



# 204

# MVS

*September 2003*

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# update

# ***MVS Update***

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## DFHSM error report from daily SMF extract

After publishing the article entitled *A simple DFHSM report writer (MVS Update, issue 191, August 2002)* I received several requests to make the code more applicable. The key point of these requests was that the code was written in the SAS language, which is not available at all mainframe sites.

What was really needed was a simple, flexible, and yet fast reporting tool that would quickly identify work that has not been completed by DFHSM.

In order to provide a simple and effective way to gather the data needed to monitor, analyse, and correct what went wrong during DFHSM activity, an easy-to-use reporting procedure has been created.

### CODE

The code is a two-part stream. In the first part, COPYSMF, selected DFHSM SMF records of the FSR type are copied from the SMF dataset to a file that can be used as a base for the archived records. By default, DFHSM writes SMF id 241 for functional statistics records (FSR and WWFSR). A detailed description of the layout of FSR records and its fields can be obtained from the DFSMS manual *DFSMSHsm Implementation and Customization Guide*.

In the second part of the code, HSMERR, the captured records are formatted and two reports are produced.

The first report (DFHSM Function Error Statistics report) shows the summary of DFHSM errors along with a breakdown of this summary data by DFHSM function (BACKUP, SPILL, Primary to Level 1, etc), date, and return/reason codes, as well as the totals of the numbers of each error. The second report is a detailed dataset-level activity, reporting *all* activity that failed.

Note: a non-zero return code indicates that the function did not

complete successfully. The meanings of the return codes are documented in *DFHSM Messages*, SH35-0094. Use this field along with the function field to find out what function was running. Then, see message ARC0734I to determine what the return code means.

```
//DEL EXEC PGM=IDCAMS
//SYSPRINT DD SYSOUT=X
//SYSIN DD *
    DELETE hlq.SMF241
    SET MAXCC=0
/*
//COPYSMF EXEC PGM=IFASMFDP, REGION=0M
//INDA1 DD DSN=sysid.SMFDUMPW, DISP=SHR * weekly/daily SMF file
//OUTDA DD DSN=hlq.SMF241, DISP=(NEW,CATLG), UNIT=SYSDA,
//      SPACE=(CYL,(x,y),RLSE),
//      DCB=(sysid.SMFDUMPW)
//SYSPRINT DD SYSOUT=X
//SYSIN DD *
    INDD(INDA1,OPTIONS(DUMP))
    OUTDD(OUTDA,TYPE(241))
/*
//HSMERR EXEC PGM=ICETOOL, REGION=0M
//TOOLMSG DD SYSOUT=X
//DFSMSG DD SYSOUT=X
//FUNCERR DD SYSOUT=X
//DETAIL DD SYSOUT=X
//SMF DD DSN=hlq.SMF241, DISP=SHR
//SHOW DD DSN=*&TEMPD, SPACE=(CYL,(15,15)), UNIT=SYSDA
//SMFC DD DSN=*&TEMPV, SPACE=(CYL,(15,15)), UNIT=SYSDA
//TOOLIN DD *
    COPY FROM(SMF) TO(SMFC) USING(SMFI)
*
* Print a summary report showing total number of errors
*
OCCUR FROM(SHOW) LIST(FUNCERR) -
TITLE('DFHSM Function Error Statistic Report') DATE TIME -
HEADER('Function') ON(1,13,CH) -
HEADER('Total number of errors') ON(VALCNT) -
BLANK
*
* Print a summary report showing number of errors by RC/RS
*
OCCUR FROM(SHOW) LIST(FUNCERR) -
TITLE('DFHSM Function Error Report by RC & RS') DATE TIME -
HEADER('Function') ON(1,13,CH) -
HEADER('Ret. Code') ON(26,4,BI) -
HEADER('Reas. Code') ON(30,4,BI) -
```

```

    HEADER('Total number of errors')  ON(VALCNT)      -
    BLANK
*
* Print a summary report showing number of function errors by date
*
OCCUR FROM(SHOW) LIST(FUNCERR)                      -
TITLE(' DFHSM Function Error Report by Function') DATE TIME -
  HEADER(' Function' )                ON(1, 13, CH)   -
  HEADER(' Date' )                    ON(18, 4, DT1, E' 9999/99/99' ) -
  HEADER(' Number of errors' )        ON(VALCNT)       -
  BLANK
*
* Print a summary report showing number of function errors by date
*
OCCUR FROM(SHOW) LIST(FUNCERR)                      -
TITLE(' DFHSM Function Error Report by Date') DATE TIME -
  HEADER(' Date' )                    ON(18, 4, DT1, E' 9999/99/99' ) -
  HEADER(' Function' )                ON(1, 13, CH)   -
  HEADER(' Number of errors' )        ON(VALCNT)       -
  BLANK
*
* Print a detailed report of function errors by date
*
DISPLAY FROM(SHOW) LIST(DETAIL)                    -
TITLE(' DFHSM Function Error Report - Detailed') DATE TIME -
  HEADER(' Si d' )                    ON(14, 4, CH)   -
  HEADER(' Time' )                    ON(22, 4, TM1, E' 99: 99: 99' ) -
  HEADER(' Function' )                ON(1, 13, CH)   -
  HEADER(' RC' )                      ON(26, 4, BI)   -
  HEADER(' REAS' )                    ON(30, 4, BI)   -
  HEADER(' Last Ref. ' )              ON(34, 4, DT1, E' 9999/99/99' ) -
  HEADER(' Last Moved' )            ON(38, 4, DT1, E' 9999/99/99' ) -
  HEADER(' DS Age' )                  ON(62, 2, BI)   -
  HEADER(' Request by' )              ON(42, 8, CH)   -
  HEADER(' From Vol. ' )              ON(56, 6, CH)   -
  HEADER(' To Vol. ' )                ON(50, 6, CH)   -
  HEADER(' Dsorg ' )                  ON(116, 4, CH)  -
  HEADER(' Data Set Name' )          ON(64, 44, CH)   -
  HEADER(' Mcl ass' )                ON(108, 8, CH)   -
  BREAK(18, 4, DT1, E' 9999/99/99' ) -
  BTITLE(' : Dai ly Error Report' ) -
  BLANK
/*
//SMFICNTL DD *
  OPTION COPY, VLSHRT
  OUTFIL FNAMES=SHOW, CONVERT,          * Build a new output record
  INCLUDE=(6, 1, BI, EQ, X' F1' , AND, 109, 4, BI, GT, X' 0000' ),
  OUTREC=(43, 1, CHANGE=(13,
    X' 01' , C' Migrate I0/I1' , * primary to level 1 migration

```

```

X' 02' ,C' Migrate 11/12' , * level 1 to level 2 migration
X' 03' ,C' Migrate 10/12' , * primary to level 2 migration
X' 04' ,C' Recall 11/10 ' , * recall from level 1 to primary
X' 05' ,C' Recall 12/10 ' , * recall from level 2 to primary
X' 06' ,C' Delete mig. ' , * delete migrated dataset
X' 07' ,C' Daily backup ' , * daily backup
X' 08' ,C' Spill backup ' , * spill backup
X' 09' ,C' Recovery ' , * recovery from backup
X' 0A' ,C' Recycle bk. ' , * recycle backup volume
X' 0B' ,C' Del by age ' , * dataset deletion by age
X' 0C' ,C' Recycle ml2 ' , * recycle migration volume
X' 0D' ,C' Volume dump ' , * full volume dump
X' 0E' ,C' Restore ' , * volume or dataset restore
X' 0F' ,C' Abackup ' , * abackup function
X' 10' ,C' Arecover ' , * arecover function
X' 11' ,C' Expire ' , * expire primary or mig. dataset
X' 12' ,C' Partrel ' , * partial release
X' 13' ,C' Expired bk. ' , * expire incremental backup ver.
X' 14' ,C' Delete bk. '), * delete incremental backup ver.
15, 4, * System identification
11, 4, * Date of function being performed
7, 4, * Time
109, 4, * Return code from mwe
113, 4, * Reason code from mwe
157, 4, * Dataset last referenced date
161, 4, * Dataset last moved date
35, 8, * Id of the user requesting the function
89, 6, * Receiving volume
99, 6, * From volume
185, 2, * Dataset age
45, 44, * Dataset name
215, 8, * Management class
177, 2, CHANGE=(4, * Dsorg (as char field)
X' 8000' ,C' IS ' , * change to indexed seq
X' 4000' ,C' PS ' , * change to physical seq
X' 2000' ,C' DA ' , * change to direct
X' 0008' ,C' VSAM' , * change to vsam
X' 0200' ,C' PDS ' , * change to partitioned
X' 4100' ,C' PSU ' , * change to ps unmoveable
X' 0080' ,C' GS '), * change to graphic
NOMATCH=(C' UNDF' )) * change to undefined

```

/\*

It is to be hoped that DASD/Storage Management personnel now have the ability to easily identify and analyse any action taken by DFHSM that failed.

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## Archival data warehousing designs

The underlying objectives of data warehouse designs can be summarized as:

- Storing a large volume of data.
- The fastest and cheapest retrieval of data
- OLTP interface for user-friendly querying.
- Controlling system load.

Business applications running on legacy systems often need such a design to query archival/back-up data, without restrictions on size. *Ad hoc* reporting techniques available on MVS systems (see *Handling ad hoc reports querying large volumes of data – a case study, MVS Update*, issue 201, June 2003) can be used to build a primitive data warehouse that meets these underlying principles.

### DATA ARCHIVAL DESIGN

Generation data groups can be used for taking back-ups daily. Because GDGs support a maximum of only 255 versions, we must collect data differently, based on a volume of data that we can browse sequentially. Below are two ways to do this.

- Monthly, appending data daily to the current version. We use a file layout with the date of archival as the file's key, which helps to search data by date. This is when a month's volume of data is suitable for sequentially browsing.
- Separate GDGs for each month with 28/29/30/31 versions based on month and year. We need two-tier JCL with INCLUDE groups (see previous article). First the JCL runs a program with built-in logic to read the date of the data archival/system date and converts it to a GDG version number. Thus the correct generation is located to store the back-up data. This information is passed to the second JCL

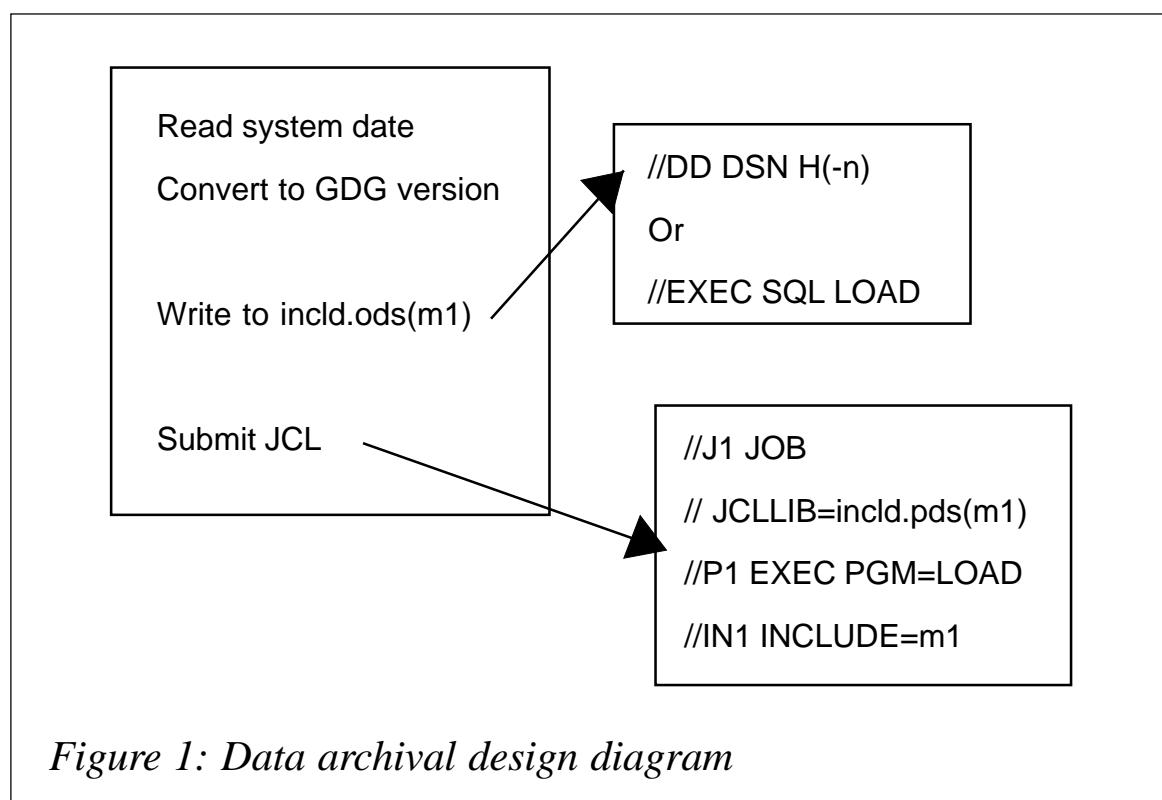
with the INCLUDE feature. The second JCL uses this INCLUDE member and stores the data. This pair runs daily to collect archival data. This JCL pair helps when a month's volume is too great for sequential browsing during data retrieval. We can have as a file key the back-up hour too, for those applications where back-ups are taken hourly.

Thus the design of the data storage layer aims at zero maintenance by storing archival data in GDGs.

It directly influences the speed of data retrieval by:

- Using two-tier JCL with INCLUDE groups.
- Choosing a file key of date/hour.

Figure 1 shows a data archival design diagram.



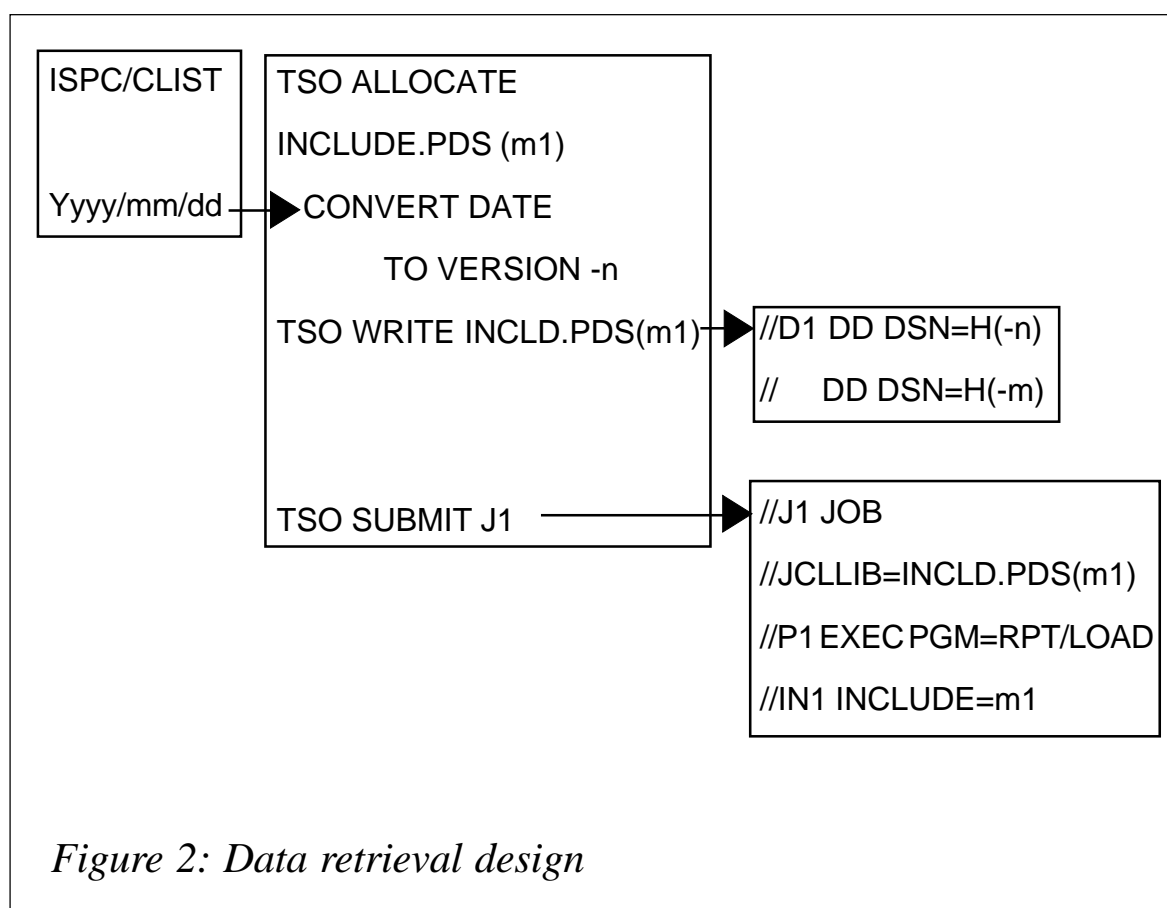
## DATA RETRIEVAL DESIGN

As discussed above, we can embed in an ISPF/CLIST front end, built-in date-version conversion logic, prepare INCLUDE



members, and trigger the JCL to run using the INCLUDE dataset group. Such an application runs queries single-threaded. Using a design with two-tier JCL, using job scheduler switches (see previous article), means that JCL waits in the MVS input queue until the preceding job pair completes. This multi-threads keying-in requests. Batch query jobs still run single-threaded. By preparing separate INCLUDE members instead, a common member query can be multi-threaded.

Queries can be made much more flexible by preparing INCLUDE members with DB2 load control statements instead of JCL DD statements. Thus data warehousing designs can be created with much more advanced front-end designs on top of layers of the fastest and cheapest archival/retrieval designs.



Data retrieval design is illustrated in Figure 2.

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## Locating strings in files

For most purposes, the Searchfor utility is an acceptable way to search for strings within files. However, it lacks support for VSAM files and a way to limit the scope of the search. For example, I need to search for a name that should exist in a *name* field, but eventually might also exist in other fields, like *address*, etc. If I know the position of the field I want to search within each record, then I may wish to restrict my search to that area. Or I may choose to search only a limited set of records, and not the entire file.

To satisfy these needs, I wrote a search program that accepts both sequential and VSAM files, allows the search zone to be limited to a range of columns within each record or to a range of records, permits the search string to be specified in hexadecimal, and also permits case to be ignored when performing comparisons (this applies only to strings specified as characters and to standard a-z characters, but you can change the translation table within the program, if you like).

The application consists of an Assembler program and a front-end formed by a REXX EXEC and an ISPF panel. It looks like this:

```
+----- Locate a string in a file -----+
|
| File.....: AARCF3.TEST.BA300
|
| Search for: x'03f5f4c1a3
|   (Enter text string or begin with X' for hexadecimal)
|
| Ignore case - only valid for text (Y,N): Y
|
| Search columns           Search records
| First column.   12       First record.   10000
| Last column.   25       Last record.    15000
|
| Enter - execute                               PF3/15 cancel
|
+-----+

```

In this example, we search for a hexadecimal string (since it begins with x', the ending quote is optional and can be omitted) with column boundaries 12 through to 25 and record boundaries 10,000 through to 15,000.

Since the string is hexadecimal, the *ignore case* field has no meaning. The program launches a job, and the result of the search can be found in sysprint: it states the record numbers where a match was found, the total number of record matches, and the total number of records searched:

```
String found in record number 00000012034
String found in record number 00000012035
String found in record number 00000012036
Number of records where string found: 00000000003
Number of records searched. . . . : 00000005001
```

## LOCATE ASSEMBLER PROGRAM

```
*=====*
* LOCATE - Locates and counts the number of record matches of a      *
*           string in a file. The file can be sequential or VSAM.    *
*                                                                 *
* Format of the parameter received:                                  *
* Offset Name and meaning                                           *
* 0   Col1 initial position within records                          *
* 4   Col2 final position within records                            *
* 8   Rec1 initial search record                                    *
* 16  Rec2 final search record                                      *
* 24  Flag X-Hexa string Y-ignore case otherwise respect case     *
* 25  String to search.                                           *
*                                                                 *
* DDnames: Infile, Sysprint                                        *
*                                                                 *
* This program reads an input file and searches for a string in each *
* record. The search can be limited to a column range within each  *
* record (Col1 to Col2) or to a set of records (Rec1 to Rec2).     *
* If flag = 'X', string is represented in hexadecimal.            *
* If flag = 'Y', string and searched areas are uppercased before  *
* comparison takes place (case is ignored). This applies only to  *
* a-z standard EBCDIC characters.                                  *
*=====*
&PROGRAM SETC 'LOCATE'
&PROGRAM AMODE 31
&PROGRAM RMODE 24
&PROGRAM CSECT
```

```

SAVE (14, 12)
LR R12, R15
USI NG &PROGRAM, R12
ST R13, SAVEA+4
LA R11, SAVEA
ST R11, 8(R13)
LR R13, R11
B GETPARMS
DC CL16' &PROGRAM 1. 1'
DC CL8' &SYSDATE'

```

\*

\*=====\*

\* Check and validate parameters \*

\*=====\*

\*

```

GETPARMS DS 0H
LR R2, R1 Copy parm pointer to R2.
L R2, 0(0, R2) Load parm address
LH R3, 0(R2) Load parm length in R3
OPEN (SYSPRINT, OUTPUT) Open sysprint (for error msgs)
LTR R3, R3 Any parm entered?
BZ EXIT1 No, error

```

\*

```

LR R6, R2
AR R6, R3 R6: point after end of parms
LA R6, 2(0, R6) Skip 2 bytes of parmlength
LA R2, 2(0, R2)
XR R9, R9 Clear length counter

```

\*

```

CONVERT EQU *
LA R9, 3 4 byte length parms
EX R9, EXPACK Execute pack
CVB R7, PARMPACK Convert to binary into R7
S R7, =F' 1' Turn limit to offset
ST R7, PARM1 And store it

```

\*

```

LA R2, 4(0, R2) Inc parm pointer
EX R9, EXPACK Execute pack
CVB R7, PARMPACK Convert to binary into R7
ST R7, PARM2 And store it

```

\*

```

LA R2, 4(0, R2) Inc parm pointer
LA R9, 7 8 byte length parms
EX R9, EXPACK Execute pack
CVB R7, PARMPACK Convert to binary into R7
ST R7, PARM3 And store it

```

\*

```

LA R2, 8(0, R2) Inc parm pointer
EX R9, EXPACK Execute pack

```

	CVB	R7, PARMPACK	Convert to binary into R7
	ST	R7, PARM4	And store it
*			
	LA	R2, 8(Ø, R2)	Inc parm pointer
	MVC	FLAG, Ø(R2)	Store flag parameter
*			
	LA	R2, 1(Ø, R2)	Inc parm pointer to string
	SR	R6, R2	Length of string to search
	SH	R6, =H' 1'	Length of string to search
	EX	R6, EXMOVE	Move to string
	STH	R6, STRINGL	Keep string length ( - 1)
	CLI	FLAG, C' X'	Hexadecimal string?
	BNE	VALIDPR	No, jump ahead
*			
	MVC	STR1, STRING1	Convert string to real hexadecimal
	MVC	STR2, STRING2	characters in three 12 byte parts
	MVC	STR3, STRING3	
	NC	STR1, XAND	
	TR	STR1, XTRN	
	NC	STR2, XAND	
	TR	STR2, XTRN	
	NC	STR3, XAND	
	TR	STR3, XTRN	
	PACK	STRP1, STR1(13)	
	PACK	STRP2, STR2(13)	
	PACK	STRP3, STR3(13)	
	MVC	STRING(6), STRP1	
	MVC	STRING+6(6), STRP2	
	MVC	STRING+12(6), STRP3	
*			
	LH	R6, STRINGL	Get original length (-1)
	LA	R6, 1(Ø, R6)	Add 1
	SRL	R6, 1	Divide length by two
	SH	R6, =H' 1'	Ready for comparisons (-1)
	STH	R6, STRINGL	Store it
*			
VALIDPR	EQU	*	Validate parameters
	CLC	PARM1, PARM2	
	BH	ERRMSG1	
	CLC	PARM3, PARM4	
	BH	ERRMSG2	
	L	R7, PARM2	
	S	R7, PARM1	
	CR	R6, R7	
	BH	ERRMSG3	
*			
	CLI	FLAG, C' Y'	Ignore case specified?
	BNE	OPENACB1	No, jump ahead
	LH	R1Ø, STRINGL	Load string length (-1)

```

        LA      R10, 1(0, R10)          Reset correct length
        LA      R5, STRING
        EX      R5, EXTRAN             Execute uppercase translation
*
*=====*
* Check whether file is VSAM or sequential. If VSAM, check for ESDS *
*=====*
*
OPENACB1 EQU *
        OPEN   INFILEA                Open ACB for VSAM input file
        LTR    R15, R15                If error, go open the file as
        BNZ    OPENDCB1                sequential.
        TESTCB ACB=INFILEA,           File is VSAM, check for ESDS      X
              ATRB=ESDS
        BNE    READFILE
*
ESDSFIL1 EQU *
        MODCB  RPL=INFILER,           Set RPL for ESDS                  X
              OPTCD=ADR
        B      READFILE
*
OPENDCB1 EQU *
        OPEN   (INFILED, INPUT)       Open sequential file
        LTR    R15, R15
        BNZ    ERRMSG4
        MVI    FILETYP1, C' S'        Set flag sequential
        LA     R2, INFILED             R2: IHADCB of input file
        USING  IHADCB, R2
*
*=====*
* Read loops (VSAM loop or sequential loop) and compare subroutine *
*=====*
*
READFILE EQU *
        XR     R7, R7                  Record counter
        L      R8, PARM3               First record
        L      R9, PARM4               Last record
        CLI    FILETYP1, C' V'        VSAM file?
        BNE    LOOPSEQ                No, go to sequential
*
LOOPVSA  EQU *
        LA     R7, 1(0, R7)            Increment record counter
        GET    RPL=INFILER            Read VSAM file
        LTR    R15, R15                End of file?
        BNZ    EXIT0
        CR     R7, R8                  First record attained?
        BL     LOOPVSA                 No, read again
        CR     R7, R9                  Last record attained?
        BH     EXIT0                   Yes, exit

```

*				
	L	R4, VAREA1	Get address of data in R4.	
	SHOWCB	RPL=INFILER, AREA=LRECL1, LENGTH=4, FIELDS=RECLEN		X X X
	L	R3, LRECL1	Get record length in R3	
	BAL	R10, COMPARE	Call compare	
	B	LOOPVSA		
*				
LOOPSEQ	EQU	*	Sequential loop	
	LA	R7, 1(0, R7)	Increment record counter	
	GET	INFILED	Read sequential	
	LR	R4, R1	R4: address of record	
	CR	R7, R8	First record attained?	
	BL	LOOPSEQ	No, read again	
	CR	R7, R9	Last record attained?	
	BH	EXIT0	Yes, exit	
	LH	R3, DCBLRECL	Load R3 with record length.	
	BAL	R10, COMPARE	Call compare	
	B	LOOPSEQ		
*				
COMPARE	EQU	*	Compare subroutine	
	L	R5, PARM1	Left limit (offset)	
	L	R6, PARM2	Right limit	
	CR	R3, R6	Record smaller than last position?	
	BNL	COMPARE0	No, continue.	
	LR	R6, R3	Yes, switch limit to record limit	
*				
COMPARE0	EQU	*		
	SR	R6, R5	R6: length of searchable area	
	AR	R5, R4	R6: position to compare	
	CLI	FLAG, C' Y'	Ignore case specified?	
	BNE	COMPARE9	No, jump ahead	
	EX	R6, EXTRAN	Execute uppercase translation	
*				
COMPARE9	EQU	*		
	SH	R6, STRINGL	R6: number of searches in the line	
	C	R6, =F' 0'	If R6<0, string is greater than search area, so skip this record	
	BNH	COMPARE3		
*				
COMPARE1	EQU	*		
	LH	R11, STRINGL	length of search string	
	EX	R11, EXCOMPAR	Execute compare	
	BNE	COMPARE2	If strings not equal, exit	
	LR	R0, R7		
	BAL	R11, UNPACK		
	MVC	MSGFND2, OUT10		
	LR	R0, R9		
	PUT	SYSPRINT, MSGFND	String found, send message	

```

        L      R11, TOTFOUND      Inc rec found counter
        LA     R11, 1(Ø, R11)
        ST     R11, TOTFOUND
        B      COMPARE3          And return
*
COMPARE2 EQU *
        LA     R5, 1(Ø, R5)      Increment compare position
        BCT   R6, COMPARE1      Loop to next
*
COMPARE3 EQU *
        BR    R1Ø              return
*
*=====*
* Send final messages, close files, and exit *
*=====*
*
EXITØ   EQU *
        L      RØ, TOTFOUND
        BAL   R11, UNPACK
        MVC   MSGTOT2, OUT1Ø
        PUT   SYSPRINT, MSGTOT
        SR    R7, R8
        LR    RØ, R7
        BAL   R11, UNPACK
        MVC   MSGFIM2, OUT1Ø
        PUT   SYSPRINT, MSGFIM
*
EXIT1   EQU *
        CLOSE INFILED
        CLOSE INFILEA
        CLOSE SYSPRINT
        L     R13, SAVEA+4
        LM   R14, R12, 12(R13)
        XR   R15, R15
        BR   R14
*
*=====*
* Other subroutines, execute instructions and work areas *
*=====*
*
EXCOMPAR EQU *
        CLC   Ø(Ø, R5), STRING
*
EXMOVE   EQU *
        MVC   STRING, Ø(R2)
*
EXPACK   EQU *
        PACK  PARMPACK, Ø(Ø, R2)
*
EXTRAN   EQU *

```



```

      TR      Ø(Ø, R5), TRTAB
*
UNPACK  EQU   *
        CVD  RØ, REGDECIM
        UNPK OUT12, REGDECIM
        BR   R11
*
ERRMSG1 EQU   *
        PUT  SYSPRINT, MSG1
        B    EXITØ
*
ERRMSG2 EQU   *
        PUT  SYSPRINT, MSG2
        B    EXITØ
*
ERRMSG3 EQU   *
        PUT  SYSPRINT, MSG3
        B    EXITØ
*
ERRMSG4 EQU   *
        PUT  SYSPRINT, MSG4
        B    EXITØ
*
MSG1    DC    CL8Ø' Error: Parm2 smaller than parm1'
MSG2    DC    CL8Ø' Error: Parm3 smaller than parm4'
MSG3    DC    CL8Ø' Error: String length smaller than search interval'
MSG4    DC    CL8Ø' Error opening input file'
MSG5    DC    CL8Ø' Record smaller than compare position'
MSGFND  DS    ØCL8Ø
MSGFND1 DC    C' String found in record number  '
MSGFND2 DC    CL5Ø'  '
MSGTOT  DS    ØCL8Ø
MSGTOT1 DC    C' Number of records where string found:  '
MSGTOT2 DC    CL5Ø'  '
MSGFIM  DS    ØCL8Ø
MSGFIM1 DC    C' Number of records searched. . . . . :  '
MSGFIM2 DC    CL5Ø'  '
*
STRINGL DS    H
        DS    CL2
STRING  DS    ØCL36
STRING1 DS    CL12
STRING2 DS    CL12
STRING3 DS    CL12
        DS    CL4
STR1    DS    CL12
        DC    C'  '
STR2    DS    CL12
        DC    C'  '

```



```

FLAG      DS      C
PARMPACK  DS      D
PARM1     DS      F
PARM2     DS      F
PARM3     DS      F
PARM4     DS      F
          DS      ØD
REGDECIM  DS      CL9
          DS      ØF
OUT12     DS      ØCL12
OUT1Ø     DS      CL1Ø
          DS      CL2
*
          LTORG
          DCBD  DSORG=PS
          YREGS
          END

```

## LOCATE REXX EXEC

```

/* REXX MVS *=====*/
/* LOCATE - Locates a string within a file. */
/*          Optional argument: file to search. */
/*          */
/* This application consists of this EXEC, LOCATE ISPF panel, */
/* and LOCATE Assembler program. The load module should reside */
/* in the library indicated by the loadlib variable below. */
/*=====*/
arg file .
file = strip(file, , " ")
loadlib = "loadlib.with.locate.module"
tempfile = userid()||".TEMP.FILE"
I = "Y"
do forever
  address ispexec
  'addpop row(1) column(1)'
  'display panel (locate)'
  if rc = 8 then exit
  'rempop'
  address tso
  msg = ""
  hexastring = Ø
  if col1 = "" then col1 = 1
  if col2 = "" then col2 = 9999
  if rec1 = "" then rec1 = 1
  if rec2 = "" then rec2 = 99999999
  if col2 < col1 then do
    msg="Column2 cannot be smaller than column1"

```

```

        iterate
    end
    if rec2 < rec1 then do
        msg="Record2 cannot be smaller than record1"
        iterate
    end
    str = strip(stri)
    lstr= length(str)
    if left(str,2) = "x' " | left(str,2) = "X' " then do
        hexastring = 1
        str = strip(str,"T"," ")
        str = translate(substr(str,3))
        lstr= length(str)
        if datatype(str,"X") <> 1 then do
            msg="Invalid hexadecimal string"
            iterate
        end
        lok = lstr // 2
        if lok <> 0 then do
            msg="Odd number of characters in hexadecimal string"
            iterate
        end
        lstr = lstr / 2
    end
    if lstr > col2 - col1 + 1 then do
        msg="String is longer than column search zone"
        iterate
    end
    if msg = "" then leave
end

if hexastring = 1 then l = 'X'
parm = "" right(col1,4,"0") || right(col2,4,"0") ||,
        right(rec1,8,"0") || right(rec2,8,"0") ||,
        l || str""

xx = msg(off)
"free dd (temp1)"
"alloc da('tempfile') dd(temp1) new reuse blksize(8000),
    lrecl(80) recfm(f,b) dsorg(ps) space(1 1) tracks delete"
if rc <> 0 then do
    say "Error "rc" allocating" tempfile
    exit
end

queue "//userid()0 JOB LOCATE,MSGCLASS=X,CLASS=A"
queue "//STEP0 EXEC PGM=LOCATE,"
queue "// PARM="parm
queue "//STEPLIB DD DISP=SHR,DSN="loadlib
queue "//INFILE DD DISP=SHR,DSN="file

```

```
queue "//SYSPRINT DD SYSOUT=*"
queue ""
```

```
"execio * diskw temp1 (finis"
"submit 'tempfile'"
"free dd (temp1)"
say "Job" userid()"Ø submitted"
exit
```

## LOCATE ISPF PANEL

```
)ATTR
_ TYPE(INPUT) CAPS(ON) JUST(LEFT) COLOR(RED)
$ TYPE(INPUT) CAPS(OFF) JUST(LEFT) COLOR(RED)
# TYPE(INPUT) CAPS(ON) JUST(RIGHT) COLOR(RED)
? TYPE(TEXT) INTENS(HIGH) SKIP(ON) COLOR(PINK)
% TYPE(TEXT) INTENS(HIGH) SKIP(ON) COLOR(YELLOW)
+ TYPE(TEXT) INTENS(LOW) SKIP(ON) COLOR(GREEN)
! TYPE(OUTPUT) CAPS(OFF) SKIP(ON) COLOR(WHITE)
)BODY WINDOW(7Ø, 17)
+
? File.....:_FILE +
+
? Search for:$STRI +
? (Enter text string or begin with X' for hexadecimal)
+
+ Ignore case - only valid for text (Y,N).:_I
+
% Search columns Search records
+ First column.#COL1+ First record.#REC1 +
+ Last column.#COL2+ Last record.#REC2 +
+
!MSG
+ Enter - execute PF3/15 cancel
)INIT
&ZWINTTL = 'Locate a string in a file'
)PROC
&ver='Y,N'
VER(&FILE, nonblank, dsname)
VER(&STRI, nonblank)
VER(&STRI, nonblank)
VER(&I, NONBLANK, listv, &ver)
VER(&COL1, num)
VER(&COL2, num)
VER(&REC1, num)
VER(&REC2, num)
)END
```

## Calling the ANTRQST macro

The following program, CQPROG, is an example of how to call the ANTRQST macro. This macro issues calls to the System Data Mover API, and can be used to issue commands for XRC, PPRC, SNAPSHOT, and FLASHCOPY functions.

This particular program issues PPRC CQUERY commands. These can also be issued as TSO CQUERY commands (or alternatively ISMF can obtain the same information), but the program has several advantages over this:

- The TSO command can be quite slow if you are issuing a lot of commands.
- The TSO command mirrors all of its displays on the system console – this can be quite a problem if you are repeatedly issuing a lot of these commands.
- This program uses an easily-tailored input file to select the volumes you want to report on.

The program reads a list containing addresses (required) and volume names (optional) that you want to issue CQUERY commands against, while using PPRC to mirror DASD volumes. TSO variables are updated and displayed using a small piece of REXX. CQPROG will need to be authorized in the IKJTSOx member.

### CQUERY PROGRAM

```
TITLE 'CQPROG - ISSUE "CQUERY" COMMANDS VIA "ANTRQST" MACRO'
*****
* CQPROG:  ISSUE CQUERY COMMANDS AGAINST SELECTED DEVICES, USING      *
*          THE 'ANTRQST' MACRO (SEE 'DFSMSDFP ADVANCED SERVICES').    *
*                                                                 *
*          THIS WILL PREVENT THE OUTPUT OF 'CQUERY' COMMANDS FROM   *
*          FLOODING THE SYSLOG.                                       *
*                                                                 *
*          THE CALLER WILL NEED READ ACCESS TO THE FACILITY CLASS    *
*          PROFILE 'STGADMIN.ANT.PPRC.COMMANDS' .                     *
```

```

*
*          THE 'ANTAS000' ADDRESS SPACE MUST BE ACTIVE.
*
* PARS:    NONE
*
* INPUT:   DATASET CONTAINING ADDRESSES OF DISKS TO BE 'CQUERY' ED.
*
* OUTPUT:  TSO 'PPRQNN' VARIABLES, FOR PANEL OR REXX DISPLAY.
*          RETURN CODES -
*
*          0 - SUCCESSFUL COMPLETION
*          4 - AT LEAST ONE NON-ZERO RETURN CODE FROM 'ANTROST'
*          8 - UNABLE TO OPEN INPUT FILE 'VOLIN'
*          12 - UNABLE TO UPDATE TSO VARIABLE
*****
          PRINT NOGEN
*****
* HOUSEKEEPING. . .
*****
CQPROG    CSECT
CQPROG    AMODE 31
CQPROG    RMODE 24
          BAKR R14,0          SAVE CALLER DATA ON STACK
          LR   R12,R15        GET ENTRY POINT
          USING CQPROG,R12    ADDRESSABILITY
*
          OPEN (VOLIN,(INPUT)) OPEN LIST OF ADDRESSES
          LTR  R15,R15        OK?
          BNZ  BADOPEN        NO...GO AND SET RC=8
          LA   R7,VOLINREC    ADDRESSABILITY TO INPUT RECS
          USING VOLINDSC,R7
          LA   R6,QRYINFO
          USING ANTDATA,R6    ADDRESSABILITY TO RETURNED DATA
READLOOP  DS    0H
*
          GET  VOLIN,VOLINREC READ LIST OF ADDRESSES
*
          AP   PPRTOT,P01     NUMBER OF INPUT RECORDS READ
          MVC  VALDEVN,VOLADDR SAVE ADDRESS
          BAL  R9,CONVCUU     CONVERT EBCDIC ADDRESS TO HEX
*
          ANTROST I LK=PPRC,          PPRC CALL
          REQUEST=PQUERY,            CQUERY TYPE REQUEST
          DEVN=DEVN,                 FIELD CONTAINING DEVICE ADDR.
          QRYSIZE=QRYSIZE,           LENGTH OF BUFFER FOR RESPONSE
          QRYINFO=QRYINFO,           BUFFER FOR RESPONSE
          RETINFO=RETINFO,           RETURNCODE INFO
          ALET=0,                     X
          ASYNCH=NO,                 X

```

```

                ECB=NO_ECB,                                X
                BITMAP=NOBITMAP,                          X
                PATHS=NO,                                  X
                WAITTIME=Ø
*
MVC    RETCD, RETINFO          SAVE RETURNCODE
MVC    RSNCD, RETINFO+4        SAVE REASONCODE
CLC    RETCD, =A(RQST_PQUERY_QRYSIZE_BIG_ENOUGH)        CALL OK?
BE     CALLOK
*-----*
* CALL TO 'ANTRQST' GAVE NONZERO RETURN CODE. THERE WILL BE AN ERROR *
* MESSAGE IN THE RESPONSE - DISPLAY THAT... *
*-----*
ANTFAIL DS    ØH
        MVC    RETC, RC4          SET RC=4
        MVC    VALDEVN+5(5), =C' ANTPØ' SET UP MSG PREFIX
        XR     R2, R2             CLEAR R2
        IC    R2, RETINFO+8      GET LENGTH OF FOLLOWING MSG
        C     R2, =F' 56'        TOO LONG FOR MSG AREA?
        BL    MOVEIT             NO... MOVE MSG
        LA    R2, 56             YES.. TRUNCATE...
MOVEIT  DS    ØH
        EX    R2, EXMVC          MOVE MSG TO VARIABLE AREA
        B     UPDATEIT          NOW GO AND UPDATE
*-----*
* CALL TO 'ANTRQST' OK - UPDATE TSO VARIABLES... *
*-----*
CALLOK DS    ØH
        CLC    ANTSTAT(6), =CL6' DUPLEX' #VOLUMES IN DUPLEX STATE
        BE     UPDTDUPL
        CLC    ANTSTAT(7), =CL7' SIMPLEX' #VOLUMES IN SIMPLEX STATE
        BE     UPDTSIMP
        CLC    ANTSTAT(7), =CL7' SUSPEND' #VOLUMES IN SUSPEND STATE
        BE     UPDTSUSP
        CLC    ANTSTAT(7), =CL7' PENDING' #VOLUME IN SUSPEND STATE
        BE     UPDTPEND
        B     UPDTTSO           UNKNOWN STATUS - SKIP COUNTS
UPDTDUPL DS    ØH
        MVC    VALASTER(3), =C' (*)'  HIGHLIGHT DUPLEX PAIRS
        AP    DUPLCT, PØ1        BUMP DUPLEX COUNT
        B     UPDTTSO
UPDTSIMP DS    ØH
        AP    SIMPCT, PØ1        BUMP SIMPLEX COUNT
        B     UPDTTSO
UPDTPEND DS    ØH
        AP    PENDCT, PØ1        BUMP PENDING COUNT
        B     UPDTTSO
UPDTSUSP DS    ØH
        AP    SUSPCT, PØ1        BUMP SUSPEND COUNT

```



```

UPDTTSDS      ØH
MVC      VALDEVN, ANTDEVN
MVC      VALLEVEL, ANTLEVEL
MVC      VALSTAT, ANTSTAT
MVC      VALVOLID, VOLVOLID
BAL      R9, CHKPATHS          CHECK PATHS
UPDATEIT DS      ØH
UNPK     NAMESTEM(3), PPRTOT
OI       NAMESTEM+2, X' FØ'    SET CORRECT SIGN
BAL      R9, UPDTPVAR
LTR      R15, R15              OK?
BNZ      SETRC12               NO... SET RC=12
MVI      VALDEVN, C' '         CLEAR OUT DETAIL AREA
MVC      VALDEVN+1 (VALUELN-1), VALDEVN
B        READLOOP

*-----*
* CLOSE INPUT DATASET...      *
*-----*
CLOSE    DS      ØH
         CLOSE VOLIN

*-----*
* RETURN TO CALLER WITH RELEVANT RC... *
*-----*
RETURN   DS      ØH
         L       R15, RETC      LOAD RETURN CODE
         PR      ,             RESTORE CALLER DATA, RETURN

*-----*
* COULDN'T OPEN INPUT FILE 'VOLIN'... *
*-----*
BADOPEN  DS      ØH
         MVC     RETC, RC8      SET RC=8
         B       RETURN

*-----*
* ERROR UPDATING TSO VARIABLE(S)... RC=12... *
*-----*
SETRC12 DS      ØH
         MVC     RETC, RC12     SET RC=12
         B       CLOSE

*-----*
* E-O-F ON INPUT - SET TSO VARIABLES FOR # IN SIMPLEX, # IN DUPLEX, *
* # PENDING, # SUSPENDED AND TOTAL INPUT RECORDS READ... *
*-----*
EOFVOLIN DS      ØH
         MVC     NAME(5), =C' PPSIM'
         MVC     NAMELEN, =A(5)
         UNPK    VALDEVN(3), SIMPCT
         OI      VALDEVN+2, X' FØ'
         MVC     VALUELEN, =A(3)
         BAL     R9, UPDTPVAR   UPDATE PPSIM

```

```

MVC NAME(5), =C' PPDUP'
MVC NAMELEN, =A(5)
UNPK VALDEVN(3), DUPLCT
OI VALDEVN+2, X' FØ'
MVC VALUELEN, =A(3)
BAL R9, UPDVAR UPDATE PPDUP
MVC NAME(5), =C' PPPND'
MVC NAMELEN, =A(5)
UNPK VALDEVN(3), PENDCT
OI VALDEVN+2, X' FØ'
MVC VALUELEN, =A(3)
BAL R9, UPDVAR UPDATE PPPND
MVC NAME(5), =C' PPSUS'
MVC NAMELEN, =A(5)
UNPK VALDEVN(3), SUSPCT
OI VALDEVN+2, X' FØ'
MVC VALUELEN, =A(3)
BAL R9, UPDVAR UPDATE PPSUS
MVC NAME(5), =C' PPTOT'
MVC NAMELEN, =A(5)
UNPK VALDEVN(3), PPRTOT
OI VALDEVN+2, X' FØ'
MVC VALUELEN, =A(3)
BAL R9, UPDVAR UPDATE PPTOT
MVC NAME(7), =C' PPRPATH'
MVC NAMELEN, =A(7)
MVC VALDEVN(1), PPRPATH
MVC VALUELEN, =A(1)
BAL R9, UPDVAR UPDATE PPRPATH
B CLOSE GO AND CLOSE VOLIN

```

\*\*\*\*\*

```

*           + + S U B R O U T I N E + + +           *
* CONVERT CHARACTER CUU (EG 'Ø94F') INTO ITS BINARY EQUIVALENT, THEN *
* PLACE IT IN THE 'DEVN' FIELD. . . *
*****

```

```

CONVCUU DS ØH
TR VOLADDR(4), TRTAB CONV. C' A->F' INTO X' A->F'
XC DWORD, DWORD CLEAR OUT WORKAREA
PACK DWORD+4(4), VOLADDR(5) REMOVE ZONES
L R8, DWORD+4 LOAD 'ØØCCUUØØ'
SRL R8, 8 SHIFT OUT TRAILING 'ØØ'
STH R8, DEVN SAVE BINARY CUU VALUE
BR R9 RETURN FROM SUBROUTINE

```

\*\*\*\*\*

```

*           + + S U B R O U T I N E + + +           *
* CHECK THE PATHS' FILEDS IN CASE THERE IS A POTENTIAL PROBLEM. AS WE *
* ARE USING ONLY 2 PATHS WE' LL CHECK ONLY PATH1 AND PATH2. . . *
*****

```

```

CHKPATHS DS ØH

```

```

MVC VALPATHS, =15C' ' RESET FIELD
CLC ANTSTAT1, =C' ' STATUS1 = SPACES?
BE PATH10K YES..GOOD
CLC ANTSTAT1, ESTABLISHED STATUS1 = ESTABLISHED?
BE PATH10K YES..GOOD
MVC VALPATHS, PATHCHEK NO...SHOW POSSIBLE ERROR
MVI PPRPATH, C' Y' SHOW POSSIBLE ERROR
BR R9 RETURN FROM SUBROUTINE
PATH10K DS ØH
CLC ANTSTAT2, =C' ' STATUS2 = SPACES?
BER R9 YES..RETURN
CLC ANTSTAT2, ESTABLISHED STATUS2 = ESTABLISHED?
BER R9 YES..RETURN
MVC VALPATHS, PATHCHEK NO...SHOW POSSIBLE ERROR
MVI PPRPATH, C' Y' SHOW POSSIBLE ERROR
BR R9 RETURN FROM SUBROUTINE
*****
* + + S U B R O U T I N E + + + *
* CALL 'IKJCT441' TO UPDATE TSO VARIABLES... *
*****
UPDTVAR DS ØH
*
LINK EP=IKJCT441, PUT VALUE INTO VARIABLE X
PARAM=(ECODE, X
NAMEPTR, X
NAMELEN, X
VALUEPTR, X
VALUELEN, X
TOKEN), X
VL=1
*
BR R9 RETURN
*-----*
*
LTORG LITERAL POOL
*-----*
* REGISTERS EQUATES, ETC... *
*-----*
EXMVC MVC VALDEVN+1Ø(Ø), RETINFO+9 EXECUTED MOVE
SIMPCT DC PL2' Ø'
DUPLCT DC PL2' Ø'
PENDCT DC PL2' Ø'
SUSPCT DC PL2' Ø'
PPRTOT DC PL2' Ø'
PØ1 DC PL1' 1'
DEVN DC XL2' ØØØØ'
NO DC CL3' NO '
RETCD DS F

```

```

RSNCD    DS    F
FWORD    DS    F
DWORD    DS    D
RETC     DC    F' 0'
RC4      DC    F' 4'
RC8      DC    F' 8'
RC12     DC    F' 12'
RETINFO  DS    XL100
QRYINFO  DS    XL1900
QRYINFOL EQU  *-QRYINFO
QRYSIZE  DC    Y(QRYINFOL)
NOBITMAP DC    CL3' NO '
PATHCHEK DC    CL13' <CHECK PATHS>'
PPRPATH  DC    CL1' N'
ESTABLISHED DC CL2' 01'

```

\*

\*

```

          0 1 2 3 4 5 6 7 8 9 A B C D E F
TRTAB    DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 0
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 1
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 2
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 3
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 4
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 5
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 6
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 7
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 8
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' 9
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' A
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' B
          DC    X' FF0A0B0C0D0E0FFFFFFFFFFFFFFFFFFFF' C (ABCDEF)
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' D
          DC    X' FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF' E
          DC    X' 00010203040506070809FFFFFFFFFFFF' F (0123456789)

```

\*

\*-----\*

\* IKJCT441 PARMLIST (TSO VARIABLE ACCESS ROUTINE)... \*

\*-----\*

```

NAME      DC    C' PPRQ'          TSO VARIABLE NAME (FIXED)
NAMESTEM  DS    CL3              TSO VARIABLE NAME (VARIABLE)
NAMELN    EQU   *-NAME          VARIABLE LENGTH
NAMELEN   DC    A(NAMELN)       "          "
VALDEVN   DC    CL4' '          DEVICE NUMBER (IE ADDRESS)
          DC    CL6' '
VALLEVEL  DC    CL9' '          LEVEL
          DC    CL5' '
VALSTAT   DC    CL7' '          STATUS
          DC    CL3' '
VALASTER  DC    CL3' '          HIGHLIGHT IF DUPLEX
          DC    CL10' '

```

VALVOLID	DC	CL6' '	VOLID
	DC	CL2' '	
WHEREPAT	EQU	*-VALDEVN	OFFSET TO VALPATHS FIELD
VALPATHS	DC	CL13' '	POSSIBLE PATH PROBLEM
VALUELN	EQU	*-VALDEVN	VALUE LENGTH
VALUELEN	DC	A(VALUELN)	" "
NAMEPTR	DC	A(NAME)	POINTER TO VARIABLE NAME
VALUEPTR	DC	A(VALDEVN)	POINTER TO VARIABLE VALUE
TOKEN	DC	F' Ø'	TOKEN (UNUSED HERE)
ECODE	DC	A(TSVEUPDT)	ENTRY CODE FOR 'SET'
*-----*			
* INPUT DCB AND BUFFER... *			
*-----*			
VOLIN	DCB	DDNAME=VOLIN, RECFM=FB, DSORG=PS, MACRF=(GM), EODAD=EOFVOLIN	X X X X
VOLINREC	DS	CL8Ø	
*-----*			
* REGISTER EQUATES, DSECTS, ETC... *			
*-----*			
*			
YREGS			
*			
ANTRQSTL			
*			
IKJTSVT			
*			
ANTDATA	DSECT		
ANTDEVN	DS	CL4	DEVICE NUMBER
	DS	CL1	
ANTLEVEL	DS	CL9	LEVEL
	DS	CL1	
ANTSTAT	DS	CL1Ø	STATUS
	DS	CL1	
ANTPTHST	DS	CL8	PATH STATUS
	DS	CL1	
ANTPSSID	DS	CL4	PRIMARY SSID
	DS	CL1	
ANTPCCA	DS	CL2	PRIMARY CCA
	DS	CL1	
ANTPSER#	DS	CL12	PRIMARY SERIAL NUMBER
	DS	CL1	
ANTSSSID	DS	CL4	SECONDARY SSID
	DS	CL1	
ANTSCCA	DS	CL2	SECONDARY CCA
	DS	CL1	
ANTSSER#	DS	CL12	SECONDARY SERIAL NUMBER

```

        DS      CL1      ' , '
        DS      CL1
        DS      CL1      ' , '
        DS      CL1
        DS      CL1      ' , '
ANTPATHS DS      CL1      NUMBER OF PATHS
        DS      CL1
ANTSADE1 DS      CL8      SAID/DEST (1)
        DS      CL1
ANTSTAT1 DS      CL2      PATH STATUS (1)
        DS      CL1
ANTSADE2 DS      CL8      SAID/DEST (2)
        DS      CL1
ANTSTAT2 DS      CL2      PATH STATUS (2)
        DS      CL1
ANTSADE3 DS      CL8      SAID/DEST (3)
        DS      CL1
ANTSTAT3 DS      CL2      PATH STATUS (3)
        DS      CL1
ANTSADE4 DS      CL8      SAID/DEST (4)
        DS      CL1
ANTSTAT4 DS      CL2      PATH STATUS (4)
*
VOLINDSC DSECT
        DS      CL1
VOLADDR  DS      CL4      ADDRESS
        DS      CL1
VOLDEVT  DS      CL4      DEVI CETYPE
        DS      CL1
VOLALLOC DS      CL7      ONLINE/ALLOC
        DS      CL1
VOLVOLID DS      CL6      VOLID!
        DS      CL1
*
        END

```

## PPRQ REXX

The following REXX code will execute the CQUERY program. It expects a sysid (from a valid list), which it then uses as a qualifier for a dataset name to hold the list of addresses that you want to issue the PPRC commands for.

```

/* Rexx tso issue CQUERY-type commands */
lastct = 0

sys = "PRD1 PRD2 PRD3 TST1 TST2"

```

```

lin = "======"

Parse Upper Arg dsn1 whatsthis
If whatsthis <> "" Then Do
  Say "Extraneous info passed: '"whatsthis'" "
  Return 8
End
If Length(dsn1) = 4 Then
  sysid = dsn1
Else
Parse Value dsn1 WITH pref '.' pprc '.' sysid '.' stuff

If Pos(sysid,sys) = 0 Then Do
  Say "==> Invalid SYSid passed: '"sysid'" "
  Return 8
End

Address "TS0"

x = Outtrap("null.", 10, "NOCONCAT")
"FREE FI (VOLIN)"
dsn = "" SYSG.PPRC."sysid".INPUT"
"ALLOC FI (VOLIN) DA("dsn") SHR"          /* Alloc list of addresses */

Do Forever
  Address "TS0"
  "CALL 'SYSG.LINKLIB(CQPROG)' "
  cqrc = rc

  If cqrc = 8 Then
    pprq001 = "Unable to open input dataset "dsn"... "

  If cqrc = 12 Then Do
    lin = "NB: Unable to update at least one TS0 variable..."
    aa = pptot + 3
    aa = Right("00"aa, 3)
    Interpret "PPRQ"aa" = lin"
  End

  pprdt = Date('E')
  ppmt = Time()
  ppsy = sysid
  If cqrc = 0 Then Do
    ppch = ppdup||" ("ppdup-lastct)" /* Calc change since last time */
    ppch = Strip(ppch, "L", "0")      /* Remove leading 0s */
    lastct = ppdup
    aa = pptot + 1
    aa = Right("00"aa, 3)
    Interpret "PPRQ"aa" = lin"      /* Underline last one */

```

```

    ppsim = Strip(ppsim, "L", "0")          /* Remove leading 0s */
    pppnd = Strip(pppnd, "L", "0")
    ppsus = Strip(ppsus, "L", "0")
    pptot = Strip(pptot, "L", "0")
End
pathmsg = " "
If pprpath = "Y" Then /* If prog found path problem, flag this */
    pathmsg = "**** POSSIBLE PATH PROBLEMS - PLEASE CHECK ****"
Address "ISPEXEC"
"ISPEXEC DISPLAY PANEL(PPRCQPAN)"
If rc = 8 Then Return /* PF3 = exit */
End

Address "TS0"
x = Outtrap("null.", 10, "NOCONCAT")
"FREE FI(VOLIN)"

Return /* That's it... */

```

## PPRCQPAN PANEL

```

)ATTR DEFAULT(%+_)
    ! TYPE(OUTPUT) INTENS(LOW) CAPS(OFF)
    ~ TYPE(OUTPUT) INTENS(HIGH)
    $ TYPE(OUTPUT) INTENS(HIGH) COLOR(RED) HILITE(BLINK)
    # AREA(SCRL) EXTEND(ON)
)BODY WIDTH(80) EXPAND(@@)
~pprdt  %@-@ PPRC CQUERY DISPLAY (-ppsy%) @-@-pprtm  %
+
+          $pathmsg
+
+   SIMPLEX: ~ppsim+DUPLICATE: ~ppch
+PENDING: ~pppnd+SUSPEND: ~ppsus+TOTAL: ~pptot
+
+          =DEVICE=      ==LEVEL==      ===STATUS===      =VOLSER=
#SCRLAREA
#
#
#
#
#
#
#
#
#
#
#
#
+
+Press%PF3+to%End,%Enter+to Refresh...
)AREA SCRLAREA
+
+   ~PPRQ001
+

```



```

+          -PPRQ002
+          -PPRQ003
+          -PPRQ004
+          -PPRQ005
          .
          .
          .
    and so on up to
          .
          .
          .
+          -PPRQ398
+          -PPRQ399
+          -PPRQ400
)INIT
)PROC
)END

```

## INPUT

Format of the input dataset. The addresses start in column 2, the volids in column 20; all other information is ignored.

```

1033 3390 ALLOC  PFPP30  ( 3, 339)  PRI VATE
1034 3390 ALLOC  PFPP31  ( 3, 339)  PRI VATE
1035 3390 ALLOC  PFPP10  ( 3, 339)  PRI VATE
1036 3390 ALLOC  PFPP11  ( 3, 339)  PRI VATE

```

---

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*Systems Programmer (UK)*

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---

## Using Pretty Good Privacy (PGP) to encrypt/decrypt and sign your MVS data

### INTRODUCTION TO CRYPTOGRAPHY

Pretty Good Privacy (PGP) is a freeware public key cryptography program developed by Phillip Zimmerman (1991) under the GNU licence, which is now available for z/OS.

Cryptography is the science of using mathematics to encrypt and decrypt data. Cryptography enables you to store sensitive information or transmit it across insecure networks (like the Internet) so that it cannot be read by anyone except the intended recipient.

A cryptographic algorithm, or cipher, is a mathematical function used in the encryption and decryption process. A cryptographic algorithm works in combination with a key (a word, number, or phrase) to encrypt the plaintext. The same plaintext encrypts to different ciphertext with different keys. The security of encrypted data is entirely dependent on two things – the strength of the cryptographic algorithm and the secrecy of the key.

Cryptanalysis is the science of analysing and breaking secure communication.

## Encryption

### *Symmetric-key encryption*

In conventional cryptography, also called secret-key or symmetric-key encryption, one key is used both for encryption and decryption.

The Data Encryption Standard (DES) is an example of a conventional cryptosystem that has been widely deployed by the US Government and the banking industry.

An extremely simple example of conventional cryptography is a substitution cipher. A substitution cipher substitutes one piece of information for another. This is most frequently done by offsetting letters of the alphabet.

So, starting with ABCDEFGHIJKLMNOPQRSTUVWXYZ and sliding everything up by three, you get DEFGHIJKLMNOPQRSTU VWXYZABC – where D=A, E=B, F=C, and so on.

Using this scheme, the plaintext SECRET encrypts as VHFUHW.

To allow someone else to read the ciphertext, you tell them that the key is 3.

Conventional encryption has benefits – it is quite simple and very fast.

But it has some major limitations – for a sender and recipient to communicate securely using conventional encryption, they must agree on a key and keep it secret between themselves. If they are in different physical locations, they must trust a courier, the phone, or some other secure communications medium to prevent the disclosure of the secret key during transmission. Anyone who overhears or intercepts the key in transit can later read, modify, and forge all information encrypted or authenticated with that key.

The persistent problem with conventional encryption is key distribution – how do you get the key to the recipient without someone intercepting it?

### *Asymmetric-key encryption*

The problems of key distribution are solved by public-key cryptography, the concept of which was introduced by Whitfield Diffie and Martin Hellman in 1975.

Public-key cryptography uses a pair of keys – a public key, which encrypts data, and a corresponding private key, for decryption.

Because it uses two keys, it is also called asymmetric cryptography. You publish your public key to the world while keeping your private key secret. Anyone with a copy of your public key can then encrypt information that only you can read.

It is computationally infeasible to deduce the private key from the public key.

Anyone who has a public key can encrypt information but cannot decrypt it. Only the person who has the corresponding private key can decrypt the information.

The main benefit of public-key cryptography is that it allows people who have no pre-existing security arrangement to exchange messages securely.

The need for sender and receiver to share secret keys via some secure channel is eliminated; all communications involve only public keys, and no private key is ever transmitted or shared.

Some examples of public-key cryptosystems are :

- Elgamal (named for its inventor, Taher Elgamal).
- RSA (named for its inventors, Ron Rivest, Adi Shamir, and Leonard Adleman).
- Diffie-Hellman (named, you guessed it, for its inventors).
- DSA, the Digital Signature Algorithm (invented by David Kravitz).

### *How PGP works*

PGP combines some of the best features of both conventional and public-key cryptography – PGP is a hybrid cryptosystem.

When a user encrypts plaintext with PGP, PGP first compresses the plaintext.

Data compression saves transmission time and disk space and, more importantly, strengthens cryptographic security. Most cryptanalysis techniques exploit patterns found in the plaintext to crack the cipher. Compression reduces these patterns in the plaintext, thereby greatly enhancing resistance to cryptanalysis.

Because the public key encryption algorithm is much slower than conventional single-key encryption, encryption is better accomplished by using a high-quality fast conventional single-key encryption algorithm to encipher the message.

In a process invisible to the user, PGP creates a temporary random key (the session key, which is a one-time-only secret key) just for this 'session', which is used to conventionally encipher the plaintext file.

The session key works with a very secure, fast, conventional encryption algorithm to encrypt the plaintext; the result is ciphertext.

Once the data is encrypted, the session key is then encrypted to the recipient's public key. This public key-encrypted session key is transmitted along with the ciphertext to the recipient.

Decryption works in the reverse. The recipient's copy of PGP uses his or her private key to recover the session key, which PGP then uses to decrypt the conventionally encrypted ciphertext.

The combination of the two encryption methods combines the convenience of public-key encryption with the speed of conventional encryption.

Conventional encryption is about 10,000 times faster than public-key encryption. Public-key encryption in turn provides a solution to key distribution and data transmission issues. Used together, performance and key distribution are improved without any sacrifice in security.

### Digital signatures

A major benefit of public key cryptography is that it provides a method for employing digital signatures.

Digital signatures let the recipient of information verify the authenticity of the information's origin, and also verify that the information was not altered while in transit.

Thus, public key digital signatures provide authentication and data integrity.

The basic manner in which digital signatures are created is:

- The signature algorithm uses your private key to create the signature and the public key to verify it.
- If the information can be decrypted with your public key, then it must have originated from you.

This system has some problems. It is slow, and it produces an enormous volume of data – at least double the size of the original information.

An improvement on this scheme is the addition of a one-way hash function in the process.

A one-way hash function takes variable-length input and produces a fixed-length output. The hash function ensures that, if the information is changed in any way (even by just one bit) an entirely different output value is produced.

### *How PGP works*

PGP uses a cryptographically strong hash function on the plaintext the user is signing. This generates a fixed-length data item known as a message digest.

Then PGP uses the digest and the private key to create the 'signature'. PGP transmits the signature and the plaintext together. PGP can also encrypt the plaintext using the recipient's public key.

Upon receipt of the message, the recipient uses PGP to recompute the digest, thus verifying the signature.

## PGP INSTALLATION ON Z/OS

### Downloading PGP for OS/390 from the Web

PGP for OS/390 and MVS is freeware which can be downloaded from the Alan Nichols Web site at <http://s390.nichols.de/pgp/index.html>.

The current version of PGP for OS/390 (in June 2003) is 2.6.3is. You have to download a PAX file (pgp263is.pax.Z – 1.16 MB) on your PC.

### Allocating and mounting an HFS file on z/OS

At this point, you should allocate and mount an HFS file on your z/OS system to upload PGP:

```
//STEP01 EXEC PGM=IEFBR14
//*
//MKFS DD DISP=(,CATLG),DSN=SYS2.OMVS.U.PGP.HFS,
// SPACE=(CYL,(005,01,1)),UNIT=SYSALLDA,
// DSNTYPE=HFS
//*
```

```
//STEP02 EXEC PGM=IKJEFT01
//SYSPROC DD DISP=SHR,DSN=SYS1.SBPXEXEC
//SYSTSPRT DD SYSOUT=*
//SYSTSIN DD *
  oshell mkdir /u/pgp
  mount filesystem('SYS2.OMVS.U.PGP.HFS') type(HFS) +
    mode(rdwr) mountpoint('/u/pgp')
//*
```

## Uploading the PAX file to z/OS

Now you have to use FTP to transfer the PAX file to your z/OS system:

```
D:\Migration\Dustbin>ftp 126.9.41.65
Connect to 126.9.41.65.
220-FTPD1 IBM FTP CS V1R4 at smvs.ctrcepl.caisse-epargne.fr, 11:33:38
on 2003-06-13.
220 Connection will close if idle for more than 5 minutes.
Utilisateur (126.9.41.65: (none)) : SXSP001
331 Send password please.
Mot de passe :
230 SXSP001 is logged on. Working directory is "SXSP001.".
ftp> cd /u/pgp
250 HFS directory /u/pgp is the current working directory
ftp> bin
200 Representation type is Image
ftp> put pgp263is.pax.Z
200 Port request OK.
125 Storing dataset /u/pgp/pgp263is.pax.Z
250 Transfer completed successfully.
ftp : 1217601 octets sent in 21,16 seconds at 57,55 Ko/sec.
ftp> quit
221 Quit command received. Goodbye.
```

```
D:\Migration\Dustbin>
```

## Uncompressing the PAX file

You should now use a Telnet session to log on to your USS system and uncompress the PAX file:

```
EZYTE27I Login: SXSP001
EZYTE28I SXSP001 Password:
IBM
Licensed Material - Property of IBM
5694-A01 (C) Copyright IBM Corp. 1993, 2001
```

(C) Copyright Mortice Kern Systems, Inc., 1985, 1996.

(C) Copyright Software Development Group, University of Waterloo, 1989.

All Rights Reserved.

U.S. Government users - RESTRICTED RIGHTS - Use, Duplication, or Disclosure restricted by GSA-ADP schedule contract with IBM Corp.

IBM is a registered trademark of the IBM Corp.

```
SXSP001:/: >cd /u/pgp
```

```
SXSP001:/u/pgp: >uncompress pgp263i.s.pax.Z
```

```
SXSP001:/u/pgp: >pax -rf pgp263i.s.pax
```

```
SXSP001:/u/pgp: >ls -al
```

```
total 5912
```

```
drwx----- 3 SXSP001 OMVSGRP      8192 Jun 13 12:45 .
drwx----- 8 SXSP001 OMVSGRP      8192 Jun  5 18:04 ..
drwxrwxrwx  5 SXSP001 OMVSGRP      8192 Sep  4 2002 pgp
-rw-r----- 1 SXSP001 OMVSGRP 2999808 Jun 13 12:34 pgp263i.s.pax
```

```
SXSP001:/u/pgp: >
```

```
SXSP001:/u/pgp: >cd pgp
```

```
SXSP001:/u/pgp/pgp: >ls -al
```

```
total 512
```

```
drwxrwxrwx  5 SXSP001 OMVSGRP      8192 Sep  4 2002 .
drwx----- 3 SXSP001 OMVSGRP      8192 Jun 13 12:45 ..
-r-----  1 SXSP001 OMVSGRP      5024 Oct  6 1998 config.txt
drwxrwxrwx  6 SXSP001 OMVSGRP      8192 Oct  6 1998 contrib
-r-----  1 SXSP001 OMVSGRP      5597 Jan  2 1996 de.hlp
drwxrwxrwx  2 SXSP001 OMVSGRP      8192 May 29 2001 doc
-r-----  1 SXSP001 OMVSGRP      4152 Jan  2 1996 en.hlp
-r-----  1 SXSP001 OMVSGRP      4526 Jan  2 1996 es.hlp
-r-----  1 SXSP001 OMVSGRP      4651 Jan  2 1996 fr.hlp
-r-----  1 SXSP001 OMVSGRP     14274 Jan 15 1996 keys.asc
-r-----  1 SXSP001 OMVSGRP    105573 Jan 18 1996 language.txt
-r-----  1 SXSP001 OMVSGRP      2035 Jan 12 1996 pgp.hlp
-r-----  1 SXSP001 OMVSGRP     19651 Jan 18 1996 readme.1st
-rwx----- 1 SXSP001 OMVSGRP       734 Sep  4 2002 readme.mvs
-r-----  1 SXSP001 OMVSGRP      2526 Jan 18 1996 readme.usa
-r-----  1 SXSP001 OMVSGRP     13453 Jan 11 1996 setup.doc
drwxrwxrwx  2 SXSP001 OMVSGRP     16384 Sep  4 2002 src
```

```
SXSP001:/u/pgp/pgp: >
```

## Compiling PGP C/C++ source

First, you should check the value of several USS environment variables in your /etc/profile:

```
export _CC_CNAME=BCCDVR          - The OS/390 C/C++ driver program
export _CC_PLIB_PREFIX=SYS1     - prefix for SYS1.SCEELKEX
```



Then, you should compile the PGP C/C++ sources:

```
SXSP001: /u/pgp/pgp/src: >make clean
rm -f *.o pgp core a.out tags *.err CEEDUMP*
SXSP001: /u/pgp/pgp/src: >make os390
make all CPP=/lib/cpp OBJS_EXT="c370.o" CFLAGS=" -O -DC370 -DHIGHFIRST"
cc -O -DC370 -DHIGHFIRST -DHIGHFIRST pgp.c
cc -O -DC370 -DHIGHFIRST -DHIGHFIRST crypto.c
cc -O -DC370 -DHIGHFIRST keygmt.c
cc -O -DC370 -DHIGHFIRST fileio.c
WARNING CBC3236 ./c370.h:114 Macro name ENOENT has been redefined.
FSUM3065 The COMPILE step ended with return code 4.
cc -O -DC370 -DHIGHFIRST mdfile.c
cc -O -DC370 -DHIGHFIRST more.c
cc -O -DC370 -DHIGHFIRST armor.c
cc -O -DC370 -DHIGHFIRST mpi lib.c
cc -O -DC370 -DHIGHFIRST mpi io.c
cc -O -DC370 -DHIGHFIRST genprime.c
cc -O -DC370 -DHIGHFIRST rsagen.c
cc -O -DC370 -DHIGHFIRST random.c
cc -O -DC370 -DHIGHFIRST idea.c
cc -O -DC370 -DHIGHFIRST passwd.c
cc -O -DC370 -DHIGHFIRST md5.c
cc -O -DC370 -DHIGHFIRST system.c
cc -O -DC370 -DHIGHFIRST language.c
cc -O -DC370 -DHIGHFIRST getopt.c
cc -O -DC370 -DHIGHFIRST keyadd.c
cc -O -DC370 -DHIGHFIRST config.c
cc -O -DC370 -DHIGHFIRST keymaint.c
cc -O -DC370 -DHIGHFIRST keymaint.c
cc -O -DC370 -DHIGHFIRST charset.c
cc -O -DC370 -DHIGHFIRST randpool.c
cc -O -DC370 -DHIGHFIRST noise.c
cc -O -DC370 -DHIGHFIRST zbits.c
cc -O -DC370 -DHIGHFIRST zdeflate.c
cc -O -DC370 -DHIGHFIRST zfile_io.c
cc -O -DC370 -DHIGHFIRST zglobals.c
cc -O -DC370 -DHIGHFIRST zinflate.c
cc -O -DC370 -DHIGHFIRST zip.c
cc -O -DC370 -DHIGHFIRST zipup.c
cc -O -DC370 -DHIGHFIRST ztrees.c
cc -O -DC370 -DHIGHFIRST zunzip.c
cc -O -DC370 -DHIGHFIRST rsaglue1.c
cc -O -DC370 -DHIGHFIRST c370.c
cc -o pgp pgp.o crypto.o keygmt.o fileio.o mdfile.o more.o armor.o
mpilib.o mpi
io.o genprime.o rsagen.o random.o idea.o passwd.o md5.o system.o
language.o g
etopt.o keyadd.o config.o keymaint.o charset.o randpool.o noise.o
```

```
zbi ts.o zdef
late.o zfile_io.o zglobals.o zinflate.o zip.o zipup.o ztrees.o zunzip.o
rsaglue
1.o c370.o
SXSP001: /u/pgp/pgp: >
```

This step produces the PGP module:

```
SXSP001: /u/pgp/pgp/src: >ls -al pgp
-rwxrwxrwx  1 SXSP001  OMVSGRP  765952 Jun 13 18:06 pgp
SXSP001: /u/pgp/pgp/src: >
```

You can now call PGP:

```
SXSP001: /u/pgp/pgp/src: >pgp -h
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/13 17:09 GMT
```

Usage summary:

```
To encrypt a plaintext file with recipient's public key, type:
  pgp -e textfile her_userid [other userids] (produces textfile.pgp)
To sign a plaintext file with your secret key:
  pgp -s textfile [-u your_userid] (produces textfile.pgp)
To sign a plaintext file with your secret key, and then encrypt it
with recipient's public key, producing a .pgp file:
  pgp -es textfile her_userid [other userids] [-u your_userid]
To encrypt with conventional encryption only:
  pgp -c textfile
To decrypt or check a signature for a ciphertext (.pgp) file:
  pgp ciphertextfile [-o plaintextfile]
To produce output in ASCII for email, add the -a option to other
options.
To generate your own unique public/secret key pair:  pgp -kg
For help on other key management functions, type:  pgp -k
SXSP001: /u/pgp/pgp/src: >
```

For ease of use, you should have the PGP directory to the PATH environment variable in /etc/profile:

```
export PATH=$PATH: /u/pgp/pgp/src
```

### If do not have a C/C++ compiler

If you have no C/C++ compiler, the executable PGP module is included in the src directory.

So omit the **make clean** and **make os390** commands.

## PGP online documentation

PGP documentation can be found in the directory `/u/pgp/pgp/doc/`.

## PGP KEYS MANAGEMENT

Now that PGP is installed on your z/OS system, in order to use it, you need to:

- Create your own keypair – PGP requires a keypair: a private key and a public key.
- Exchange your public keys with other PGP users.

### Create your private/public keypair

You have to create you own private/public keypair, which will be stored in two 'keyrings'.

A keyring is a file that stores keys.

Each user has two keyrings – one public and one private:

- The private keyring (`secring.pgp`) stores your private key.
- The public keyring (`pubring.pgp`) stores public keys; yours, and those that you receive from other PGP users.

To generate your own unique public/secret key pair, you should use the following command from the USS prompt:

```
SXSP001: /u/SXSP001: >pgp +nomanual -kg
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/16 10:17 GMT
```

Pick your RSA key size:

- 1) 512 bits- Low commercial grade, fast but less secure
- 2) 768 bits- High commercial grade, medium speed, good security
- 3) 1024 bits- "Military" grade, slow, highest security

Choose 1, 2, or 3, or enter desired number of bits: 2048  
2048

Generating an RSA key with a 2048-bit modulus.

You need a user ID for your public key. The desired form for this user ID is your name, followed by your E-mail address enclosed in <angle brackets>, if you have an E-mail address.

For example: John Q Smith 12345.6789@compuserve.com

Enter a user ID for your public key:

MyOrg Site1

MyOrg Site1

You need a pass phrase to protect your RSA secret key.

Your pass phrase can be any sentence or phrase and may have many words, spaces, punctuation, or any other printable characters.

Enter pass phrase: pass phrase for MyOrg Site1

Enter same pass phrase again: pass phrase for MyOrg Site1

Note that key generation is a lengthy process.

We need to generate 1536 random bits. This is done by measuring the time intervals between your keystrokes. Please enter some random text on your keyboard until you hear the beep:

1536

1528 ff

1504 fffgeez

1448 ehptytp

1376 fptoto

1320 ftppepep

1248 rprprp

1192 totptp

1136 tptptp

1080 tpttpp

1016 tpppp

984 evstcy

928 tptptp

872 toprp

824 rtpvm

776 tpbpbp

720 ptptp

672 fpfpfp

616 rprpr

568 dptptp

512 vpvvpv

456 rprprp

400 yyoy

```

360 dpdpf
312 rprp
272 dpdp
232 vpvp
192 vmvpeeer
128 vmerpe
72 cpfpf
24 epep
  * -Enough, thank you.
... **** ..... ****
Pass phrase is good. Just a moment....
Key signature certificate added.
Key generation completed.
SXSP001: /u/SXSP001: >

```

The 'pass phrase' is a kind of password to protect your secret key. Each time you need to use your secret key (to decrypt a file, to sign a message...), you will have to specify your pass phrase.

This command produces the public and the private keyrings:

```

SXSP001: /u/SXSP001: >ls -al *.pgp
-rw-rw-rw- 1 SXSP001 OMVSGRP      963 Jun 16 11:22 publi ng.pgp
-rw-rw-rw- 1 SXSP001 OMVSGRP     1489 Jun 16 11:22 secrei ng.pgp
SXSP001: /u/SXSP001: >

```

## Exchange your public keys with other PGP users

After you have created a keypair, you can begin corresponding with other PGP users.

But first, you will need to:

- Extract and send a copy of your public key to other PGP users.
- Add other PGP users' public keys to your public keyring.

### *Extracting your public key*

To extract your public key from your public ring, you should enter:

```

SXSP001: /u/SXSP001: >pgp -kx MyOrg Site1
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.

```

Current time: 2003/06/16 10:52 GMT

Extracting from key ring: 'pubring.pgp', userid "MyOrg".

Key for user ID: MyOrg Site1  
2048-bit key, key ID E59DA965, created 2003/06/16

Key extracted to file 'Site1.pgp'.  
SXSP001:/u/SXSP001: >

This command produces a copy of your public key (Site1.pgp), which can be sent (in binary) to other PGP users.

### *Adding a public key to your public keyring*

When you receive a public key from another PGP user, you have to use the following command to add it to your public keyring:

```
SXSP001:/u/SXSP001: >pgp -ka Site1.pgp
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/16 09:00 GMT
```

```
Looking for new keys...
pub 2048/E59DA965 2003/06/16 MyOrg Site1
```

```
Checking signatures...
pub 2048/E59DA965 2003/06/16 MyOrg Site1
sig! E59DA965 2003/06/16 MyOrg Site1
```

```
Keyfile contains:
  1 new key(s)
```

```
One or more of the new keys are not fully certified.
Do you want to certify any of these keys yourself (y/N)? y
y
```

```
Key for user ID: MyOrg Site1
2048-bit key, key ID E59DA965, created 2003/06/16
Key fingerprint = 70 CF C1 A4 8E F6 28 01 D8 65 4A E7 90 B5 90 EC
This key/userID association is not certified.
Questionable certification from:
MyOrg Site1
```

```
Do you want to certify this key yourself (y/N)? y
```

y

Looking for key for user 'MyOrg Site1':

Key for user ID: MyOrg Site1

2048-bit key, key ID E59DA965, created 2003/06/16

Key fingerprint = 70 CF C1 A4 8E F6 28 01 D8 65 4A E7 90 B5 90 EC

READ CAREFULLY: Based on your own direct first-hand knowledge, are you absolutely certain that you are prepared to solemnly certify that the above public key actually belongs to the user specified by the above user ID (y/N)? y

y

You need a pass phrase to unlock your RSA secret key.

Key for user ID: MyOrg Site2

512-bit key, key ID 3EFAE341, created 2003/06/11

Enter pass phrase: MyOrg Site2

Pass phrase is good. Just a moment....

Key signature certificate added.

Make a determination in your own mind whether this key actually belongs to the person whom you think it belongs to, based on available evidence. If you think it does, then based on your estimate of that person's integrity and competence in key management, answer the following question:

Would you trust "MyOrg Site1"

to act as an introducer and certify other people's public keys to you?

(1=I don't know. 2=No. 3=Usually. 4=Yes, always.) ? 4

4

SXSP001: /u/SXSP001: >

## PGP USAGE

You can use PGP to:

- Exchange an encrypted file from Site2 to Site1.
- Sign a message to send it from Site2 to Site1.

In order to use PGP, the users must have their PGP ring files in their home directory.

On Site2, we will use the following sample message file to send it to Site1:

```

BROWSE -- /u/SXSP001/text_msg.txt ----- Line 00000000 Col 001 025
Command ==> Scroll ==> PAGE
***** Top of Data *****
line 01 from MyOrg Site2
line 02 from MyOrg Site2
line 03 from MyOrg Site2
line 04 from MyOrg Site2
line 05 from MyOrg Site2
line 06 from MyOrg Site2
line 07 from MyOrg Site2
line 08 from MyOrg Site2
line 09 from MyOrg Site2
line 10 from MyOrg Site2
line 11 from MyOrg Site2
***** Bottom of Data *****

```

## Exchanging an encrypted file

### *Encrypting a file on Site2*

To encrypt a plaintext file with the recipient's Site1 public key, you should enter:

```

SXSP001:/u/SXSP001: >pgp -e text_msg.txt MyOrg Site1
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/16 09:21 GMT

```

```

Recipients' public key(s) will be used to encrypt.
Key for user ID: MyOrg Site1
2048-bit key, key ID E59DA965, created 2003/06/16

```

```

Key for user ID: MyOrg Site1
2048-bit key, key ID E59DA965, created 2003/06/16

```

```

.
Ciphertext file: text_msg.txt.pgp
SXSP001:/u/SXSP001: >

```

This command produces the encrypted file, which is not readable!

```

BROWSE -- /u/SXSP001/text_msg.txt.pgp ----- Line 00000000 Col 001 080
Command ==> Scroll ==> PAGE
***** Top of Data *****
e...a...V.z.....LL..?b..r..X.e..~/
7..E$. ;i.g.....e.....rB~..6.....eL...)J".

```



```
....P.J...//w. n. |. X.)....X....r.)..6.v..m..Z..'V..*...../...\. -..IB. -
\gs \&.eb
***** Bottom of Data *****
```

### *Decrypting a file on Site1*

To decrypt the message, the recipient on Site1 should use the following command:

```
SXSP001:/u/SXSP001: >pgp -d text_msg.txt.pgp
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/16 11:51 GMT
```

```
File is encrypted. Secret key is required to read it.
Key for user ID: MyOrg Site1
2048-bit key, key ID E59DA965, created 2003/06/16
```

```
You need a pass phrase to unlock your RSA secret key.
Enter pass phrase: pass phrase for MyOrg Site1
Pass phrase is good. Just a moment.....
Plaintext filename: text_msg.txt
SXSP001:/u/SXSP001: >
```

The result is the 'clear' file:

```
BROWSE -- /u/SXSP001/text_msg.txt ----- Line 00000000 Col 001 025
Command ==> Scroll ==> PAGE
***** Top of Data *****
line 01 from MyOrg Site2
line 02 from MyOrg Site2
line 03 from MyOrg Site2
line 04 from MyOrg Site2
line 05 from MyOrg Site2
line 06 from MyOrg Site2
line 07 from MyOrg Site2
line 08 from MyOrg Site2
line 09 from MyOrg Site2
line 10 from MyOrg Site2
line 11 from MyOrg Site2
***** Bottom of Data *****
```

### *Using PGP in batch*

You can also use PGP as a batch job using BPXBATCH:

```
//INET EXEC PGM=BPXBATCH, REGION=4096K,
// PARM='SH pgp -z ''pass phrase fo My0rg Site1'' - pass phrase
// -d /u/SXSP001/text_msg.txt.pgp'
//SYSERR DD PATH='/tmp/pgpbatch.syserr',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
//STDOUT DD PATH='/tmp/pgpbatch.stdout',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
//STDERR DD PATH='/tmp/pgpbatch.stderr',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
//SYSOUT DD PATH='/tmp/pgpbatch.sysout',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
```

This first example shows how to specify the pass phrase using the `-z` parameter.

The next example will use the `PGPPASS` environment variable:

```
//INET EXEC PGM=BPXBATCH, REGION=4096K,
// PARM='SH pgp -d /u/SXSP001/text_msg.txt.pgp'
//SYSERR DD PATH='/tmp/pgpbatch.syserr',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
//STDOUT DD PATH='/tmp/pgpbatch.stdout',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
//STDERR DD PATH='/tmp/pgpbatch.stderr',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
//SYSOUT DD PATH='/tmp/pgpbatch.sysout',
// PATHOPTS=(OWRONLY, OCREAT, OTRUNC),
// PATHMODE=SI RWXU
//STDENV DD *
PGPPASS=pass phrase fo My0rg Site1 - pass phrase
/*
```

## Signing a message

### *Signing a 'clear' message on Site2*

To sign a plaintext file with your secret key and have the output readable to people:

```
SXSP001:/u/SXSP001: >pgp -sta text_msg.txt -u My0rg Site2
No configuration file found.
```

Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.  
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18  
International version - RSA Patent Expired. Does not use RSAREF.  
Current time: 2003/06/20 15:19 GMT

A secret key is required to make a signature.  
You need a pass phrase to unlock your RSA secret key.  
Key for user ID: MyOrg Site2  
2048-bit key, key ID E59DA965, created 2003/06/16

Enter pass phrase: pass phrase for MyOrg Site2  
Pass phrase is good. Just a moment...  
Clear signature file: text\_msg.txt.asc  
SXSP001: /u/SXSP001: >

The result of this command is the following 'signed message':

```
BROWSE -- /u/SXSP001/text_msg.txt.asc ----- Line 00000000 Col 001 064  
Command ==> Scroll ==> HALF
```

```
***** Top of Data *****
```

```
----BEGIN PGP SIGNED MESSAGE----
```

```
line 01 from MyOrg Site2  
line 02 from MyOrg Site2  
line 03 from MyOrg Site2  
line 04 from MyOrg Site2  
line 05 from MyOrg Site2  
line 06 from MyOrg Site2  
line 07 from MyOrg Site2  
line 08 from MyOrg Site2  
line 09 from MyOrg Site2  
line 10 from MyOrg Site2  
line 11 from MyOrg Site2
```

```
----BEGIN PGP SIGNATURE----
```

```
Version: 2.6.3i  
Charset: noconv
```

```
i QEVAwUBPvMI 1J+cDqrI nali AQEdSgf/XK52Ji I u3z0GFVJpvnbc7MOI I fhs7XD  
b1f3VUr8G09B1/z3xHa7aRqaVePZZbdcl yrhB0tQ3oLRZtanPDV1Y9o6I EXbazhI  
xla/5XxGJzqSM9sDfj hr7TYi WC2hpyAtoYci r03HFV0088BqhLP/QZUVbzE5FNB+  
EG+bl 0kYAXaN01CVk2l 0exQYtE27DahqRxvQ7qZ0rI EmvTj 0SI 43UPJXnj H09dR2  
h51c4I EHLS4h7CPYyQvKpw1oQgVF9taUHJAcvKJI rxUzuZLD18z6N00gQrj mI BVQ  
3BDD95ZPOMKJkW5eCEwvaXzVYGXC0NkZpTGJa3VHuvFn2Rb2nSj BUg==  
=24Fu
```

```
----END PGP SIGNATURE----
```

```
***** Bottom of Data *****
```

## *Verifying the signature of the 'clear' message from Site2 on Site1*

To check the signature integrity of a signed file:

```
SXSP001:/u/SXSP001: >pgp text_msg.txt.pgp
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/20 12:43 GMT
.
File has signature. Public key is required to check signature.
.
Good signature from user "MyOrg Site2".
Signature made 2003/06/20 14:35 GMT using 2048-bit key, key ID E59DA965

Plaintext filename: text_msg.txt
SXSP001:/u/SXSP001: >
```

## *Signing an 'encrypted' message on Site2*

To sign a plaintext file with your secret key, and then encrypt it with the recipient's public key:

```
SXSP001:/u/SXSP001: >pgp -es text_msg.txt MyOrg Site1 -u MyOrg Site2
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/20 16:01 GMT

A secret key is required to make a signature.
You need a pass phrase to unlock your RSA secret key.
Key for user ID: MyOrg Site2
2048-bit key, key ID E59DA965, created 2003/06/16

Enter pass phrase: pass phrase for MyOrg Site2
Pass phrase is good. Just a moment....

Recipients' public key(s) will be used to encrypt.
Key for user ID: MyOrg Site1
2048-bit key, key ID 5A0B2D55, created 2003/06/20

Key for user ID: MyOrg Site1
2048-bit key, key ID 5A0B2D55, created 2003/06/20

Key for user ID: MyOrg Site2
2048-bit key, key ID E59DA965, created 2003/06/16
.
```

Ciphertext file: text\_msg.txt.pgp  
SXSP001: /u/SXSP001: >

### *Verifying the signature of the 'encrypted' message from Site2 on Site1*

To decrypt an encrypted file and to check the signature integrity of a signed file:

```
SXSP001: /u/SXSP001: >pgp text_msg.txt.pgp
No configuration file found.
Pretty Good Privacy(tm) 2.6.3i - Public-key encryption for the masses.
(c) 1990-96 Philip Zimmermann, Phil's Pretty Good Software. 1996-01-18
International version - RSA Patent Expired. Does not use RSAREF.
Current time: 2003/06/20 14:03 GMT
```

```
File is encrypted. Secret key is required to read it.
Key for user ID: MyOrg Site1
2048-bit key, key ID 5A0B2D55, created 2003/06/20
```

```
You need a pass phrase to unlock your RSA secret key.
Enter pass phrase: pass phrase for MyOrg Site1
Pass phrase is good. Just a moment.....
File has signature. Public key is required to check signature.
```

```
.
Good signature from user "MyOrg Site2".
Signature made 2003/06/20 16:01 GMT using 2048-bit key, key ID E59DA965
```

```
Plaintext filename: text_msg.txt
SXSP001: /u/SXSP001: >
```

## PGP LINKS

Alan Nichols' OS/390 PGP page: <http://s390.nichols.de/pgp/index.html>.

International PGP Home page: <http://www.pgpi.org/>.

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## Using TAR and JAR files on MVS

I recently needed to transmit a large number of TIF images (held in MVS datasets) to a server running a Windows environment. The Unix TAR command can be used under Unix System Services on MVS to create an archive file that can then be transmitted to the server and expanded using the familiar Windows ZIP utility.

In the following example the images are first copied from a PDS to an HFS directory using the OPUTX EXEC.

```
/**
/** use the OPUTX exec to copy each member of the
/** PDS into the HFS directory
/**
//OPUTX EXEC PGM=IKJEFT01
//SYSEXEC DD DSN=SYS1.SBPXEXEC,DISP=SHR
//SYSPROC DD DISP=SHR,DSN=SYS1.SISPLIB
//ISPLIB DD DISP=SHR,DSN=SYS1.SISPLIB
//          DD DISP=SHR,DSN=SYS1.SISPLIB
//ISPLIB DD DISP=SHR,DSN=SYS1.SISPLIB
//          DD DISP=SHR,DSN=SYS1.SISPLIB
//ISPLIB DD DISP=SHR,DSN=SYS1.SISPLIB
//          DD DISP=SHR,DSN=SYS1.SISPLIB
//ISPLIB DD DISP=SHR,DSN=SYS1.SISPLIB
//          DD DISP=SHR,DSN=SYS1.SISPLIB
//ISPLIB DD DISP=SHR,DSN=SYS1.SISPLIB
//          DD DISP=SHR,DSN=SYS1.SISPLIB
//ISPLIB DD DSN=&&ISPLIB,SPACE=(TRK,(5,1,2)),LIKE=&userid.ISPLIB,
//          DISP=(NEW,PASS)
//SYSTSPRT DD SYSOUT=*
//SYSTSIN DD *
        OPUTX 'your-PDS-of-images' '/your-dir' ASIS BINARY CONVERT(NO) +
        MODE(644) SUFFIX(tif)
/**
```

The TAR command is then used to create the archive dataset – note that in this particular case the compression flag is not used because TIF images are already in a compressed format.

```
/**
/** use Unix System Services TAR command
/** to create an archive dataset
/**
//TAR EXEC PGM=BPXBATCH,
```

```

//          PARM='sh tar -cUvf //tiffs.tar /your-dir'
//STDOOUT DD PATH='/tmp/stdout' ,
//          PATHOPTS=(OCREAT, OWRONLY), PATHMODE=SI RWXU,
//          PATHDI SP=KEEP
//STDERT DD PATH='/tmp/stderr' ,
//          PATHOPTS=(OCREAT, OWRONLY), PATHMODE=SI RWXU,
//          PATHDI SP=KEEP
//*
```

The resulting MVS dataset, &userid.TIFFS.TAR, can be sent (obviously as a binary transmit) to the server and unzipped to restore the individual image files. Note that the Windows ZIP program will recognize files with an extension of .tar.

If you want to include text files in your archive, there is a slight problem in that we are using a binary transfer – somewhere along the line the data will need to be translated into ASCII. The following job will take an MVS text file (ie EBCDIC character data) and convert it to ASCII. The TEXT option in this case will result in the addition of an LF character being added to delimit the end of each record. This will be OK for a Unix system; however, the standard Windows text file expects CR LF or X'0D0A' as the delimiter.

```

//OCOPY EXEC PGM=IKJEFT01
//I NDD DD DSN=your-text-file, DI SP=OLD
//OUTHFS DD PATH='/your-dir/myfile.txt' ,
//          PATHDI SP=(KEEP, DELETE),
//          PATHOPTS=(OWRONLY, OCREAT), PATHMODE=SI RWXU
//SYSTSPRT DD SYSOUT=*
//SYSTSIN DD *
OCOPY INDD(INDD) OUTDD(OUTHFS) TEXT CONVERT((BPXFX311)) +
FROM1047
/*
```

While the Windows ZIP utility is quite happy to work with Unix TAR files created on MVS, the reverse is not true. To cater for the situation where you may want to package up some files on a Windows environment and upload them to MVS as a single dataset, you can make use of the Java archive utility, which provides a similar ZIP functionality, and, being Java, provides platform-independence. Further information about Java archive files can be found on the Internet (see the Sun Java tutorial pages).

Assuming that you have made Java available on your MVS system (mounting the HFS containing the Java software and setting the environment variables), the following JCL illustrates the use of the JAR command to expand a Java archive file that was created on a Windows PC:

```
//* use JAR to unload the archive
//*
//JAR      EXEC PGM=BPXBATCH,
//        PARM='sh cd work;jar xf /work/fixes.jar'
//STDOUT  DD PATH='/tmp/stdout',
//        PATHOPTS=(O_CREAT, O_WRONLY), PATHMODE=SI_RWXU,
//        PATHDISP=KEEP
//STDERR  DD PATH='/tmp/stderr',
//        PATHOPTS=(O_CREAT, O_WRONLY), PATHMODE=SI_RWXU,
//        PATHDISP=KEEP
//*
```

The OGETX EXEC can then be used to move files into a standard MVS PDS:

```
OGETX /work '&userid.MVS.PDS' LC SUFFIX(bin) BINARY
```

With many vendors providing fixes over the Internet as downloads, one possible use for this kind of job would be to move a number of fixes from a directory on your PC up to an MVS dataset.

---

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## **Finding CSECTs within LPA load modules in virtual storage**

In the course of development of certain types of system application, the ability to programmatically locate CSECT addresses within LPA-resident modules can sometimes prove useful. Unfortunately, the information to facilitate such location, ie the displacement of a given CSECT within the load module, while present in the linkage editor control information for the module, is not retained in virtual storage when the module is



actually loaded. Thus, it becomes necessary to refer back to the linkage editor CESD (Composite External Symbol Dictionary) control records, which form a portion of the load module's contents in its LPALST library on DASD. From these, the desired CSECT can be located, and its load module displacement extracted. This value, when added to the output of a system service call (which identifies the module's virtual storage load address), precisely identifies the virtual storage address of the aforesaid CSECT.

FNDCSCT is a statically or dynamically called routine which performs the above function. It is passed the load module and CSECT names as parameters, and, upon completion of a successful location operation, returns the associated virtual storage address in register 0 and a return code of 0 in register 15. WTO-type error messages and a non-zero return code in register 15 are generated upon recognition of any conditions which preclude successful completion. Note that the location technique employed will result in the appearance of one or more IEC141I 013-18 exception messages if the desired module is not found in the first of two or more LPALST libraries. These do not signal error conditions if the module is found in a subsequent LPALST library. If the module is not found in any LPALST library, then IEC141I 013-18 messages will appear for all such libraries, and a further 'FNDCSCT007E Load module loadmodulename was not found' message will signal the error condition.

FNDCSCT should be linkedited as a stand-alone load module into any desired load library. No specific linkage editor attributes need be assigned.

The calling sequences for FNDCSCT are as follows:

Static call:

```

          call FNDCSCT, (loadmod, csect)
          ltr  r15, r15
          bnz  errorrtn
loadmod  dc   cl 8' loadmodulename'
csect    dc   cl 8' Csectname'

```

## Dynamic call:

```
link    ep=FNDCSCT, (loadmod, csect)
ltr     r15, r15
bnz     errorrtn
loadmod dc    c1 8' loadmodul ename'
csect   dc    c1 8' Csectname'
```

## FNDCSCT

```
fndcsct amode 31
fndcsct rmode 24
fndcsct csect
r0      equ    0
r1      equ    1
r2      equ    2
r3      equ    3
r4      equ    4
r5      equ    5
r6      equ    6
r7      equ    7
r8      equ    8
r9      equ    9
r10     equ   10
r11     equ   11
r12     equ   12
r13     equ   13
r14     equ   14
r15     equ   15
        stm    r14, r12, 12(r13)           entry linkage
        lr     r12, r15
        using fndcsct, r12
        st    r13, savearea+4
        la    r15, savearea
        st    r15, 8(r13)
        lr     r13, r15
        b     fcsc0020                     branch to start
return  ds    0h
        xr    r15, r15                     exit linkage
        l     r13, 4(r13)
        l     r14, 12(r13)
        lm   r1, r12, 24(r13)
        br    r14                          return
fcsc0020 ds    0h
        lr     r9, r1                      save pal pointer
        l     r3, cvtptr                  CVT pointer
        icm   r3, 15, cvtsmext-cvt(r3)    CVT extension
        bnz   fcsc0050                     it's there
```

```

wto 'FNDCSCT001E No CVTX address found'
la r15,8
b return+2
fcsc0050 ds 0h
icm r3,15,cvteplps-cvtvstgx(r3) LPAT address
bnz fcsc0100 it's there
wto 'FNDCSCT002E No LPAT address found'
la r15,8
b return+2
fcsc0100 ds 0h
cli =cl 4' LPAT', 0(r3) really the
LPAT? 00490000
be fcsc0150 YES
wto 'FNDCSCT003E LPAT ID check failed'
la r15,8
b return+2
fcsc0150 ds 0h
icm r4,15,4(r3) number of LPAT entries
bnz fcsc0200 more than none
wto 'FNDCSCT004E No LPAT entries present'
la r15,8
b return+2
fcsc0200 ds 0h
l r2,0(r9) get member name address
mvc membrtu+6(8),0(r2) move member name to txt unit
la r3,9(r3) point to first LPAT entry
fcsc0250 ds 0h
cli 0(r3),0 end of entries?
be fcsc0900 yes, module not found
mvc dsnamtu+6(44),0(r3) no, move dsname to text unit
mvi alcverb,s99vrbal prime dynaloc fields
la r0,dsnamtu
st r0,alctua1
la r0,membrtu
st r0,alctua2
la r0,statstu
st r0,alctua3
la r0,rtdntu
st r0,alctua4
oi alctua4,x'80'
la r1,alcrbptr
dynaloc
ltr r15,r15 allocation ok?
bz fcsc0400 yes
st r15,allocrc
mvc wto00250+37(44),dsnamtu+6 move dataset name
trt wto00250+37(44),trtable search for terminating blank
mvi 0(r1),c'(' member name preparation
lr r5,r1 retain address of blank

```

```

        mvc    1(8, r5), membrtu+6           move member name
        trt    1(8, r5), trtable           search for terminating blank
        mvi    0(r1), 'c')'               end of member name
        cnop   0, 4
wto00250 wto  'FNDCSCT005E Error allocating
,
        mvc    dfdaplp, =a(alcrb)          initialize DAIRFAIL parms
        mvc    dfrcp, =a(allocrc)
        la     r1, =a(0)
        st     r1, dfj eff02
        la     r1, =x'4032'
        st     r1, dfi dp
        xc     dfcpplp, dfcpplp
        mvc    dfbufp, =a(dfbufs)
        la     r1, dfparms
        link   ep=IKJEFF18                invoke DAIRFAIL
        ltr    r2, r15                    ok?
        bz     fcsc0300                    yes
wto      'FNDCSCT006E DAIRFAIL error - return code set to DAIRFAI*
        L return code'
        lr     r15, r2
        b      return+2
fcsc0300 ds    0h
        lh     r5, dfbufl 1                extract the DAIRFAIL message
        sh     r5, =h'5'                   set/check message length
        ch     r5, =h'112'
        bnh    fcsc0320
        lh     r5, =h'112'
fcsc0320 ds    0h
        ex     r5, exmvc1                  move it to the wto
        cnop   0, 4
wto00300
wto      'FNDCSCT006E
,
,
        issue it
        clc    dfbufl 2, =a(0)            any second level message?
        be     fcsc0340                    no
        lh     r5, dfbufl 2                yes, extract as well
        sh     r5, =h'5'                   set/check message length
        ch     r5, =h'112'
        bnh    fcsc0330
        lh     r5, =h'112'
fcsc0330 ds    0h
        ex     r5, exmvc2                  move it to wto
        cnop   0, 4
wto00330
wto      'FNDCSCT006E
,
,

```

fcsc0340	ds	0h	
	l	r15, allocrc	
	b	return+2	
fcsc0400	ds	0h	
	mvc	library+dcbddnam-i hadcb(8), rtdntu+6	move ddname
	open	(library, (INPUT)), mode=31	open the library
	tm	library+(dcboflgs-i hadcb), dcbofopn	ok?
	bo	fcsc1000	yes, module has been located
	la	r3, 45(, r3)	no, next LPAT library
	b	fcsc0250	recycle
fcsc0900	ds	0h	
	l	r3, 0(, r9)	get load module name addr
	mvc	wto00900+32(8), 0(r3)	move load module name
	cnop	0, 4	
wto00900	wto	'FNDCSCT007E Load module xxxxxxxx was not found'	
	la	r15, 8	
	b	return+2	
fcsc1000	ds	0h	
	get	library	get a load module record
	lr	r3, r1	retain its address
	cli	0(r3), x' 20'	CESD record?
	bne	fcsc1000	no, get another one
	la	r4, 0	
	icm	r4, 3, 6(r3)	get record length
	la	r3, 8(, r3)	
	la	r4, 0(r3, r4)	point past last byte
	sh	r4, =h' 8'	back off sufficiently
	l	r5, 4(, r9)	get passed CSECT name addr
fcsc1150	ds	0h	
	clr	r3, r4	past upper search limit?
	bni	fcsc1000	yes, go get another record
	clc	0(8, r3), 0(r5)	no, CSECT names match?
	bne	fcsc1500	no
	tm	8(r3), x' 0f'	yes, type=SD?
	bz	fcsc2000	yes, what we're looking for
fcsc1500	ds	0h	
	la	r3, 1(, r3)	next byte
	b	fcsc1150	iterate
fcsc2000	ds	0h	
	la	r4, 0	
	icm	r4, b' 0111', 9(r3)	load the displacement
	l	r8, 0(, r9)	point to load module name
	csvquery	inename=(r8), search=LPA, outloadpt=loadpt	look for it
	ltr	r2, r15	find it?
	bz	fcsc2050	yes
	ch	r2, =h' 8'	module not found?
	be	fcsc2020	yes
	wto	'FNDCSCT008E CSVQUERY error - return code set to CSVQUER* Y return code'	

```

        lr    r15, r2
        b     return+2
fcsc2020 ds    0h
        wto   'FNDCSCT009E CSVQUERY could not find the requested load*
            module'
        la    r15, 8
        b     return+2
fcsc2050 ds    0h
        l     r0, loadpt          Load module load point
        alr   r0, r4             add CSECT
displacement      01960013
        b     return            return
*****
* DCB ABEND exit *
*****
fcsc8000 ds    0h
        lr    r3, r1             retain parameter list addr
        lr    r4, r14           retain return addr
        mvi   3(r3), 4         ignore the abend
        lr    r14, R4          restore return addr
        br    r14              return
*****
* DCB EODAD routine *
*****
fcsc9000 ds    0h
        l     r2, 4(r9)         Load CSECT name address
        mvc   wto09000+51(8), 0(r2)  move CSECT name
        cnop  0, 4
wto09000 wto   'FNDCSCT010E No CESD record found for CSECT xxxxxxxx '
        la    r15, 8
        b     return+2
*****
* Executed instructions *
*****
exmvc1  mvc   wto00300+20(0), dfbuft1
exmvc2  mvc   wto00330+20(0), dfbuft2
*****
* Data Area *
*****
savearea dc    18f' 0'
dsnamtu  dc    al 2(dal dsnam), al 2(1), al 2(44), cl 44' '
membrtu  dc    al 2(dal membr), al 2(1), al 2(08), cl 08' '
statstu  dc    al 2(dal stats), al 2(1), al 2(01), xl 01' 8'
rtddntu  dc    al 2(dal rtddn), al 2(1), al 2(08), cl 08' '
library  dcb   ddname=dummy, macrf=GL, dsorg=PS, recfm=U, lrecl=0, *
            blksi ze=32760, devd=DA, eodad=fcsc9000, exlst=exlst
exlst    dc    0f' 0', x' 11', al 3(fcsc8000)
loadpt   dc    a(0)            module load point
allocrc  dc    f' 0'          allocation return code

```

```

al crbptr dc      x' 80' , al 3(al crb)      request block pointer
al crb     ds      0f                          request block
al crbln   dc      al 1(20)                   request block length
al cVERB   dc      al 1(0)                     verb code
al cFLAG1  ds      0al 2                       flags
al cflg11  dc      al 1(0)                     first flags byte
al cflg12  dc      al 1(0)                     second flags byte
al crsc    ds      0al 4                       reason code fields
al cerror  dc      al 2(0)                     error reason code
al cinfo   dc      al 2(0)                     information reason code
al cxtpp   dc      a(al ctupl)                 tupl address
al crsv01  ds      f                           reserved
al cflg2   ds      0al 4                       authorized functions flags
al cflg21  dc      al 1(0)                     first flags byte
al cflg22  dc      al 1(0)                     second flags byte
al cflg23  dc      al 1(0)                     third flags byte
al cflg24  dc      al 1(0)                     fourth flags byte
al ctupl   ds      0f                          text unit pointer list
al ctua1   dc      a(0)                        text unit address 1
al ctua2   dc      a(0)                        text unit address 2
al ctua3   dc      a(0)                        text unit address 3
al ctua4   dc      a(0)                        text unit address 4
          i k j e f f d f
trtable   dc      256x' 0'
          org      trtable+c' '
          dc      c' '
          org      ,
*****
* Dsects *
*****
          cvt      dsect=YES, list=NO
          dcbd     dsorg=PS, devd=DA
          iefzb4d0
          iefzb4d2
          end

```

---

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---

## Reorganization of datasets

KSDS datasets with a lot of activity tend to get split. In practice, CI splits are more frequent than CA splits, and when a CA split happens the number of CI splits is normally already large.

Datasets that tend to get split should be periodically reorganized, in order to avoid performance degradation.

The program presented here reads a LISTCAT output, searching for KSDS files whose number of CI splits is greater than a specified limit. When it finds one, it writes out a five-step JCL to reorganize that file.

The steps are as follows:

- Create VSAM temp file (similar to the original)
- Repro original to the temp
- Delete and define the original file
- Repro from temp to the new file
- Delete the temp.

The volume where the new file is allocated is the same as the one where the original file was. I made no provision for multi-volume files because I felt files with such dimensions should be treated separately. This program should be used only for single volume datasets.

All the JCL is written to a single file that can later be submitted, or modified if necessary before being submitted. The program also writes to SYSTSPRT a list of the files considered for reorganization, showing how many there are.

This EXEC should be run in a two-step batch job, where the first step issues a LISTCAT ALL for a given catalog, writing the output to a file, and the second step runs the EXEC, reading the output and writing out the JCL to another file.

## SAMPLE JOB

```
//STEP1      EXEC PGM=IDCAMS
//SYSPRINT DD DI SP=(NEW, PASS), DSN=&&TEMP1, UNIT=VI O,
//           DCB=(LRECL=121, RECFM=FB, BLKSIZE=2420),
//           SPACE=(TRK, (30, 30))
//SYSOUT DD SYSOUT=*
//SYSIN DD *
```



```

LISTCAT ALL CATALOG(my.catalog)
/*
//STEP2 EXEC PGM=IRXJCL, PARM='VSREORG'
//SYSEXEC DD DISP=SHR, DSN=where.is.the.exec
//SYSTSPRT DD SYSOUT=*
//SYSPRINT DD SYSOUT=*
//FILEIN DD DISP=(OLD,DELETE), DSN=&&TEMP1
//FILEOUT DD DISP=SHR, DSN=jcl.file

```

## VSREORG REXX

```

/*= REXX =====*/
/* VSREORG - Reads a "LISTCAT ALL" listing allocated to FILEIN and */
/* generates JCL for the reorganization of KSDS datasets */
/* whose CI split number is greater than a certain */
/* threshold. The JCL is written to DDname FILEOUT. */
/* Threshold value is indicated by variable splitsci. */
/*=====*/
splitsci = 10 /* minimum number of CI splits for reorg */
suf = "TP" /* suffix for temporary datasets */
execio 1 disk filein
if rc <> 0 then do
  say "Error reading filein"
  exit
end
pull linha
cc = left(linha, 1)
data_flag = 0
c = 0
do alpha = 0
  execio 1 disk filein
  if rc <> 0 then leave
  pull linha
  if cc = 1 then linha = substr(linha, 2)
  call check_line_type
  if data_flag = 1 then do
    call extract_values
  end
end
end
call write_fileout
say
say total "datasets were considered for reorganization"
exit
/*=====*/
/* search for a DATA component */
/*=====*/
check_line_type:
select

```

```

when substr(linha, 1, 7)="CLUSTER" then do
  c = c + 1
  cluster.c = word(linha, 3)
  data_flag = 0
  aix = 0
end
when substr(linha, 4, 4)="DATA" & aix = 0 then do
  data.c = word(linha, 3)
  data_flag = 1
end
when substr(linha, 1, 3)="AIX" then do
  data_flag = 0
  aix = 1
end
when substr(linha, 4, 5)="INDEX" & data_flag=1 & aix=0 then do
  index.c = word(linha, 3)
  data_flag = 0
end
when substr(linha, 1, 7)="NONVSAM" then do
  data_flag = 0
end
otherwise nop
end
return
/*=====*/
/*          extract values from line          */
/*=====*/
extract_values:
select
  when substr(linha, 8, 7)="REC-TOT" & data_flag = 1 then do
    linha = translate(linha, " ", "-")
    rectot.c = right(word(linha, 3), 11)
    splici.c = right(word(linha, 6), 5)
    if splici.c < splitsci then do
      c = c - 1
      data_flag = 0
    end
  end
  when substr(linha, 8, 8)="SHROPTNS" & data_flag = 1 then do
    shropt.c = substr(linha, 17, 3)
    if word(linha, 5) <> "INDEXED" then do
      c = c - 1
      data_flag = 0
    end
  end
  when substr(linha, 8, 7)="REC-INS" then do
    linha = translate(linha, " ", "-")
    freeci.c = right(word(linha, 6), 2)
  end
end

```

```

when substr(linha,8,7)="REC-UPD" then do
  linha = translate(linha, " ", "-")
  freeca.c = right(word(linha,6),2)
end
when substr(linha,8,7)="REC-DEL" then do
  linha = translate(linha, " ", "-")
  spli ca.c = right(word(linha,6),3)
end
when substr(linha,8,6)="KEYLEN" then do
  linha = translate(linha, " ", "-")
  keylen.c = right(word(linha,2),2)
  alrecl.c = right(word(linha,4),5)
  alrecl.c = space(alrecl.c,0)
  cisi ze.c = right(word(linha,8),5)
  cisi ze.c = space(cisi ze.c,0)
end
when substr(linha,8,3)="RKP" then do
  linha = translate(linha, " ", "-")
  keypos.c = right(word(linha,2),2)
  mlrecl.c = right(word(linha,4),5)
  mlrecl.c = space(mlrecl.c,0)
end
when substr(linha,8,7)="SPACE-T" then do
  linha = translate(linha, " ", "-")
  sptype.c = left(word(linha,3),3)
  if sptype.c = "CYL" then sptype.c = "CYLINDERS"
  if sptype.c = "TRA" then sptype.c = "TRACKS"
  if sptype.c = "REC" then sptype.c = "RECORDS"
end
when substr(linha,8,7)="SPACE-P" then do
  linha = translate(linha, " ", "-")
  spprim.c = right(word(linha,3),5)
  spprim.c = space(spprim.c,0)
end
when substr(linha,8,7)="SPACE-S" then do
  linha = translate(linha, " ", "-")
  spseco.c = right(word(linha,3),4)
  spseco.c = space(spseco.c,0)
end
when substr(linha,8,6)="VOLSER" then do
  linha = translate(linha, " ", "-")
  volume.c = word(linha,2)
end
otherwise nop
end
return
/*=====*/
/*                Write fileout                */
/*=====*/

```

```

write_fileout:
total = 0
do k = 1 to c
  dropbuf
  if length(cluster.k". "suf) > 42 then do
    say ">>> Temporary name exceeds 42 characteres"
    say ">>> for dataset " cluster.k". "suf
    say ">>> Dataset skipped"
    iterate k
  end
  queue "//STEP" k "1 EXEC PGM=IDCAMS, COND=(0, LT)"
  queue "//SYSPRINT DD SYSOUT=*"
  queue "//SYSIN DD *"
  queue " DEFINE CLUSTER( -"
  queue "   NAME("cluster.k". "suf") -"
  queue "           "sptype.k"("sppri m.k spseco.k") -"
  queue "           RECSZ("al recl . k ml recl . k") -"
  queue "           KEYS("keyl en. k keypos. k")) -"
  queue "   DATA( -"
  queue "           NAME("cluster.k". "suf". D)) -"
  queue "   INDEX( -"
  queue "           NAME("cluster.k". "suf". I))"
  queue "/*"
  queue "//STEP" k "2 EXEC PGM=IDCAMS, COND=(0, LT)"
  queue "//INF DD DISP=SHR, "
  queue "// DSN="cluster.k
  queue "//OUTF DD DISP=SHR, "
  queue "// DSN="cluster.k". "suf""
  queue "//SYSPRINT DD SYSOUT=*"
  queue "//SYSIN DD *"
  queue " REPRO INFILE(INF) OUTFILE(OUTF)"
  queue "/*"
  queue "//STEP" k "3 EXEC PGM=IDCAMS, COND=(0, LT)"
  queue "//SYSPRINT DD SYSOUT=*"
  queue "//SYSIN DD *"
  queue " DELETE ("cluster.k") CL PURGE"
  queue " DEFINE CLUSTER( -"
  queue "   NAME("cluster.k") -"
  queue "           "sptype.k"("sppri m.k spseco.k") -"
  queue "           VOLUME("vol ume. k") -"
  queue "           RECSZ("al recl . k ml recl . k") -"
  queue "           FSPC("freeci . k freeca. k") -"
  queue "           KEYS("keyl en. k keypos. k")) -"
  queue "   DATA( -"
  queue "           NAME("data. k") -"
  queue "           CI SZ("ci si ze. k")) -"
  queue "   INDEX( -"
  queue "           NAME("i ndex. k"))"
  queue "/*"

```

```

queue "//STEP"K"4 EXEC PGM=IDCAMS, COND=(Ø, LT) "
queue "//INF DD DISP=SHR, "
queue "// DSN="cluster.k". "suf""
queue "//OUTF DD DISP=SHR, "
queue "// DSN="cluster.k
queue "//SYSPRINT DD SYSOUT=*"
queue "//SYSIN DD *"
queue " REPRO INFILE(INF) OUTFILE(OUTF)"
queue "/*"
queue "//STEP"K"5 EXEC PGM=IDCAMS, COND=(Ø, LT) "
queue "//SYSPRINT DD SYSOUT=*"
queue "//SYSIN DD *"
queue " DELETE ("cluster.k". "suf") CL PURGE"
queue "/*"
queue "/*"
queue "/*=====*"
queue "/*"
execio queued() disk fileout
total = total + 1
say "JCL generated for dataset " cluster.k
end
execio Ø disk fileout "(finis"
return

```

---

*Systems Engineer*  
(Portugal)

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---

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Information Builders, Two Penn Plaza, New York, NY 10121-2898, USA.  
Tel: (212) 736 4433.  
URL: <http://www.informationbuilders.com/products/webfocus/index.html>.

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IIBM has announced WebSphere Business Integration Adapters V2.3.0 and V2.3.1,

which introduce 11 new adapters and provide connectivity to WebSphere Application Server.

The WebSphere Business Integration Adapter portfolio has been expanded to support industry applications. This release contains new adapters to provide connectivity to IndusConnect Framework, Maximo MEA, and QAD MFG/PRO.

New technology adapters in this release provide support for healthcare industry standards, such as Health Level Seven (HL7), as well as providing connectivity to iSeries, Lotus Domino, Exchange, COM, and CORBA.

Two new mainframe adapters are also being introduced. The first provides connectivity to a Natural application by executing the Natural program and retrieving the data returned from the execution. The adapter uses JDBC connectivity to the OS/390 platform. Standard operations on the applications are executed using stored procedure calls. The second is the adapter for IDMS, which provides access to IDMS data on mainframe systems using JDBC connectivity and host RPC. Standard operations on the applications are executed using stored procedure calls.

Both mainframe adapters use the NEON Shadow client and server to achieve connectivity. These adapters are available on AIX, HP-UX, Solaris, and Windows.

For further information contact your local IBM representative.  
URL: <http://www.ibm.com/websphere/integration/wbiadapters>.

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**xephon**