAGES = ((TOTAL-LTH - 1) / 256) + 1
      PERFORM BUILD-IT
ELSE
      PERFORM SEND-CONTROL
END-IF
      .
BUILD-IT.

```

Editor's note: this article will be concluded next month.

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Listing the TMONCICS control file RLC definitions

Landmark's The Monitor for CICS/ESA is a widely used third-party monitoring software package for CICS installations. Amongst the facilities offered by The Monitor is the Resource Level Cancel (RLC) facility. This enables CICS performance staff to set limiting values for the critical processor resources consumed by a CICS transaction. If the transaction exceeds a defined limit, The Monitor attempts to cancel the transaction in order to protect the rest of the CICS system from the effects of the runaway transaction.

The resources are CPU, storage, and I/O operations. These are further broken down into task CPU time consumed and task elapsed time, task storage above and below 16MB, and I/O to DL/I, DB2, and to a user-defined database such as IDMS or ADABAS. In addition, The Monitor allows every transaction to have a limit on the number of individual EXEC CICS operations that it is permitted to execute.

To define these RLC values, Landmark offers a set of screens within the on-line portion of The Monitor, known as the Cross System Monitor or CSM, and they are stored in the control file. This is VSAM-based. There is, however, no way to list the definitions, other than by using the same set of screens.

To bypass having to use the CSM screens, and to quickly generate a listing of all the definitions currently in place, I have written the REXX program, TMCERLC, which formats the contents of the control file and writes a report as shown in Figure 1.

The report is mostly self-explanatory, the fields being exactly as defined in The Monitor documentation. The EIBFN lines show where transactions have been limited to the number of EXEC CICS functions they can execute, and show the internal CICS code, the number of that particular function that is allowed, and a description of the CICS function. The table of these descriptions does not contain all possible EXEC CICS functions, but it can be easily modified by adding function descriptions if desired. When the description 'Type not defined' appears, simply add an entry with the code shown and a description, either from The Monitor CSM RLC definition screens, or from the documentation in the relevant *CICS Diagnosis Reference* manual.

Jobname	Tran	CPU	STG	STG	Run	DLI	DB2	User DB	Excl	Excl	Excl
		sec	<16	>16	sec	I/O	I/O	I/O	Hi	Log	Wait
CICS0001	****	10	1024	1024	0	10000	10000	10000	N	N	N
CICS0001	C***	0	0	0	0	0	0	0	N	N	N
CICS0001	F1**	20	512	1024	0	5000	0	0	N	N	N
CICS0001	CSNE	0	0	0	0	0	0	0	Y	N	Y
CICS0001	CSSY	0	0	0	0	0	0	0	Y	N	Y
CICS0001	CSTP	0	0	0	0	0	0	0	Y	N	Y
CICS0001	TM**	0	0	0	0	0	0	0	Y	N	Y
CICS0002	****	99	0	0	0	0	0	0	N	N	N
CICS0002	C***	0	0	0	0	0	0	0	N	N	N
CICS0002	F2**	20	256	512	0	5000	0	0	N	N	N
EIBFN	0602		100		Read						
EIBFN	0604		100		Write						
EIBFN	0606		50		Rewrite						
EIBFN	0608		20		Delete						
CICS0002	CSNE	0	0	0	0	0	0	0	Y	N	Y
CICS0002	CSSY	0	0	0	0	0	0	0	Y	N	Y
CICS0002	CSTP	0	0	0	0	0	0	0	Y	N	Y
CICS0002	TEST	0	0	0	0	0	0	0	N	N	N
EIBFN	0602		10		Read						
EIBFN	0604		10		Write						
CICS0002	TM**	0	0	0	0	0	0	0	Y	N	Y

Figure 1: TMCERLC sample output

TMCERLC REXX SOURCE

```

/*----- REXX----- */
/* Function   : Process TMON for CICS/ESA V2.0 VTCECNTL to extract */
/*           : resource limit cancel values.                       */
/*----- */
numeric digits 21
call init_eib
say ''
say 'Jobname Tran CPU STG STG Run DLI ',
' DB2 User DB Excl Excl Excl'
say '
sec <16 >16 sec I/O ',
' I/O I/O Hi Log Wait'
done = 'n'
do while done = 'n'
"execio 1 diskr tmndat"
if rc = 0 then
do

```

```

    parse pull tmnrec
    call proc_rec
    end
else
    done = 'y'
end
exit 0
/*-----*/
/* Process a record */
/*-----*/
proc_rec:
type = substr(tmnrec,1,1)
if type = 'K' then
    do
    jobn = substr(tmnrec,2,8)
    if jobn ≠ sjob then
        say ''
    sjob = jobn
    tran = substr(tmnrec,18,4)
    select
        when tran = '00000000'x then return
        when trlc = '7FFFFFFF'x then return
    otherwise
        do
            rcpu = substr(tmnrec,49,4)
            rsa = substr(tmnrec,53,4)
            rsb = substr(tmnrec,57,4)
            rrun = substr(tmnrec,61,4)
            rdli = substr(tmnrec,65,4)
            rdb2 = substr(tmnrec,69,4)
            rusr = substr(tmnrec,73,4)
            xhi = substr(tmnrec,77,1)
            xlog = substr(tmnrec,78,1)
            xwt = substr(tmnrec,79,1)
            if rcpu ≠ '7FFFFFFF'x then
                rcpu = c2d(rcpu) * 64 / 1000000
            else
                rcpu = 0
            if rsb ≠ '7FFFFFFF'x then
                rsb = c2d(rsb) / 1024
            else
                rsb = 0
            if rsa ≠ '7FFFFFFF'x then
                rsa = c2d(rsa) / 1024
            else
                rsa = 0
            if rrun ≠ '7FFFFFFF'x then
                rrun = c2d(rrun) * 64 / 1000000
            else
                rrun = 0
            if rdli ≠ '7FFFFFFF'x then

```

```

        rdli   = c2d(rdli)
    else
        rdli   = 0
    if rdb2 = '7FFFFFFF'x then
        rdb2   = c2d(rdb2)
    else
        rdb2   = 0
    if rusr = '7FFFFFFF'x then
        rusr   = c2d(rusr)
    else
        rusr   = 0
    say jobn tran format(rcpu,4,0) format(rsb,6) format(rsa,6),
        format(rrun,4,0) format(rdli,8) format(rdb2,8),
        format(rusr,8) ' ' xhi ' ' xlog ' ' xwt
    j = 80
    do i = 1 to 32
        eibrec = substr(tmnrec,j,6)
        eib = substr(eibrec,1,2)
        select
            when eib = '0000'x then nop
            when eib = 'FFFF'x then nop
            otherwise do
                k = c2x(eib)
                say ' EIBFN' c2x(eib) ' ',
                    format(c2d(substr(eibrec,3,4)),6) ' ' eibtyp.k
            end
        end
        j = j + 6
    end
end
end
end
return
/*-----*/
/* Initialize EIB descriptions - not an exhaustive list, add */
/* additional descriptions as required, see CICS Diagnosis Reference */
/* for a list of command descriptions. */
/*-----*/
init_eib:
eibtyp.      = 'Type not defined'
eibtyp.0202  = 'Address'
eibtyp.0204  = 'Handle condition'
eibtyp.0206  = 'Handle aid'
eibtyp.0208  = 'Assign'
eibtyp.020A  = 'Ignore condition'
eibtyp.020C  = 'Push'
eibtyp.020E  = 'Pop'
eibtyp.0210  = 'Address set'
eibtyp.0402  = 'Receive'
eibtyp.0404  = 'Send'
eibtyp.0406  = 'Converse'

```

eibtyp.040C = 'Wait terminal'
eibtyp.0410 = 'Wait signal'
eibtyp.0420 = 'Allocate'
eibtyp.0422 = 'Free'
eibtyp.042C = 'Wait convid'
eibtyp.042E = 'Extract process'
eibtyp.0430 = 'Issue abend'
eibtyp.0432 = 'Connect process'
eibtyp.0602 = 'Read'
eibtyp.0604 = 'Write'
eibtyp.0606 = 'Rewrite'
eibtyp.0608 = 'Delete'
eibtyp.060A = 'Unlock'
eibtyp.060C = 'Startbr'
eibtyp.060E = 'Readnext'
eibtyp.0610 = 'Readprev'
eibtyp.0612 = 'Endbr'
eibtyp.0614 = 'Resetbr'
eibtyp.0802 = 'Writeq TD'
eibtyp.0804 = 'Readq TD'
eibtyp.0806 = 'Deleteq TD'
eibtyp.0A02 = 'Writeq TS'
eibtyp.0A04 = 'Readq TS'
eibtyp.0A06 = 'Deleteq TS'
eibtyp.0C02 = 'Getmain'
eibtyp.0C04 = 'Freemain'
eibtyp.0E02 = 'Link'
eibtyp.0E04 = 'Xctl'
eibtyp.0E06 = 'Load'
eibtyp.0E08 = 'Return'
eibtyp.0E0A = 'Release'
eibtyp.0E0C = 'Abend'
eibtyp.0E0E = 'Handle abend'
eibtyp.1002 = 'Asktime'
eibtyp.1004 = 'Delay'
eibtyp.1006 = 'Post'
eibtyp.1008 = 'Start'
eibtyp.100A = 'Retrieve'
eibtyp.100C = 'Cancel'
eibtyp.1202 = 'Wait event'
eibtyp.1204 = 'Enqueue'
eibtyp.1206 = 'Dequeue'
eibtyp.1208 = 'Suspend'
eibtyp.1402 = 'Write journalnum'
eibtyp.1404 = 'Wait journalnum'
eibtyp.1602 = 'Syncpoint'
eibtyp.1802 = 'Receive map'
eibtyp.1804 = 'Send map'
eibtyp.1806 = 'Send text'
eibtyp.1808 = 'Send page'
eibtyp.180A = 'Purge message'

```

eibtyp.180C = 'Route'
eibtyp.180E = 'Receive partn'
eibtyp.1810 = 'Send partnset'
eibtyp.1812 = 'Send control'
eibtyp.1C02 = 'Dump'
eibtyp.4A02 = 'Asktime abstime'
eibtyp.4A04 = 'Format time'
eibtyp.5602 = 'Spool open'
eibtyp.5604 = 'Spool read'
eibtyp.5606 = 'Spool write'
eibtyp.5610 = 'Spool close'
eibtyp.5E22 = 'Wait external'
eibtyp.5E32 = 'Wait CICS'
eibtyp.6402 = 'Perform security'
eibtyp.6C02 = 'Write operator'
eibtyp.6C12 = 'Issue DFHWT0'
eibtyp.7402 = 'Signon'
eibtyp.7404 = 'Signoff'
eibtyp.7406 = 'Verify password'
eibtyp.7408 = 'Change password'
eibtyp.7602 = 'Perform shutdown'
eibtyp.7E02 = 'Dump transaction'
eibtyp.7E04 = 'Dump system'
return

```

JCL TO RUN TMCERLC

Because REXX cannot read a VSAM file directly, it is necessary to reproduce the control file data into a sequential dataset before running TMCERLC. The following JCL achieves this:

```

//TMCERLC JOB
//*
//EXTRCNTL EXEC PGM=IDCAMS
//SYSPRINT DD SYSOUT=*
//IN DD DSN=TMCE.V20.VTCECNTL,DISP=SHR
//OUT DD DSN=&&CNTL,DISP=(,PASS),UNIT=SYSDA,SPACE=(CYL,(1,1)),
// DCB=(RECFM=VB,LRECL=27994,BLKSIZE=27998)
//SYSIN DD *
    REPRO INFILE(IN) OUTFILE(OUT) FROMKEY(K) TOKEY(K) /* RLC RECS */
//*
//TMCERLC EXEC PGM=IRXJCL,PARM='TMCERLC'
//SYSEXEC DD DSN=SYS1.SYSEXEC,DISP=SHR
//TMNDAT DD DSN=&&CNTL,DISP=(OLD,DELETE)
//SYSTSIN DD DUMMY
//SYSTSPRT DD SYSOUT=*
//*

```

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CICS news

InfoSpinner has announced CICS support for its ForeSite Application Server. ForeSite Transaction Server, a new System Integration Module (SIM), has been released as part of a suite of application integration products.

Transaction Server SIM is an add-on to ForeSite Application Server 3.0, and can be installed on ForeSite's PageServers to run concurrently with CICS software. It allows ForeSite to interface with transaction server software and update large databases as transactions occur on a Web site. It also provides OLTP management for applications on both IBM and non-IBM platforms.

For further information contact:
InfoSpinner, 1601 North Glenville Drive,
Suite 108, Richardson, TX 75081, USA.
Tel: (972) 479 0135.
URL: <http://www.infospinner.com>.

* * *

CICS users can benefit from Software AG's announcement of support for Java applications on OS/390 mainframes for its Bolero Application Factory.

Although Bolero creates Java Byte Code that can theoretically run on all platforms that have a Java Virtual Machine, Java implementations are different on mainframes, NT, and Unix. The OS/390 version of Bolero supports mainframes specifically, and allows Bolero applications to be used in the CICS Open Transaction

Environment and to store persistent objects in DB2 databases.

For further information contact:
Software AG (UK), Charter Court, 74/78
Victoria Street, St Albans, Herts, AL1 3XH,
UK.
Tel: (01727) 844 455.
Software AG of America, 11190 Sunrise
Valley Drive, Reston, VA 22091, USA.
Tel: (703) 860 5050.
URL: <http://www.software-ag.com>.

* * *

IBM has announced CICS Transaction Server for VSE/ESA Release 1. The product includes a new CICS Server, CICS Universal Clients, CICS Transaction Gateway, REXX for CICS, CICS Distributed Data Management, and CICS/VSE Version 2.3, including Report Controller Facility (RCF) for both levels of CICS server. Future additions will include a CICS Web Interface (CWI), and the CICS 3270 Bridge.

Release 1 includes a new facility for viewing CSMT messages in the coexistent CICS/VSE Version 2.3. Translated versions of CICS Transaction Server for VSE/ESA messages and RCF panels are provided in Japanese, simplified Chinese, and German.

For further information contact your local IBM representative.

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