



# 224

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# update

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# ***MVS Update***

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## Considerations when placing page datasets on DASD with PAV enabled

We recently experienced some issues with page dataset usage and on investigation discovered some interesting information.

The initial problem manifested itself as high auxiliary storage usage; however, on issuing a **D ASM** from the operator console we noticed that 10 of our 20 page datasets were actually 0% utilized. This obviously confused us, but on doing a little research we identified the devices that were not being used – they were non-PAV type devices. The datasets we had placed on PAV devices were being over used.

We found an APAR OA04644, which described our problem. What IBM has done is recode the Auxiliary Storage Manager from z/OS 1.3 onwards. The recoding allows a separate queue to be created for PARTE control blocks that represent page datasets residing on Parallel Access Volumes. This works as follows: if a page out occurs, the Auxiliary Storage Manager will make a search for which page dataset to place the page onto. First of all it searches the PAV device queue, before looking at other device types. If a PAV device is located with a response time of less than the average of all other device types, it will be chosen as the candidate for the page-out. This results in PAV devices getting more pages than non-PAV devices. This, according to IBM, is intentional.

The actual implementation of the separate queue for PAV devices introduces a potential risk to systems that page a lot. We overcame the problem by moving the page datasets on the non-PAV volumes to PAV volumes, and everything is balanced again. However, if you are not aware of this change, you may find that your paging environment is not acting the way you expect it to.

IBM has actually now updated its documentation to reflect the

fact that any page datasets residing on PAV devices will be examined first when Auxiliary Storage Manager is searching for a page dataset to page out to. The bottom line is that when the Auxiliary Storage Manager selects a dataset, any paging datasets that reside on Parallel Access Volume devices are examined first because of the reliability and performance characteristics of these devices.

Because of this, preference will be given to PAV devices – and it is normal to see higher utilization on the PAV-resident datasets when compared with non-PAV resident datasets.

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## **Deleted datasets at a glance**

### **PROBLEM ADDRESSED**

Every now and then it happens that someone from a site's large user base asks the 'perennial' question: "Where is my dataset? It used to be in that catalog, on that volume, but it is not there any more. Who deleted it?"

During the course of recent reviewing and restructuring of DR procedures, the very same question occurred more frequently than usual, and prompted me to search for a quick, simple, and easy-to-use solution that would supply the straightforward answer.

### **SOLUTION PROPOSED**

In a search for a solution I have looked at the SMF manual to find out whether there is any record that would allow me to determine who deleted or renamed a dataset. It came as no

surprise that actually there are such SMF records. At this point, I would suggest that you first check whether you are collecting these SMF records at all. They are fairly voluminous, and lots of installations don't bother. When enabled by the SMFPRMxx TYPE parameter, SMF creates type 17 and 18 records, which are written each time a dataset is deleted or renamed. Record type 17 is written when a non-temporary DASD dataset or a temporary DASD dataset is deleted/scratched. This record contains the dataset name, number of volumes, and volume serial numbers. The job name, time, and date that the reader recognized the JOB card (for this job) constitute the job's identification.

However, you should note that you will not see the userid (or RACF ID) in the type 17 records of the user who actually deleted the dataset. You may get tempted to look for a RACF ID at the supplied field SMF17UID, but that field is not supposed to be the RACF ID. In fact, the SMF manual states that the content of this field is taken from the SMF Common Exit Parameter Area, not from the USER= parameter on the job statement. In particular it is copied from the 8 bytes starting at offset 20. This area is initialized to blanks and may be updated by SMF exits such as IEFUJI, IEFUJV, IEFUSI, and IEFACTRT. The SMF CEPA is mapped by the JMR (Job Management Record), which is mapped by the IEFJMR macro. If you want this to contain the RACF ID of the job, it would probably be easiest to write an IEFUJI exit to take the ID from the ACEE and put it into the field in the JMR. You can find the type 17 description in macro IGGSMF17 in SYS1.MACLIB.

On the other hand, it may happen that SMF type 17 records are not being produced during a scratch DASD dataset because the check on the dataset name in IGG0290D was not consistent with the SMF manual's description of a TEMPORARY or SYSTEM-generated dataset name. In other words, DADSM was interpreting a non-temporary dataset name as a temporary dataset name: IGG0290D was checking for 'SYS' in positions 1-3, '.T' in positions 9-10, and '.' in

position 17. Thus, a dataset name like SYSX.SMF.TYPE110.\* was considered when writing SMF records during a scratch. If you experience that kind of problem, take a look at the APAR database as to find a PTF for APAR OW43581, which corrects the SMF type 17 fields. Installing APAR OW41664 corrupts the SMF type 17 fields JOB, READTIME, and LOCLINFO, and the DSNNAME has 'SYS' in columns 1–3. New APAR OW43581 corrects the problem created when the strengthening of DADSM's name checking by OW41664 destroyed the register used to build the type 17 SMF record (dataset SCRATCH).

Also please note that you may use the REC parameter, in the SMFPRMxx parmlib member, to specify that record type 17 is to be collected. REC(ALL) specifies that record type 17 is to be written for both temporary and non-temporary datasets. REC(PERM) specifies that SMF is to write record type 17 only for non-temporary datasets; it is not to be written for temporary datasets (ie datasets having names that start with SYSydd. Thhmmss, either from DSN=&&datasetname or from the absence of any dataset name).

It may happen that a particular dataset has not been deleted at all, but renamed. SMF writes record type 18 each time a non-VSAM dataset defined by a DD statement (either explicitly or implicitly) is renamed. This record contains the old dataset name, the new dataset name, the number of volumes, and the volume serial numbers.

A detailed description of the layout of SMF type 17 and 18 records can be obtained from the manual *MVS System Management Facilities (SMF)* – SA22-7630-03.

## CODE

Based on the record description obtained from the manual mentioned above, a simple report writer was written. The code is a two-part stream. In the first part (DUMP178) SMF records 17 and 18 are dumped/copied from SMF datasets to a file, which can be used as a base of archived records.

In the second part (DS1718), the captured records are formatted and four reports are produced. These reports are grouped into two sets: one set reports all deleted datasets, while the second set provides a summary of renamed datasets.

The dataset deletion report set displays the job or system task performing the deletion, along with the dataset name, volume serial of the dataset, as well as the date and time the deletion took place.

The dataset rename report set displays the job or system task performing the renaming, along with the old dataset name, new dataset name, volume serial of the dataset, and the date and time the rename took place.

```

/*-----
/* The SMF extract job.
/* This job will extract the SMF records type 17 & 18 from
/* selected SMF dataset (weekly or current)
/*-----
//DUMP178 EXEC PGM=IFASMFDP,REGION=0M
//INDD DD DSN=your.smf.ds,DISP=SHR
//OUTDD DD DSN=&&SMF1718,DISP=(NEW,PASS),
// SPACE=(CYL,(25,5)),DCB=(your.smf.weekly.ds)
//SYSPRINT DD SYSOUT=*
//SYSIN DD *
        INDD(INDD,OPTIONS(DUMP))
        OUTDD(OUTDD,TYPE(17,18))
/*
//DS1718 EXEC PGM=ICETOOL
//TOOLMSG DD SYSOUT=*
//DFSMSG DD SYSOUT=*
//RAWSMF DD DSN=&&SMF1718,DISP=SHR
//SMF1718 DD DSN=&&TEMPA,SPACE=(CYL,(9,9)),UNIT=SYSDA
//SMF17 DD DSN=&&TEMP7,SPACE=(CYL,(9,9)),UNIT=SYSDA
//SMF18 DD DSN=&&TEMP8,SPACE=(CYL,(9,9)),UNIT=SYSDA
//SYS17 DD DSN=&&TEMPS,SPACE=(CYL,(9,9)),UNIT=SYSDA
//SYS18 DD DSN=&&TEMPT,SPACE=(CYL,(9,9)),UNIT=SYSDA
//JOBDEL DD SYSOUT=*
//JOBREN DD SYSOUT=*
//SYSDEL DD SYSOUT=*
//SYSREN DD SYSOUT=*
//TOOLIN DD *
COPY FROM(RAWSMF) TO(SMF1718) USING(SMFI)
COPY FROM(SMF1718) TO(SMF17) USING(SMF7)
DISPLAY FROM(SMF17) LIST(JOBDEL) -
        TITLE('Job deleted datasets') DATE(4MD/) -

```

```

HEADER('Lpar')          ON(15,4,CH)          -
HEADER('Date')          ON(11,4,DT1,E'9999/99/99') -
HEADER('Time')          ON(7,4,TM1,E'99:99:99') -
HEADER('Jobname')       ON(19,8,CH)          -
HEADER('Volser')        ON(95,6,CH)          -
HEADER('Dataset')       ON(45,44,CH)         -
COPY FROM(SMF1718) TO(SMF18) USING(SMF8)
DISPLAY FROM(SMF18) LIST(JOBREN)            -
  TITLE('Job renamed datasets') DATE(4MD/)  -
  HEADER('Lpar')        ON(15,4,CH)          -
  HEADER('Date')        ON(11,4,DT1,E'9999/99/99') -
  HEADER('Time')        ON(7,4,TM1,E'99:99:99') -
  HEADER('Jobname')     ON(19,8,CH)          -
  HEADER('Volser')      ON(139,6,CH)         -
  HEADER('Old ds name') ON(45,44,CH)         -
  HEADER('New ds name') ON(88,44,CH)         -
COPY FROM(SMF1718) TO(SYS17) USING(SMF9)
DISPLAY FROM(SYS17) LIST(SYSDEL)           -
  TITLE('System tasks deleted datasets') DATE(4MD/) -
  HEADER('Lpar')        ON(15,4,CH)          -
  HEADER('Date')        ON(11,4,DT1,E'9999/99/99') -
  HEADER('Time')        ON(7,4,TM1,E'99:99:99') -
  HEADER('Jobname')     ON(19,8,CH)          -
  HEADER('Volser')      ON(95,6,CH)          -
  HEADER('Data set')    ON(45,44,CH)         -
COPY FROM(SMF1718) TO(SYS18) USING(SMFT)
DISPLAY FROM(SYS18) LIST(SYSREN)           -
  TITLE('System tasks renamed datasets') DATE(4MD/) -
  HEADER('Lpar')        ON(15,4,CH)          -
  HEADER('Date')        ON(11,4,DT1,E'9999/99/99') -
  HEADER('Time')        ON(7,4,TM1,E'99:99:99') -
  HEADER('Jobname')     ON(19,8,CH)          -
  HEADER('Volser')      ON(139,6,CH)         -
  HEADER('Old ds name') ON(45,44,CH)         -
  HEADER('New ds name') ON(88,44,CH)         -
/*
//SMFICNTL DD *
  OPTION SPANINC=RC4,VLSHRT
  INCLUDE COND=(6,1,BI,EQ,17,OR,6,1,BI,EQ,18)
//SMF7CNTL DD *
  OPTION SPANINC=RC4,VLSHRT
  INCLUDE COND=(6,1,BI,EQ,17,AND,
(19,4,CH,NE,C'DFSM',AND,
  19,4,CH,NE,C'IXGL',AND,
  19,4,CH,NE,C'DBTD',AND,
  19,4,CH,NE,C'DSND'))
//SMF8CNTL DD *
  OPTION SPANINC=RC4,VLSHRT
  INCLUDE COND=(6,1,BI,EQ,18,AND,
(19,4,CH,NE,C'DFSM',AND,

```



```

    19,4,CH,NE,C'IXGL',AND,
    19,4,CH,NE,C'DBTD',AND,
    19,4,CH,NE,C'DSND'))
//SMF9CNTL DD *
OPTION SPANINC=RC4,VLSHRT
INCLUDE COND=(6,1,BI,EQ,17,AND,
(19,4,CH,EQ,C'DFSM',OR,
19,4,CH,EQ,C'IXGL',OR,
19,4,CH,EQ,C'DBTD',OR,
19,4,CH,EQ,C'DSND'))
//SMF7CNTL DD *
OPTION SPANINC=RC4,VLSHRT
INCLUDE COND=(6,1,BI,EQ,18,AND,
(19,4,CH,EQ,C'DFSM',OR,
19,4,CH,EQ,C'IXGL',OR,
19,4,CH,EQ,C'DBTD',OR,
19,4,CH,EQ,C'DSND'))
/*

```

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## **Making the transition from Assembler to C on the mainframe – revisited**

### **PASSING PARAMETERS AND ARGUMENTS**

This article extends on a previous article in *MVS Update* (see issue 216, September 2004) that discussed the general problems of migrating from Assembler to C. The current article discusses the particular problems associated with passing parameters to and from C programs and functions. For completeness, some parts are taken from the original article.

### **PASSING ARGUMENTS**

All common mainframe programming languages (Assembler, C, COBOL, PL/I) differentiate between the form with which

arguments are passed to a main program (namely the program that is initially invoked, such as with the EXEC PGM= JCL statement) and to a subprogram (namely a program invoked with a CALL statement or function call). This discussion ignores any 'exotic' calling conventions, such as XPLINK.

A related topic concerns the returning of arguments. Unlike traditional Assembler, C is a stack-based language; this has consequences when arguments are returned. The stack contains temporary data maintained by the run-time environment (Language Environment). Each program that is Language Environment-conformant takes space from the stack when called and returns the space when it exits. For want of a better term, 'traditional Assembler' refers to Assembler programs that conform to the calling conventions that go back to the original OS, ie the Assembler program saves the caller's parameters in a save area provided by the caller, creates a static save area for use by lower-level called programs. Any data areas required by the Assembler program are located in static memory, typically in the same CSECT as the program code. Only in exceptional circumstances are data areas allocated dynamically, for example to provide space for very large tables. Assembler programs that use dynamic storage areas, for example Language Environment-conforming programs and re-entrant programs, are, in effect, stack-based and are subject to similar restraints as C programs.

### Main program

C programs differ from Assembler programs in not only the form of the arguments but also the method used to pass arguments. Both Assembler (also COBOL and most other mainframe languages) and C have different conventions for the arguments depending on whether a main program or a subprogram (function) is called.

Whereas an Assembler (COBOL, etc) main program receives the address (in register 1) of a pointer to the length-prefixed argument list, a C main program intrinsically receives two

parameters: the number of passed arguments and the array of pointers to the passed arguments (the invoked program name is passed as the first argument followed by the calling argument automatically parsed into blank-delimited subarguments; each argument is passed as a 0-delimited character string). This mimics the behaviour of personal computer compilers and so simplifies the porting of applications. The parsing can be suppressed by specifying the NOARGPARSE compiler option, in which case the passed array of pointers contains just two entries: the invoked program name and the specified invocation parameter, each of which is passed as a 0-delimited character string.

The C behaviour for passed arguments also has a significant difference when compared with other common programming languages (Assembler, COBOL, etc): namely, the passed arguments are automatically converted to lower case (even for the NOARGPARSE compiler option). There are, however, means of countering this behaviour.

Note: the conversion to lower case is performed character-by-character. This means any hexadecimal codes that happen to coincide with upper-case characters will also be converted, for example X'01C2' (namely, X'01',C'B') would be converted automatically to X'0192' (X'01',C'b').

### *Case 1 (standard behaviour)*

For example, the JCL statement:

```
// EXEC PGM=TESTPROG,PARM='ALPHA  BETA'
```

would cause the following parameters to be generated (note: to better illustrate the behaviour, two blanks are specified between ALPHA and BETA):

Assembler:

```
+-----+
|      |?--- register 1
+-----+
|
|
```

```

?
+-----+
|11|ALPHA  BETA|
+-----+

```

C:

```

+-----+
argv|  3  |
+-----+

+-----+
argc|    |    |    |
+-----+
      |    |    |    +-----+
      |    |    |    +-----?|testprog|#|
      |    |    |    +-----+
      |    |    |    +-----+
      |    |    |    +-----?|alpha|#|
      |    |    |    +-----+
      |    |    |    +-----+
      |    |    |    +-----?|beta|#|
      |    |    |    +-----+

```

where # represents the 0-delimiter ('\0', namely X'00').

The following C code fragment lists the passed parameters:

```

int i;

int main(int argc, char *argv[]) {
    for (i=0; i < argc; i++)
        printf("argv: %s\n", argv[i]);
}

```

*Case 2 (behaviour with the NOARGPARSE compiler option)*

For example, the JCL statement:

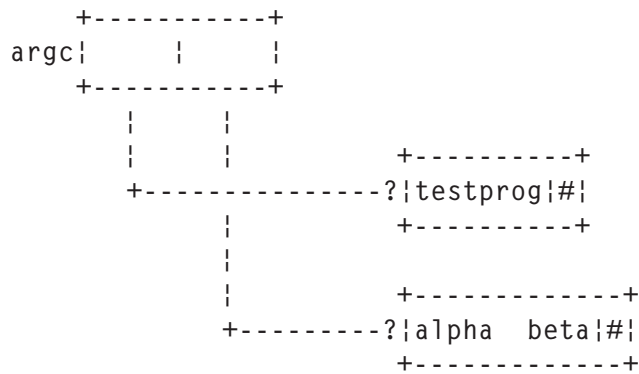
```
// EXEC PGM=TESTPROG,PARM='ALPHA  BETA'
```

would cause the following parameters to be generated:

```

+-----+
argv|  2  |
+-----+

```



where # again represents the 0-delimiter ('\0').

### Obtaining the true EXEC parameter

The [0] element from the `__osplist` built-in array (defined in `<stdlib.h>`) can be used to obtain the contents of register 1 when the program was invoked, namely, the address of a pointer to the EXEC parameter. This is the EXEC parameter as specified by the caller, ie it is not converted. The `__osplist` built-in array is briefly mentioned in the *Language Environment Programming Reference*.

```

#include <stdlib.h>
#include <stdio.h>

/* the following structure declares the form of the EXEC parameter */
struct {
    short vl;
    char vd[101];
} *pv;

main() {
    pv=(void *)__osplist[0];
    pv=(void *)pv;
    printf("vl:%hd\n",pv->vl); /* EXEC parameter length */
    printf("vd:%s\n",pv->vd); /* EXEC parameter data */
}

```

### Subprograms

Whereas, in Assembler, arguments are normally passed by reference, in C, arguments are implicitly passed by value. This means that the address-of operator (&) must be explicitly applied to pass an argument by reference (exception: strings

are always passed by reference). To ensure that the arguments are passed in the correct form, the calling program should include a prototype declaration of the called function (subprogram).

### Calling program (code fragment):

```
void FUNC(int, int *); /* function prototype */
int x1 = 1, x2 = 2;
```

```
FUNC(x1, &x2); /* invoke function */
```

Called function

```
void FUNC(int i1, int *pi2) {
    printf("p1: %d\n", i1);
    printf("p2: %d\n", *pi2);
}
```

FUNC() makes the following display:

p1: 1

p2: 2

### Calling legacy programs

Many programs make use of existing services, which can be classified as being legacy programs. In C, functions are called rather than programs. Although, as standard, such functions are invoked using C calling conventions, conventional legacy programs (subroutines) with the operating system calling convention can also be called when they have been identified with the `#pragma linkage (pgmname,OS)` instruction (*pgmname* is the name of the associated program).

The following code fragments show how a legacy program (here ISPLINK, the program interface for ISPF services) is invoked dynamically from C and Assembler. The two versions have the same functionality.

Note: for simplicity, the code to check that the program was loaded is omitted (null pointer returned from `fetch()` or zero address from the `LOAD` macro).

### C code fragment:

```
#pragma linkage (OSFUNC,OS)
typedef int OSFUNC();
```

```

int rc;                                /* return code */
OSFUNC *fptr;

char sn[] = "VCOPY ";                  /* service name */
char option[] = "MOVE ";               /* option */
char vn1[] = "(VN1 VN2)";             /* name list */
int vla[2] = {8, 8};                  /* length array */
char vna[2][8];                        /* value array */

fptr = (OSFUNC *)fetch("ISPLINK");
rc = (*fptr)(sn, vn1, &vla, &vna, option);

```

The equivalent Assembler code fragment:

```

LOAD EP=ISPLINK
LR 15,0 pass entry point address
CALL (15),(SN,VNL,VLA,VNA,MOVE),VL

SN DC CL8'VCOPY ' service name
OPTION DC CL8'MOVE ' option
VNL DC C'(VN1 VN2)'
VLA DC F'8,8'
VNA DS 2CL8

```

Special attention must be paid to the form of the arguments passed to the invoked program in C. The address-of operator (&) must be applied to the two arrays (vla and vna) to pass the required pointer. Determining the correct form of the passed arguments is the most difficult part of calling legacy programs from C.

### Returning arguments

Because local variables in a C program (and normally re-entrant Assembler programs) are contained in the stack, such variables are no longer available after the program has finished execution (or more correctly, cannot be guaranteed to be available – this is a danger, because, depending on when a variable is used, it might still be ‘present’) and so they cannot be used as the returned result. Only a value that can be returned in a register (eg short, int, float, char [a single character, but not a string – even if the string consists of just a single character], pointer, or a direct value that can be expressed as such) can be used as a direct return value, eg the field specified on the return instruction (obviously a pointer

may not contain the address of a local [non-static] variable). Instead, a static variable (ie a variable with static storage duration), a variable allocated in the heap (eg allocated with malloc()–), or a variable passed from the calling program (obviously as a pointer) must be used. Another possibility is the use of a global variable, namely a variable that is defined outside all program blocks.

The following demonstration programs all have the same functionality. The calling program passes an integer (the value 3) to the called function. The called function increments this passed value by 1 and returns this new value (4) to the calling program.

To keep the techniques as general and practical as possible, it is assumed that the calling program and the called function are separate compile units that are statically linked together. However, with the exception of referencing global variables, the techniques can also be used for dynamically loaded programs. Similarly, to avoid over-complication, the shown code fragments, although complete, are often trivial, but suffice to illustrate the concepts.

Calling program (result used directly):

```
int F02(int); /* function prototype */

void main()
{
    if (F02(3) == 4) puts("four");
}
```

Called function:

```
int F02(int i)
{
    int n; /* define local variable */
    n = i; /* transfer caller's argument to local variable */
    n++; /* increment */

    return n; /* return the local variable */
}
```



Because the calling value does not save the returned result, the function call can only be used immediately in an instruction.

Furthermore, in this particular simple case, the use of a local variable is not necessary. Because a copy of the caller's variable is passed, this copy could be incremented and returned.

Called function (alternative):

```
int F02(int i)
{
    i++;    /* increment copy of caller's argument */

    return i; /* return the incremented copy */
}
```

Calling program (result returned as a pointer):

The called function can be written in several ways:

```
int *F02(int); /* function prototype */

void main()
{
    int *pi;
    pi = F02(3);
}
```

Called function (result defined as a static variable):

```
int *F02(int i)
{
    static int n; /* static local variable */
    n = i; /* transfer caller's argument to local variable */
    n++;

    return &n; /* return address of local variable */
}
```

Note: if rather than **static int n**; merely **int n**; was written, the calling program would access the wrong data field, but because the address is probably still valid, no error would be indicated.

Called function (explicit memory allocation [heap] with malloc()):

```
#include <stdlib.h> /* header for malloc() */
```

```

int *F02(int i)
{
    int *pn;

    pn = malloc(sizeof(*pn)); /* allocate heap memory */
    *pn = i; /* transfer caller's argument to the allocated heap memory */
    (*pn)++;
    return pn; /* return the address of the allocated heap memory */
}

```

Called function (result defined as a global variable):

```

int n; /* <global variable, defined outside all program blocks */

void F02(int i, int *pi)
{
    n = i;
    n++;
    return &n;
}

```

Calling program directly accesses an external (global) variable:

If the accessed variable is defined as a global variable (as in this case), it can be accessed directly when it is declared as being external in the calling program.

Because good software engineering practice deprecates the use of global variables (unless their use is unavoidable), such a usage should be used with discretion.

```

#include <stdio.h>

extern int n;

int *F02(int);

void main()
{
    F02(3);
    printf("n:%d\n",n); /* display the result */
}

```

This is equivalent to EXTERN variables in Assembler, for example:

- L 1,=V(N) – address of external variable

- L 0,0(1) – load contents of external variable into register 0.

Calling program (result returned directly):

```
int F02(int); /* prototype */
```

```
void main()
```

```
{
    int i;
    i = F02(3);
}
```

1.4.3.1 Called function

```
int F02(int i)
```

```
{
    int n; /* local variable */
    n = i;
    n++;
```

```
    return n;
```

```
}
```

This form can be used here because the result can be passed as a register (in this particular case, as a general purpose register, but floating-point registers can also be used when appropriate).

Calling program (result returned in a variable passed by the calling program):

```
void F02(int, int *); /* prototype */
```

```
void main()
```

```
{
    int i; /* <i> will contain the returned value */
    F02(3, &i);
}
```

In this case, the calling program passes the address of a field (defined locally within the calling program) in which the called function returns its result.

Called function:

```
void F02(int i, int *pi)
```

```
{
    int n; /* work field */
```

```

n = i;
n++;
*pi = n;

return; /* no explicit return value */
}

```

### *Scope of the returned result*

The form with which a variable is defined determines its scope (possibly better known as its lifetime). The following list summarizes the various scopes. To distinguish between the calling program (main program) and the called function, the terms 'program' and 'function', respectively, are used.

- function result – the function result is transient, but can be used in the appropriate instruction.
  - function invocation – if nothing is specified, the variable is in the stack which exists only while the function is being invoked.
- static           function lifetime – the variable exists until the function is deleted from memory (implicitly or explicitly).
- malloc()        program lifetime – the variable exists while the program is loaded or until the variable is explicitly deleted from memory.
- file scope      program lifetime – the variable exists while the program is loaded. The variable is global and can be accessed directly with the extern qualifier. This is not good software engineering practice.

The programmer should be aware of the consequences of explicitly allocated variables (via malloc()), ie extensive use without the appropriate housekeeping can lead to memory leakage. Memory leakage occurs when memory is allocated but never freed after use. Because explicit memory allocation is relatively time-consuming, it should be used with caution when a large number of variables need to be allocated.

## FORWARDING A PARAMETER LIST

It is sometimes necessary to forward a parameter list, for example to call another subprogram. In Assembler, it suffices just to forward the contents of register 1; namely, the pointer to the parameter list. Because C normally needs to build the call at compilation time, a variable parameter list – processed with `argv[]` (main program) or `va_arg()` (subprogram) – must be converted into a fixed parameter list (eg an array) and then forwarded to the called subprogram, but even in this case the full flexibility of Assembler is not possible.

### Assembler program

Because the Assembler program does not need to know the format of the parameter list, the parameter list address can simply be forwarded. No differentiation needs to be made between a main program and a subprogram:

```
F01      CSECT
          BAKR 14,0  save caller's registers
          BASR 12,0  load base register
          USING *,12 set base register
          L    2,1   save caller's parmlist
...
* call subprogram
          L    1,2   reload caller's parmlist
          CALL F02   note: no parameters
...
          PR    ,    program return
```

For simplicity, the BAKR/PR instruction-pair is used to save and restore the caller's registers, although the usual calling conventions can also be used.

### C program

Here is an example of the procedure involved. The first function, `fo1()`, is called with a variable parameter list, `fo1(int, ...)`, where the second and following arguments are strings. (The number of the following arguments must be known, here it is specified as the first argument.) The arguments are passed to the second function, `fo2()`, as a fixed parameter list,

fo2(int, char \*pca[]). To simplify the code, the array has a fixed length.

```
#include <stdio.h>
#include <stdarg.h>

int fo1(int, ...);
int fo2(int, char *pca[]);

int main()
{
    fo1(2, "alpha", "beta");
    return 0;
}

int fo1(int n, ...) {
    va_list argptr;
    char *pc;
    int i;
    #define N 10 /* max. number of parms */
    char *pca[N];

    printf("N:%d\n",n); /* number of varparms */
    va_start(argptr, n);
    for (i=0; (i < n) & (i < N); i++) {
        pc = va_arg(argptr, char *);
        puts(pc);
        pca[i] = pc;
    }
    va_end(argptr);

    fo2(n,pca);
    return n;
}

int fo2(int n, char *pca[]) {
    char *pc;
    int i;

    printf("N:%d\n",n); /* number of parms */
    for (i=0; i < n; i++) {
        pc = pca[i];
        puts(pc);
    }
    return n;
}
```

---

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---

## No source, no worry – generating source code from the load module

Source code loss can strike anyone. Back-up devices may have failed or been misconfigured, or perhaps back-ups were not made at all. The loss may represent recent updates or additions – even the entire application; or other demands may require the modification of an old application for which the source code has been misplaced

It was the belief for years that you should re-write programs from scratch should the original source code become lost. Fortunately, there is now an alternative for recovering source code from object code. This article describes a formal technique for the recovery of COBOL and Assembler source code statements from the load module.

### THE LOAD MODULE ANALYSIS

First let us understand the load module and the information it can provide us with. Only load modules contain accurate information about the code running on a system. Load modules contain the machine instructions that are actually executed when a program runs. Each load module is composed of a number of CSECTs (control sections) containing the main program instructions, run-time routines supplied with compilers, and subroutines previously created by users. In addition to executable machine instructions, each CSECT also contains information – such as languages, compiler release levels, and compiler and linkage-editor options – that identifies what it does and how it was created. A load module begins with a source module written by a programmer in a high-level programming language. The source module must then be submitted to a compiler for the language the source is written in (there is a unique compiler for each programming language). The compiler changes the programming language into machine language – the machine instructions a processor can execute

– and puts them in an object module. The object module is then fed into the linkage editor, which transforms it into a load module.

When submitting a program to a compiler, programmers can specify parameters or options. The parameters selected can change the machine instructions generated from the source code. In most cases, a compiler inserts information about itself, such as version and release level, as well as information identifying what options were selected when the program was compiled into the load module. The format and location of this information may vary depending on the compiler and even the release of the compiler used. Like compilers, the linkage editor also has parameters that can be selected to vary the way a program executes. It also allows users to add IDR (Identification Record) data to load modules. And like most compilers, the linkage editor inserts identifying information about itself, parameters selected, and IDR data added into load modules as it creates them.

The linkage editor also inserts additional CSECTs into load modules for CALL statements that invoke statically linked system run-time routines or user-created subroutines to perform frequently used functions. Run-time routines are supplied with compilers or as part of the Language Environment (LE). Subroutines are user written and must have been previously compiled. Even load modules that dynamically invoke run-time routines from a resident library may have some statically linked routines, such as ‘bootstrap’ modules. Each subroutine contains its own identifying information about the compiler that created it.

## DISASSEMBLING

Disassembling, or source code recovery, is based on the following concepts or steps:

- 1 Delink
- 2 Disassembly



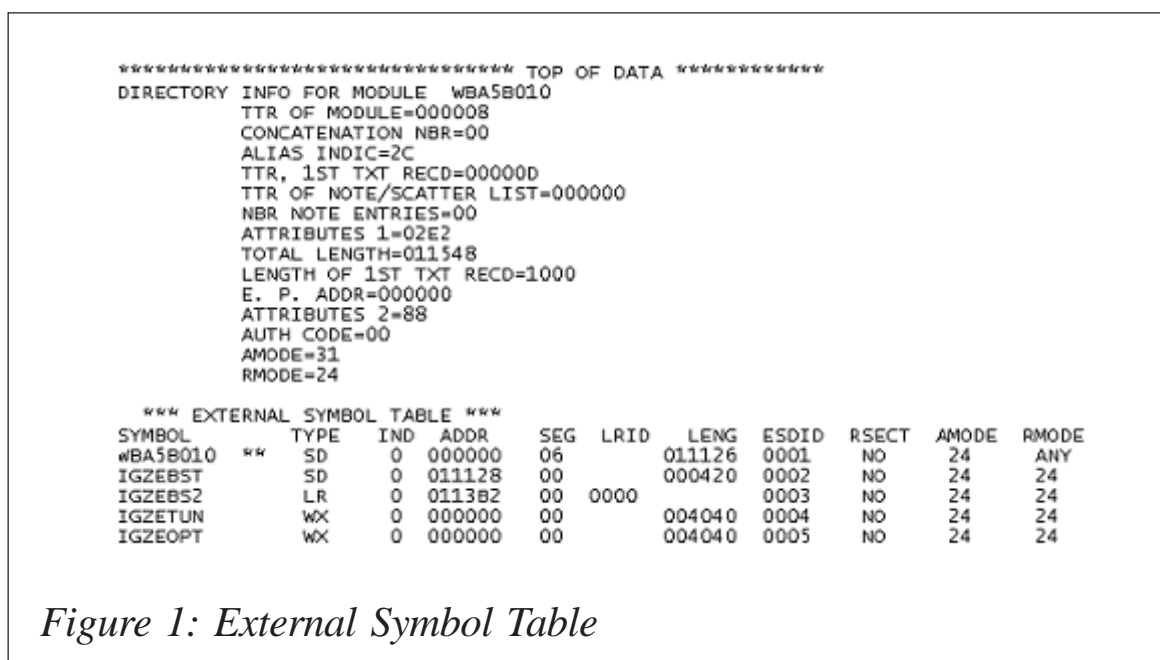
### 3 Decompile

- pattern matching
- operand analysis

### 4 Validation

#### Step 1: delink

The first step in recovering source code from an executable program is to break down the executable into its core components, known as control sections or CSECTs. This process identifies the individual modules that make up the executable. They include run-time modules provided by the operating system to incorporate system-level services such as CICS and/or DB2. Each of these CSECTs is intertwined in the executable, so in order to focus on one individual CSECT, the executable must first be delinked. From here, the basic 'program' can be disassembled. CSECT information can be obtained from the External Symbol Table in the SYSPRINT card from the spool after the first run of the disassembler on the load module. This is illustrated in Figure 1.



## Step 2: disassembly

The compiled 'program' is really a representation of ones and zeros, called binary. An example of the internals of a binary program file are shown below:

```
00000101 00011111 00001000 00000000 00000000 00000000 11010010 00010011
01100001 01100000 01100001 01001000 01011000 01000000 11010010 00110000
01000001 11100000 10111000 01101000 11010010 00000011 01000000 00100000
11000000 01001000 01011000 11110000 01000000 00001100 00000101 00011111
00000100 11000010 00000000 00000000 01000111 11110000 10111000 01110100
01011000 10110000 11000001 00111000 01000111 11110000 10110110 11111100
11010010 00000011 11010010 01100000 11010010 01100100 01000001 00000000
10111000 10000110 01010000 00000000 11010010 01100100 01000111 11110000
10110001 10001000 11010010 00000011 11010010 01100100 11010010 01100000
11010010 01011111 10100000 00000000 10000000 00000000 01001000 00000000
11000001 00001000 01011000 01000000 11010010 00101100 01000001 11100000
10111000 10101000 01011000 11110000 01000000 00001100 00000101 00011111
00001000 00000000 00000000 00000000 01011000 10110000 11000001 00111000
```

Programs represented in this manner, while accurate, are extremely difficult to inspect and interpret. At a somewhat higher level, the same program can be viewed from a byte-oriented or hexadecimal level of granularity, as shown below:

```
47F0F028 00C3C5C5 000000B0 00000014 47F0F001 98CEAC00
00000000 00000000 90ECD00C 4110F038 98EFF04C 07FF0000
00000748 000000AE 00000000 00000498 00000E38 000000CA
E2C1D4D7 D3C5F340 F2F0F0F5 F0F1F2F0 F1F4F3F1 F0F4F0F2
0000076C E0E86C4C 00800000 0288000A 80204000 00000000
00000018 00008000 40404040 0007E2C1 D4D7D3C5 F3400500
0000FFFF FFB20000 00000000 00380000 00080000 00060000
0D600000 00050000 00000000 00000000 00000000 00000000
00000108 00000434 00000476 00000498 FFFFFFFC 00001000
C940D4C1 D2C540C9 E340D7D6 E2E2C9C2 D3C5E2E8 E2D6E4E3
```

While somewhat easier to decipher, hexadecimal representation of a program is still at too low a level for most programmers. The same code can be further represented at a higher level – Assembler. The previous examples are shown below as a set of Assembler language statements and operands:

```
00040C 04 00000000 DC F'0'
000410 04 00000000 DC F'0'
000414 02 0000 DC H'0'
000416 08 D7D6E260C9404040 DC C'POS-I'
```

00041E	08	4040404040404040		DC	C'	'
000426	08	4040404040404040		DC	C'	'
00042E	06	404040404040		DC	C'	'
000434	04	58209128	A000434	L	R2,296(,R9)	
000438	06	D2012000A00C		MVC	0(2,R2),A000114	
00043E	06	D2012008A00C		MVC	8(2,R2),A000114	
000444	06	D2072010A032		MVC	16(8,R2),A00013A	
00044A	06	D2032018A014		MVC	24(4,R2),A00011C	
000450	06	D2032020A014		MVC	32(4,R2),A00011C	
000456	06	D2032028A014		MVC	40(4,R2),A00011C	
00045C	06	D2072030A02A		MVC	48(8,R2),A000132	
000462	06	D2032038A014		MVC	56(4,R2),A00011C	
000468	04	920E2040		MVI	64(R2),X'0E'	
00046C	04	920F2048		MVI	72(R2),X'0F'	
000470	04	5830D08C		L	R3,140(,R13)	
000474	02	07F3		BR	R3	
000476	04	92F18000	A000476	MVI	0(R8),C'1'	
00047A	06	D2018008A0BC		MVC	8(2,R8),A0001C4	
000480	06	D2028010A0B5		MVC	16(3,R8),A0001BD	

The code illustrated in the above three examples differs in the format in which the information is presented. Disassembly of the code, no matter what the original source language, is the first concept in source code recovery. The disassembler produces Assembler language source statements and a pseudo-listing using object code as input. You can use the Assembler language source file and listing to understand the program, debug it, and recover the lost source code.

We use DISASM, which is a one-pass disassembler that produces an Assembler language source program from a CSECT within a load module. Control cards permit specification of areas containing no instructions; allow base registers to be provided so that symbolic labels may be created during disassembly, and definition of DSECTs to be used during disassembly. Control statements permit the specification of areas containing instructions, data or uninitialized data areas, provide base registers so that symbolic labels are created during disassembly, and define the DSECTs used during disassembly.

Registers are denoted thus:

- Access Registers are denoted by A0, A1,...A15.

```

-----
//DISASM EXEC PGM=ASMDASM,PARM='options' 1
//SYSLIB DD DSN=user.loadlib,DISP=SHR 2
//SYSPRINT DD SYSOUT=A,DCB=BLKSIZE=121 3
//SYSPUNCH DD DSN=user.command.asm,DISP=(,CATLG), 4
// UNIT=SYSDA,DCB=BLKSIZE=3200,
// SPACE=(TRK,(5,2),RLSE)
//SYSIN DD *
module_name csect_name 5
other control statements
.
.
.
/*COPYLIB DD DSN=user.copy.name,DISP=SHR 6
-----

Sample Disassembler MVS JCL

```

*Figure 2: JCL statements used*

- Control Registers are denoted by C0, C1,...C15.
- Floating Point Registers are denoted by F0, F1,...F15.
- General Purpose Registers are denoted by R0, R1,...R15.
- Vector Registers are denoted by V0, V1,...V15.

The disassembler provides informational comments for recognized SVCs, and for various branch instructions to aid in creating a documented source program. You invoke the disassembler as a batch program using JCL. Figure 2 shows the JCL statements you can use.

### Step 3: decompile

#### *Pattern matching*

The next phase in the recovery of COBOL is that every COBOL verb will generate a distinct pattern of machine instructions that can be analysed and then programmatically stored, retrieved, and applied to the programs to be recovered. These patterns may be as simple as a single machine instruction or they may encompass many dozens of machine

instructions. Just as solving a jigsaw puzzle requires a sharp eye for subtle differences in the shapes of the puzzle's pieces, source code recovery requires an intimate knowledge of COBOL verbs and the patterns of machine code that they produce. For source code recovery, COBOL compiler research begins by creating source code files that provide examples of all the various programming elements in all their flavours, particularly Procedure Division verbs. Next, these research files are compiled with the PMAP option or, in the case of COBOL II, the LIST option. Finally, the compile listings are analysed to determine what verbs generate what particular patterns of machine instructions.

The PERFORM verb provides a good example of this type of research. The source code version of a simple PERFORM verb followed by the LIST option of the compile listing for that same PERFORM is shown below:

```
PERFORM 1000-OPEN-FILES THRU 1000-EXIT.
```

```

D2 03 D 214 D 234          MVC  214(4,13),234(13)          PSV=2  VN=02
41 00 B 032                LA   0,032(0,11)                GN=02
50 00 D 234                ST   0,234(0,13)                VN=02
58 B0 C 058                L   11,058(0,12)          PBL=1
47 F0 B 11A                BC  15,11A(0,11)        PN=03
D2 03 D 234 D 214          MVC  234(4,13),214(13)          VN=02  PSV=2

```

Analysing this pattern, and others like it, leads to the following conclusions: the pattern begins and ends with MVC instructions and the operands of these instructions are reversed. Finally, two instructions prior to the last statement in the pattern there is an unconditional branch to the paragraph to be performed. These instructions are distinctive enough for us to conclude that any identical patterns found in a COBOL program we are recovering must be generated by a simple PERFORM. Having done this type of analysis for all different verbs, and for many different COBOL compilers and compiler options, we can develop a DECOMPILER that can readily generate source code from the load module. Below is a simple REXX program to identify the starting and ending offset of a simple PERFORM pattern in the disassembled code:

```

/* MACRO TO COLLECT ALL THE MVC STATEMENTS FROM THE DISASSEMBLED CODE */
"ISREDIT MACRO"
ADDRESS ISPEXEC VGET ('OUTPDS') PROFILE
ADDRESS ISREDIT "EXCLUDE ALL"
ADDRESS ISREDIT "FIND ALL ' MVC '"
ADDRESS ISREDIT "DELETE ALL X"
ADDRESS ISREDIT "RESET"
ADDRESS ISREDIT "CRE .ZF .ZL &OUTPDS"
ADDRESS ISREDIT "CANCEL"
EXIT

```

```

/***** REXX *****/
/* THIS PROGRAM READS THE FILE WITH MVC STATEMENTS AND IDENTIFIES THE
PERFORM PATTERN.THE START AND END OFFSET OF THE PERFORM PATTERN IS
STORED IN PERFORM_BEGIN.I PERFORM_END.I */

```

```

ADDRESS TSO "ALLOC DA("OUTPDS") F(INFILE) SHR REUSE"
ADDRESS TSO "EXECIO * DISKR INFILE (FINIS STEM IN1."
COUNT1 = 0
DO I=1 TO IN1.0
  DO J=1 TO 2
    IF J=1 THEN
      DO
        X = LEFT(IN1.I,72) ; ORIGINAL = X
      END
    ELSE
      DO
        K = I + 1 ; X = LEFT(IN1.K,72)
      END
    NO_OF_WORDS = WORDS(X)
    COLUMN_OPERAND = NO_OF_WORDS
    COLUMN_INST = NO_OF_WORDS - 1
    OPERAND = WORD(X,COLUMN_OPERAND)
    INST = WORD(X,COLUMN_INST)
    IF J=1 THEN
      PARSE VALUE OPERAND WITH VAL1 '(' TEMP_1 ',' VAL2 '('
    ELSE
      PARSE VALUE OPERAND WITH VAL3 '(' TEMP_1 ',' VAL4 '('
  END /* J */
  IF VAL1 = VAL4 & VAL2 = VAL3 THEN
    DO
      TEMP_1 = WORD(ORIGINAL,1)
      IF PROGRAM_START_OFFSET < X2D(TEMP_1) THEN
        DO
          COUNT1 = COUNT1+1
          PERFORM_BEGIN.COUNT1 = TEMP_1
          PERFORM_END.COUNT1 = WORD(X,1)
          A_VAL2.COUNT1 = VAL2
        END
      END
    END
  END

```

```
END /* I */
```

```
ADDRESS TSO "FREE DD (INFILE)"
```

### *Operand analysis*

The pattern description language does a good job of recovering COBOL PROCEDURE DIVISION verbs, but does not provide much information about other elements in a COBOL program, particularly data items in file descriptions, WORKING STORAGE, and the LINKAGE SECTION. These items are recovered by applying the next concept in source code recovery, which is referred to as operand analysis. Operand analysis is the detailed examination and eventual analysis of the operands of the machine instructions making up the PROCEDURE DIVISION. This examination and analysis allows us to determine the relative placement of a data item and the PICTURE clause that is most appropriate for the data item. Data item recovery relies on analysis of research programs in much the same manner as the PROCEDURE DIVISION verb recovery. The compile listings are analysed and permit the determination of the types of PICTURE clauses that generate specific machine instructions. The following example illustrates that type of research and involves an elementary data item that is defined as a PIC 9:

```
Ø1 NUM1 PIC 999.
```

```
COMPUTE NUM1 = NUM1 + 1
```

F212 D1A8 9ØØØ	PACK 424(2,13),Ø(3,9)	TS2=Ø
96FØ D1A9	OI 425(13),X'ØF'	TS2=1
FA1Ø D1A8 AØ53	AP 424(2,13),83(1,1Ø)	TS2=Ø
F321 9ØØØ D1A8	UNPK Ø(3,9),424(2,13)	NUM1
96FØ 9ØØ2	OI 2(9),X'FØ'	NUM1+2
47FØ B134	BC 15,3Ø8(Ø,11)	GN=6(ØØØ3E8)

```
-----  
Ø1 NUM1 PIC 999.
```

```
Ø1 NUM2 PIC 999 COMP-4.
```

```
MOVE NUM1 TO NUM2.
```

F272 D180 9000	PACK 384(8,13),0(3,9)	TS2=0
960F D187	OI 391(13),X'0F'	TS2=7
4F30 D180	CVB 3,384(0,13)	TS2=0
4030 9008	STH 3,8(0,9)	NUM2

Our research has determined that whenever a PIC 9 item is involved in an arithmetic operation the final machine instruction generated will be determined with:

OI < >,X'0F'.

Finding this type of machine instruction during source code recovery allows us to conclude that the data item being operated on is an elementary item of type PIC 9. In addition the PACK and UNPK instructions indicate the length in bytes of the item, which can then be expressed as the number of 9s in the PICTURE clause.

PACK 424(2,13),0(3,9)

Here 3 indicates the length of the NUM1.

The actual operand analysis occurs during both the disassembly and the decompilation of the program being recovered. As each operand is encountered, the probable PICTURE clause value based on the machine instruction acting on the operand is determined and stored for later use during the source code generation phase of the decompiler.

Operand analysis allows you to recover accurately any data item that is referenced within the program.

#### Step 4: validation

The last step of the source code recovery methodology is simply to prove that the source code has been recovered accurately. In order to do this, we take the resulting recovered source code and compile it using the same version of the compiler (and compile options) that created the original executable program. The compile options that were originally specified are key for validation. The newly-created object module is then disassembled and compared with the disassembly of the original executable. While a 100% identical



object module is seldom expected, the technician must compare the original object with the newly-created object and verify that the resulting program is 'functionally' equivalent.

## SUMMARY

The process of recovering source code from object code is a formal process that results in source code that is guaranteed to be 100% functionally equivalent to the original code regardless of the original COBOL dialect, the options under which the program was compiled, or the operating system that the application is running on; and the decompiler forms the basis for detecting and recreating the corresponding source code.

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## Get the most out of subchannels: a STSCH illustration

### SUBCHANNEL

A subchannel number is a system-unique half-word value within a channel subsystem. The subchannel is addressed by the basic I/O instructions (SSCH, MMSCH, RSH, etc) and its status and configuration information is maintained in a special channel subsystem block – the subchannel-information block (SCHIB). You can easily retrieve that block using a STORE SUBCHANNEL (STSCH) instruction.

Let's take a look at the SCHIB. It has a twelve-word structure that contains three major fields:

- The path management control word (PCMW)
- The subchannel status word (SCSW)
- The model dependent area/measurement block.

Most of the information contained in the SCHIB, particularly the PCMW, is propagated by IOS to UCB. For example the Last Path Used Mask (LPUM), which contains the channel path that was last used for communicating between the channel subsystem and the device, is also maintained in the UCB IOS extension.

To illustrate my point, the following code will scan each subchannel (from 0 to 65535) and issue a STSCH to get the SCHIB and process path masks so it can print a simple path configuration for each device (only operational and last used path). Before issuing a STSCH, general register 1 must contain the subsystem-identification word (SID). This SID has been extended to provide both the CSS ID and MIF ID to which the subchannel is configured, but this information is implicitly provided at run-time by the LPAR hypervisor in a transparent way, so programs don't need to code the entire SID, only the last word (binary number one and a half-word subchannel number).

## I/O MEASUREMENT

Regarding I/O measurements, with the advent of z/990 and z/architecture comes a new 64-byte ECMB (replaces the old CMB) and the Extended Measurement Word (EMW). Both these measurement blocks contain device connect, disconnect, pending, busy time, and control unit queueing time, etc.

EMW is an extension of the IRB, and measurement data is provided on an I/O operation basis, while ECMB has a more sampling approach.

In order to access ECMB, the subchannel must have the following bits on in the related SCHIB:

- Measurement Mode Enable.
- Measurement Block Format Control (format 1).
- Measurement Block Update Enable.

It can then be pointed to by an absolute storage address – the measurement block address within the SCHIB. A much simpler way to access the ECMB is to get the ALET for the dataspace, the measurement block index from the UCB, and the compute displacement (see in the code, only the SSCH count is printed).

This sample program also gets the Configuration Data Record (CDR) of a device finally to provide a configuration report.

In short, SCHIB and ECMB, plus the UCB, can be easily combined to code an efficient device and path I/O monitor or a configuration report utility.

## CSSSCAN

```

CSSSSAN TITLE 'Scan Subchannels '
*****
* Name      : CSSSCAN
* Requires  : APF
* Function  : Scan subchannel (0=>X'ffff') and get SCHIB.
*           : Get PATH configuration through LPUM/PAM  SCHIB values
*           : Check Measurements Mode
*           : Access IOS data space to obtain Measurements Block info
*           : (only the accumulated SSCH count is printed).
*           : Obtain CDR to inform about manufacturer and model
*****
CSSSSAN CSECT
        COPY EQU
CSSSSAN AMODE 31
CSSSSAN RMODE 24
        BAKR  R14,0           Save
        LR    R12,R15        Using R12
        USING CSSSSAN,R12    Addressability
CONN   EQU    *
        OPEN (OUT,OUTPUT),MODE=31  Open sysout
        PUT  OUT,TITLE           Put 3  header
        PUT  OUT,TITLE2         lines
        PUT  OUT,TITLE3
T00    EQU    *

```

```

L      R5,=XL4'0000FFFF'          Load 65535 X'ffff'
XR     R6,R6                       Clear R6
*****
* Scan Subchannel 0=>65535          *
*****
T01    EQU      *
      MODESET  MODE=SUP,KEY=ZERO    To mode Sup
      STH     R6,SCIW+6             Store subchannel numb
      L      R1,SCIW+4             Load Sub Ident Word
      STSCH   SCHIB2               Get Schib
      LA     R2,SCHIB2             SCHIB addr
      USING  SCHIB,R2             addr for SCHIB
      STH     R6,ZSCN              Edit
      UNPK   ZSCN(5),ZSCN(3)       Subchannel Number
      TR     ZSCN(5),TAB

*

L      R3,SCHIP                    Select SCH with a
LTR    R3,R3                       non null Interrup.
BZ     NEXTSCH                     (Same as UCB)
MVC   WUCBPTR(4),SCHIP            Save UCB pointer
TM    SCHFLG2,SCHV               Check Device Validity
BO    OKDEVNO                     If Device validy bit off
B     NEXTSCH                     then skip this sch
OKDEVNO EQU      *
      MVC   WRKDEVD(2),SCHDEVNO    Save to access CDR
      MVC   ZDEVNUM(2),SCHDEVNO    Save Device numbr
      UNPK  ZDEVNUM(5),ZDEVNUM(3)  Edit Device numbr
      TR   ZDEVNUM(5),TAB

GETPATH EQU      *
      LA   R9,ZCHPTAB              Load CHPID edit table
      LA   R10,SCHCHPID            Chpids from SCHIB
      MVC  ZCHPTAB(24),=CL24' '    Clear print zone
      MVC  WRKCHPD(2),=H'0'        Zeroes
      LA   R3,8                    Max chpid per device = 8
PATH00 EQU      *
      XR   R8,R8                   Clear R8

**=====**
* Check PAM Path Available Mask
**=====**
      IC   R8,SCHPAM               Path available Mask byte
      LR   R4,R3                   Compute displacement
      BCTR R4,0                    within the mask with
      SRL  R8,0(R4)                8-loop count -1 - shift
      N    R8,=XL4'00000001'       Check Last bit only
      LTR  R8,R8                   If zero this path is
      BZ   NOPATH                  Not available to z/OS
      XR   R8,R8                   Clear

**=====**
* Print Chpid
**=====**

```

```

XR      R7,R7                      Clear R7
IC      R7,Ø(R1Ø)                  Get chpid
ST      R7,PARM                    Store
UNPK   WKZ(2*L'PARM+1),PARM(L'PARM+1)  format
TR      WKZ(2*L'PARM),TAB          Edit chpid number
MVC    Ø(2,R9),WKZ+6
MVI    2(R9),C' '

**=====
* Check PUM 'Last Path Used Mask'
**=====
IC      R8,SCHLPUM                 Last Used Path Mask byte
SRL    R8,Ø(R4)                   Shift with loop index
N      R8,=XL4'ØØØØØØØ1'         Last bit only
LTR    R8,R8                       if null
BZ     NXTPATH                     Process next chpid
MVC    ZLAST(2),WKZ+6             edit last used chpid
MVC    WRKCHPD+1(1),Ø(R1Ø)
B      NXTPATH                     Process next Chpid
NOPATH EQU *
MVC    Ø(2,R9),=CL2'..'         Indicates not available
NXTPATH EQU *
LA     R9,3(R9)                   Next Chp zone in output
LA     R1Ø,1(R1Ø)                 Next Chp in SCHIB TAB
BCT    R3,PATHØØ                 To next chpid
CLC    ZCHPTAB(24),=CL24' '     if at least one
BNE    GETCDR                     path found, getcdr
MVC    ZCHPTAB(24),=CL24'** No Operational Path **
B      GETCMB                     Print detail line

**=====
* GetCDR : Get Configuration data report with noio (<= from z/OS)
**=====
GETCDR EQU *
GETMAIN RU,LV=65535               GetMain
LR     R4,R1                      Keep storage addr
IOSCDR DEVN=WRKDEVD,CHPID=WRKCHPD+1,CDRAREA=(R4),CDRLEN=65535,*
      READ=NOIO,STATUS=CDRSTAT
LTR    R15,R15                    Indicates No CDR
BNZ    NOCDR                       if request failed
CLI    NEDFLAGS-NED(R4),X'CØ'     Process NED only
BE     NOCDR
MVC    ZCDR(3),NEDMANUF-NED(R4)   Format like
MVI    ZCDR+3,C'-'
MVC    ZCDR+4(6),NEDTYPEN-NED(R4) IBM-ØØ9Ø32
MVI    ZCDR+1Ø,C'-'              IBM-ØØ9Ø32-
MVC    ZCDR+11(3),NEDMODN-NED(R4) IBM-ØØ9Ø32-ØØ5
FREEMAIN RU,LV=65535,A=(R4)       Free Storage
B      GETCMB
NOCDR EQU *
MVC    ZCDR(16),=CL16'N/A'       Indicates no CDR
FREECDR EQU *

```

\*\*=====\*

\* GetCMD : Get Configuration data report with noio (<= from z/OS)

\*\*=====\*

```

GETCMB  EQU  *
        LA  R2,SCHIB2           Restore addr for SCHIB
        TM  SCHFLG2,SCHMM       Is MM Bit on
        BNO NOMM               Measurement off- skip
        L   R2,16               To CVT
        L   R2,CVTOPCTP-CVT(R2) To RMCT
        L   R3,RMCTCMCT-RMCT(R2) To CMCT
        CLI CMCTVERS-CMCT(R3),X'02' Level 2 => ECMB support
        BNE NOECMB             Not z/990 support
        CLI CMCTECMBFLAGS-CMCT(R3),CMCTECMBMODE See 0A06164
        BNE NOECMB             Not z/990 support
        IOSCMXA UCBPTR=WUCBPTR,UCBCXPTR=WUCBCXT To UCB Comm Ext
        L   R4,WUCBCXT          Save to addr
        L   R4,UCBIEXT-UCBETI(R4) UCB ios Extension
        LH  R9,UCBMBI-UCBRSTEM(R4) Get MBI
        LR  R7,R9               *64 to reach entry
        SLL R7,6                ECMB
        L   R1,16               To CVT
        L   R1,CVTOPCTP-CVT(R1) To RMCT
        L   R2,RMCTCMCT-RMCT(R1) To CMCT
        L   R8,CMCTECMBPTR-CMCT(R2) To ECMB pointer
        LAM R9,R9,CMCTECMBALET-CMCT(R2) ALET
        SAC 512                 Switch AR mode
        LA  R9,0(R7,R8)         Get ECMB addr
        USING ECMB,R9
        L   R1,ECMBSschRschCount Get ssch/rsch count
        SAC 0                   Switch AR mode
        CVD R1,PL8              Convert and
        MVC MASK1(10),=X'402020202020202020202020' print ssch count
        ED  MASK1(10),PL8+3      Edit
        MVC ZSSCH(10),MASK1
        B   WRITEIT             Print line
NOMM    EQU  *
        MVC ZSSCH(10),=CL10'MM-off' Measurement Mode off
        B   WRITEIT
NOECMB  EQU  *
        MVC ZSSCH(10),=CL10'Not supp ' Not Supported
WRITEIT EQU  *
        MODESET MODE=PROB,KEY=NZERO Back to prob mode
        PUT  OUT,ZSCN           Print detail
        MVI ZSCN,X'40'         Raz output zone
        MVC ZSCN+1(159),ZSCN
NEXTSCH EQU  *
        LA  R6,1(,R6)          Next subchannel
        BCT R5,T01             Loop
        B   EXIT               Exit
    
```

```

      CNOP  0,4
EXIT   EQU   *
      PR
      CNOP  0,4
**=====**
** SubChannel-Identification Word => bit                               *
** CSSID and MIF ID are provided by LPAR Hypervisor                   *
** see IBM J. RES. & DEV. VOL. 48 NO. 3/4 MAY/JULY 2004             *
** see z/Architecture Principles of Operation SA22-7832-02           *
**=====**
SCIW   DC   A(0)
      DC   XL2'0001',XL2'00'
SCHIB2 DC   12F'0'
OUT    DCB  MACRF=PM,DSORG=PS,RECFM=FB,LRECL=80,DDNAME=OUT
      CNOP  0,4
CDRSTAT DC  X'00'
      CNOP  0,4
WUCBPTR DC  F'0'
WUCBCXT DC  F'0'
MASK1  DC   X'40212020202020206B202020'
*=====*
* Edit zone                                                           =
*=====*
*
TITLE  DC  C'Subc Dev  Chpids                Last Activity'
TITLE10 DC C' Physical                          '
TITLE2  DC C'Num Adr  Available              Used (ECMB) '
TITLE20 DC C' Description                        '
TITLE3  DC C'-----+-----+-----+-----+-----+'
TITLE30 DC C'+-----+-----+-----+-----+-----+'
*
ZSCN   DC  CL4' ',C' '                SubChannel Number
ZDEVNUM DC CL4' ',C' '                Device number
ZCHPTAB DC CL24' ',C' '              Chpid Table
ZLAST  DC  CL4' ',C' '                Last Used Channel
ZSSCH  DC  CL10' ',C' '              Accumulated SSH count
ZCDR   DC  CL16' ',C' '              CDR returned from system
TAB    DC   15XL16'00'
      DC   C'0123456789ABCDEF'
*=====*
*=====*
WRKCHAN DS   F
WRKDEVD DS   H
WRKCHPD DS   H
PARM    DC   F'0'
PARM2   DC   F'0'
PL8     DS   PL8
*
      CNOP  0,4
RETCODE DC   F'0'

```

```

RSNCODE   DC   F'0'
WRKUCB    DS   CL100
WKZ       DS   CL8
REALECMB  DS   CL64
          CNOP 0,4
          LTORG
          IOSDDEVI
          IEFUCBOB LIST=YES,PREFIX=YES
          IHACDR
          IOSDUCBP
          IHASCHIB DSECT=YES
          IRACMCT
          IRARMCT
          CVT   DSECT=YES

```

```

*=====
* ECMB layout =
*=====

```

```

ECMB      DSECT
ECMBSschRschCount  DS F Number of SSCH/RSCH
ECMBSampleCount    DS F Number of SSCH/RSCH collected
ECMBConnectTime    DS F Summation of deviceconnect times
ECMBPendingTime    DS F Summation of SSCH/RSCH request pending
ECMBDisconnectTime DS F Summation of subchannel disconnect times
ECMBCUqueueingTime DS F Summation of control queueing times
ECMBDeviceActiveOnly DS F Summation of device active-onlytimes
ECMBDeviceBusyTime DS F Summation of device busy times
ECMBInitialCmdRespT DS F Initial commandresponse time
          END

```

## SAMPLE OUTPUT

Subc Num	Dev Adr	Chpids Available	Last Activity Used (ECMB)	Physical Description
0000	0100	AD 2E AB 86 33 32 2F 87	87	171 N/A
000D	010D	AD 2E AB 86 33 32 2F 87	33	1624 IBM-003390-B3C
0012	0112	AD 2E AB 86 33 32 2F 87	2F	1624 IBM-003390-B3C
0100	0200	9F 1C B3 34 C3 80 CF C7	9F	1136879 STK-003490-B40
0140	0421	9A .. .. .		MM-off N/A
0277	06E8	C6 .. .. .	C6	MM-off IBM-002084-CTC
042C	0F8C	F8 .. .. .	F8	MM-off IBM-001732-001
0431	0FF0	2B .. .. .	2B	MM-off IBM-009032-005
0C1D	428A	.. 9C 0C A8 CC AE .. .	AE	39620 HTC-003390-B3C

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## SMS status check REXX routine

We have recently had some issues with adding volumes into SMS, replacing previously replaced volume serial numbers. We found that in most cases these had been left in status QUINNEW or DISNEW by the work performed to remove them previously.

This was easily corrected, but did cause some confusion and also meant Help Desk calls were logged.

To avoid this situation I have written a small REXX EXEC that runs weekly and flags any volumes that have been removed, but are not in an ENABLE SMS status. We can then go in and re-enable these so that if volumes are added over old definitions they are immediately enabled.

This routine ensures that if the person re-using the volumes forgets to check the status we do not get Help Desk calls saying volumes cannot be utilized because of the old status.

The REXX EXEC that we run is shown below. Note that the following line of code is not testing for spaces but actually has hex '0000' coded. However, it prints as follows:

```
IF SUBSTR(DCOLREC1.C,145,4) = '      '
```

You can code the hex '0000' by entering HEX mode in ISPF edit.

Here's the SMSVOLCK EXEC:

```
/*REXX*/
```

```
/*          ***** */
/*      * PROGRAM NAME: SMSVOLCK.              * */
/*      *                               * */
/*      * PROGRAMMER:   JOHN BRADLEY.         * */
/*      *                               * */
/*      * PURPOSE:     THIS PROGRAM READS DCOLLECT VL TYPE * */
/*      *               RECORDS AND THEN LOOKS TO SEE WHETHER * */
/*      *               A VOLUME CODED IN A STORAGE GROUP IS * */
/*      *               NON-EXISTENT. IF IT IS IT CHECKS      * */
```

```

/*      *      WHETHER IT IS IN ANY STATE OTHER      *      */
/*      *      THAN 'ENABLE'. IF IT IS THEN A      *      */
/*      *      RECORD IS WRITTEN TO A REPORT SO IT      *      */
/*      *      CAN BE ENABLED.      *      */
/*      *****      */
/* READ ALL RECORDS FROM DCOLLECT OUTPUT.      */

TOTAL = 0      /* INITIALIZE TOTAL.      */
"EXECIO * DISKR DCOLIN (STEM DCOLREC1. FINIS)"

DO C = 1 TO DCOLREC1.0      /* LOOP FOR NO. OF RECORDS.      */
IF SUBSTR(DCOLREC1.C,5,2) = 'VL' THEN /* ONLY PROCESS TYPE 'VL'.      */
DO      /* IF IT IS THEN LOOP.      */
VOLSER = SUBSTR(DCOLREC1.C,27,6) /* GET THE VOLSER.      */

/* GET STORAGE GROUP NAME.      */

DVLSTGRP = LEFT(STRIIP(SUBSTR(DCOLREC1.C,91,8),T,'00'X),8)
DVLCSMSS = SUBSTR(DCOLREC1.C,121,1) /* GET STATUS.      */

SELECT      /* START SELECT.      */
WHEN DVLCSMSS = '00'X THEN /* STATUS 00 SET TO 'NONE'.      */
DVLCSMSS = "NONE "
WHEN DVLCSMSS = '01'X THEN /* STATUS 01 SET TO 'ENABLE'.      */
DVLCSMSS = "ENABLE "
WHEN DVLCSMSS = '02'X THEN /* STATUS 02 SET TO 'QUIESCE'.      */
DVLCSMSS = "QUIESCE"
WHEN DVLCSMSS = '03'X THEN /* STATUS 03 SET TO 'QUINEW'.      */
DVLCSMSS = "QUINEW "
WHEN DVLCSMSS = '04'X THEN /* STATUS 04 SET TO 'DISABLE'.      */
DVLCSMSS = "DISABLE"
WHEN DVLCSMSS = '05'X THEN /* STATUS 05 SET TO 'DISNEW'.      */
DVLCSMSS = "DISNEW "
END      /* END SELECT.      */

IF SUBSTR(DCOLREC1.C,145,4) = ' ' THEN /* IF HEX 0000 UNITADDR.      */
DO /* LOOP.      */
IF DVLCSMSS ^= 'ENABLE' THEN /* IF NOT = 'ENABLE'      */
DO
OREC = ' ' /* BLANK OREC FIELD.      */
OREC = ' ||VOLSER|| ' ||DVLSTGRP|| ' ||DVLCSMSS
QUEUE OREC /* OUTPUT RECORD.      */
TOTAL = TOTAL + 1 /* ADD 1 TO TOTAL.      */
END /* END LOOP.      */
END /* END LOOP.      */
END /* END LOOP.      */
END /* END LOOP.      */

IF TOTAL ^= 0 THEN /* RECORDS PROCESSED?      */
DO /* YES START LOOP.      */

```

```

HDR = ' VOLSER  STORGRP  STATUS' /* SET HEADER. */
PUSH HDR /* PUSH ONTO STACK. */
"EXECIO 1 DISKW DCOLOUR" /* WRITE HDR. */
HDR = ' -----' /* SET HEADER. */
PUSH HDR /* PUSH ONTO STACK. */
"EXECIO 1 DISKW DCOLOUR" /* WRITE HDR. */
HDR = ' ' /* SET HEADER. */
PUSH HDR /* PUSH ONTO STACK. */
"EXECIO 1 DISKW DCOLOUR" /* WRITE HDR. */
"EXECIO * DISKW DCOLOUR (FINIS)" /* ALL RECORDS DONE, WRITE. */
END /* END LOOP. */
EXIT 0 /* EXIT WITH 0. */

```

The JCL required to run the DCOLLECT to obtain input into the REXX and into later steps that run the EXEC itself are shown below:

```

//JXB7884S JOB (JXB),'JXB,J.BRADLEY',CLASS=A,MSGCLASS=H
/* *****
/* * GENERATE DCOLLECT RECORDS FROM ACTIVE SCDS. *
/* *****
/* *****
/* * DELETE DATASETS FROM LAST RUN. *
/* *****
/* *****
//STEP1 EXEC PGM=IDCAMS
//SYSPRINT DD SYSOUT=*
//SYSIN DD *
DELETE JXB.JB.DCOLIN
DELETE JXB.JB.DCOLOUR
/*
/* *****
/* * RUN DCOLLECT AGAINST ACTIVE SMS SCDS. *
/* *****
//STEP2 EXEC PGM=IDCAMS
//SYSPRINT DD SYSOUT=*
//OUTDD DD DSN=JXB.JB.DCOLIN,
// DCB=(RECFM=VB,LRECL=32576,BLKSIZE=0,DSORG=PS),
// SPACE=(CYL,(20,10),RLSE),
// DISP=(,CATLG,DELETE)
//SYSIN DD *
DCOLLECT -
OUTFILE(OUTDD) -
SMSDATA(ACTIVE)
/*
/* *****

```

```

//*      *
//*      * EXECUTE SMSVOLCK REXX EXEC.
//*      *
//*      *
//*      *
//STEP3      EXEC  PGM=IRXJCL,PARM='SMSVOLCK',REGION=4M
//SYSTSPRT   DD    SYSOUT=*,RECFM=FBA,LRECL=133,BLKSIZE=0
//SYSEXEC    DD    DSN=JXB.JB.CNTL,DISP=SHR
//SYSTSIN    DD    DUMMY
//DCOLIN     DD    DSN=JXB.JB.DCOLIN,DISP=SHR
//DCOLOUT    DD    DSN=JXB.JB.DCOLOUT,
//            DISP=(NEW,CATLG,DELETE),
//            SPACE=(CYL,(1,1),RLSE),
//            DCB=(RECFM=FB,LRECL=80,BLKSIZE=0,DSORG=PS)

```

Finally, below is an example of the output from the REXX EXEC.

VOLSER	STORGRP	STATUS
PRDA01	PROD3390	DISNEW
PRDA02	PROD3390	DISNEW
PRDA03	PROD3390	DISNEW
PRDA04	PROD3390	DISNEW
PRDA05	PROD3390	DISNEW
PRDA06	PROD3390	DISNEW
PRDA17	PROD3390	DISNEW
PRDA18	PROD3390	DISNEW
PRDA20	PROD3390	DISNEW
PRDA21	PROD3390	DISNEW

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## Syslog analysis using LOGLYZER

I needed to analyse my syslog data. Aside from using SORT and REXX, I couldn't find any tools that really made the job any easier. SORT allowed me to copy and select certain records, as long as what I was looking for was always in exactly the same place. REXX, of course, allows me to do whatever I

want. What I wanted was a way to search for data in each logical syslog record.

Each syslog record is actually a Write to Operator (WTO) message. Some WTOs are Multi-Line WTOs (MLWTO). When the MVS WRITELOG command is used, it externalizes the syslog. Then an external writer is used to copy it to a sequential dataset. This process causes many records in the syslog records to get broken across physical records. Things can get complicated when you are searching for data in the continuation records.

To deal with this, I wrote the LOGLYZER REXX EXEC. LOGLYZER will read a sequential syslog and rejoin all the records. It provides search parameters to select and/or reject records based on the value at a specific word offset. Multiple selection criteria can be used to reduce the amount of output. The qualifying records will be copied to the OUTPUT DD. Additionally, a statistical analysis of the messages is always produced after the processing is completed.

In LOGYLZER, each blank-delimited word in a syslog record becomes a 'token'. The token is the number of the word in the record. Using the following \$HASP100 syslog record as an example, each word would have the following token value.

```
      (1)  (2)  (3)  (4)  (5)  (6)  (7)  (8)
CS01N 02000000 CS01 05026 00:00:01.10 STC14981 00000091 $HASP100
      (9)  (10) (11)
SMFUDUMP ON STCINRDR
```

So, if I wanted to find all the records that were \$HASP100 messages for SMFUDUMP, I would want to find token 8 equal to \$HASP100 and token 9 equal to SMFUDUMP. Using my input format, this would be done using the following format:

```
8=$HASP100 9=SMFUDUMP
```

Token=value (using the '=' sign) is for selection, and token-value (using the '-' sign) is for rejection. All criteria are 'ANDed' together. Each value is 'wildcarded' against the actual value in the syslog record. Therefore, you can use the entire value

or just prefixes, suffixes, or substrings of the actual values.

If I want to find all DSNT501I messages (a seven line MLWTO) from DB2 subsystem DB2B with reason code 00C9008E, but not from job MYTEST, in this example DSNT501I message:

```
(1)      (2)      (3)          (4)      (5)          (6)      (7)      (8)
CS01M 40000000 CS01      05026 20:11:32.54 STC15059 00000090 DSNT501I -
(9)      (10)
DB2B DSNILMCL

      (11)      (12)      (13)          (14)          (15)
RESOURCE UNAVAILABLE 293 CORRELATION-ID=MYTEST CONNECTION-ID=DB2CALL
(16)      (17)
LUW-ID=* REASON

(18)      (19)      (20)      (21)      (22)          (23)
00C9008E TYPE 00000302 NAME CUSDB001.TSCUS      .X'100756'
```

I would use the following search criteria:

```
8=DSNT501I 9=DB2B 18=00C9008E 14-MYTEST
```

This would select all DSNT501I records from DB2B with reason code 00C9008E and ignore records from job MYTEST.

Here is sample JCL for LOGLYZER:

```
//jobcard...
//LOGLYZER EXEC PGM=IKJEFT01,DYNAMNBR=999,
//          PARM='LOGLYZER * 8=DSNT501I 9=DB2B 18=00C9008E 14-MYTEST'
//SYSEXEC DD DSN=your.rexx.pds,DISP=SHR
//SYSLOG DD DSN=your.syslog.dsn,DISP=SHR,DCB=BUFNO=60
//OUTPUT DD SYSOUT=*
//SYSTSPRT DD SYSOUT=*
//DIAGMSG DD SYSOUT=*
//SYSTSIN DD DUMMY
```

Here is some sample LOGLYZER output:

```
OPSR0ZLF ---- LOGLYZER started 26 Jan 2005 19:32:29 on SY01 --- JOB 4346
```

Search criteria:

```
Token 8 equals Value DSNT501
```

```
SYSLOG used: SYS0.SYSLOG.DAILY.SY01.SMR05026
```

```
559,765 records read from SYSLOG in 3.990338 seconds
```

SYSLOG starts at: 26 Jan 2005 00:00:00.01  
SYSLOG ends at : 26 Jan 2005 17:00:00.79

366,256 expanded records generated in 65.883703 seconds

8 records matched search criteria, output is in OUTPUT DD

457 unique Message ID's found

MSGID Counts:

Sorted Descending by Frequency			Sorted Alpha by MSGID		
SARSTC22	83662	33%	!AOP0000	23	<1%
\$HASP250	17211	6%	\$HASP000	123	<1%
\$HASP100	11517	4%	\$HASP003	20	<1%
IEF196I	7084	2%	\$HASP097	38	<1%
DTM1460I	6743	2%	\$HASP100	11517	4%
TSS7053I	6111	2%	\$HASP112	1	<1%
\$HASP373	4953	1%	\$HASP119	1	<1%
\$HASP395	4937	1%	\$HASP125	3	<1%
SVTM052I	4783	1%	\$HASP150	1404	<1%
\$HASP540	4724	1%	\$HASP160	779	<1%
IEF404I	4677	1%	\$HASP250	17211	6%
.					
ARC6018I	1	<1%	WER169I	1	<1%
ARC0415I	1	<1%	WER177I	1	<1%
ARC6019I	1	<1%	WER211B	1	<1%
ARC1540I	1	<1%	WER246I	1	<1%
ARC1541I	1	<1%	WER410B	2	<1%
IKT122I	1	<1%	WER416B	3	<1%

MSGID Total: 248,712

OPSR0ZLF-LOGLYZER ended 26 Jan 2005 19:33:53 83.7 on SY01 RC=0-JOB 4346

## LOGLYZER REXX EXEC

```
/*  
/*                               REXX                               */  
/*  
/* Purpose: Find strings in the SYSLOG                               */  
/*-----*/
```

```

/* Syntax: LOGLYZER readlimit token1 token2 token3 token4 token5... */
/*-----*/
/* Parms: READLIMIT - The number of records to read ('*' means all) */
/*          TOKENn   - Unlimited number of tokens                      */
/*          */
/* Notes: TOKENs are a pair of the positional word and the value      */
/* ie      8=IEA630I 10-OPSR0Z                                        */
/*          */
/* Use '=' for an equals comparison or '-' for a not equal compare    */
/* The leading characters of a value can be used for wildcarding      */
/* ie      8=IEA630I 10-OPS                                          */
/*          */
/* OPS would match OPS, OPSX, OPSM02, OPSR0Z, OPSR0ZXX, etc.        */
/*          */
/* Optional keyword tokens:      (only support with '=' comparison)  */
/* LPAR=SYSA                      (replaced with the target LPAR)    */
/* DATE=TODAY                      (replaced with today's julian date) */
/* DATE=YESTERDAY                  (replaced with yesterday's julian date) */
/* DATE=05085                      (any valid julian date)           */
/* TIME=HH:MM:SS.hh                (any valid time)                 */
/* TIME=10:                          (would get 10:00:00.00-10:59:59.99) */
/* MSGID=DSNT501                   (would get all DSNT501I messages) */
/* TASKID=JOB12345                 (can be jobnum, console, internal/instream) */
/* CMD=$C                          (can be any command or prefix)    */
/*          */
/* All WTOs (single or multi-line) are searched as a single record   */
/* Single line WTOs are copied as a single record to the OUTPUT DD   */
/* Multi-line WTOs are copied as separate lines to the OUTPUT DD    */
/*          */
/* Return Codes:                                                       */
/*      00 - Everything worked                                          */
/*      01 - Records truncated while writing to OUTPUT DD              */
/*      02 - Unexpected record types found in the SYSLOG              */
/*      04 - No records matched the search criteria                   */
/*      08 - All parms missing                                        */
/*      10 - Tokens missing                                          */
/*      11 - Parm or Token format error                              */
/*      12 - TSO ALLOC or EXECIO error                               */
/*          */
/* Sample JCL:                                                         */
/* //Jobcard...                                                       */
/* //LOGLYZER EXEC PGM=IKJEFT01,                                       */
/* //          PARM='LOGLYZER * 8=IEA630I 10-OPSR0Z'                  */
/* //SYSEXEC DD DSN=your.exec.pds,DISP=SHR EXEC PDS                   */
/* //SYSLOG DD DSN=your.syslog.dsn,DISP=SHR Syslog Input             */
/* //SYSLOGX DD DSN=your.syslogx.dsn,DISP=NEW Optional Output       */
/* //SYSTSPRT DD SYSOUT=* TS0/REXX Output                             */
/* //OUTPUT DD SYSOUT=* Matching records                             */
/* //DIAGMSGs DD SYSOUT=* Error messages                             */
/* //SYSTSIN DD DUMMY TS0/REXX Input                                 */

```



```

/*                                                                 */
/* The SYSLOGX DD is optional and can be used to capture expanded */
/* SYSLOG records. Start with an LRECL=1000 and use RECFM=VB. If  */
/* you get RC=1, it means some records were truncated.            */
/*****/
/*                          Change Log                            */
/***** @REFRESH BEGIN START      2005/04/03 12:55:22 *****/
/* Standard housekeeping activities                               */
/*****/
call time 'r'
parse arg parms
signal on syntax name trap
signal on failure name trap
signal on novalue name trap
probe = 'NONE'
vardump = 'NONE'
modtrace = 'NO'
modspace = ''
call stentry 'DIAGMSGS'
module = 'MAINLINE'
push trace() time('L') module 'From:' 0 'Parms:' parms
if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
call modtrace 'START' 0
/*****/
/* Set local estoeric names                                     */
/*****/
@vio   = 'VIO'
@sysda = 'SYSDA'
/***** @REFRESH END      START      2005/04/03 12:55:22 *****/
/* Set defaults                                               */
/*****/
version = '1.0.0'
/*****/
/* Default starting position of the syslog (record type location) */
/* N, S, M, D, E, L, O, etc. Set this to the offset of the record */
/* type if not position 1 (SMR requires logstart=5)            */
/*****/
logstart = 1
/*****/
/* Get yesterday's julian date (in case it is needed)          */
/*****/
yestjul = g2j(date('S',(date('B',date('S'),'S')-1),'B'))
/*****/
/* Accept, validate, and print the arguments                   */
/*****/
parse arg parms
if parms = '' then call rcexit 8 'No parms provided'
parse var parms readlimit parms
if readlimit <> '*' then
do

```

```

        if datatype(readlimit,'W') = 0 then
            call rcexit 11 'READLIMIT must be numeric or "*"
        end
    if parms = '' then call rcexit 10 'No tokens provided'
    origparm = parms
    tcount = words(parms)
    say
    say execname version 'Search criteria:'
    say
    do i=1 to tcount
        select
            when pos('=' ,word(parms,1)) <> 0 then
                do
                    parse var parms token '=' value parms
                    ro = 'equals'
                end
            when pos('-' ,word(parms,1)) <> 0 then
                do
                    parse var parms token '-' value parms
                    ro = 'not equal'
                end
            otherwise call rcexit 11,
                'Parm' i 'is invalid, must be n=value or n-value'
        end
    if datatype(token,'W') <> 1 then
        do
            select
                when token = 'LPAR' then nop
                when token = 'DATE' then nop
                when token = 'TIME' then nop
                when token = 'TASKID' then nop
                when token = 'MSGID' then nop
                when token = 'CMD' then nop
                when token = 'STRING' then nop
                otherwise call rcexit 11 'Token' i '('token') is invalid'
            end
        end
    if value = '' then
        call rcexit 11 'Value' i 'is missing'
    say 'Token' token ro 'Value' value
end
say
/*****
/* Confirm required files are available */
/*****
call ddcheck 'OUTPUT'
call ddcheck 'SYSLOG'
/*****
/* Identify the file being used */
/*****

```

```

if LRC = 0 then
  say 'SYSLOG used:' sysdsname
else
  say 'SYSLOG used: Can not determine DSN for tape datasets'
  say
  /*****
  /* Read the original SYSLOG
  /*****
  fcount = 0
  call tsotrap "EXECIO" readlimit "DISKR SYSLOG (STEM SYSLOG. FINIS"
  readtime = format(time('e'),,1)
  say commafy(syslog.0) 'records read from SYSLOG in' readtime 'seconds'
  say
  /*****
  /* Date and Time range
  /*****
  first = 2
  do forever
    parse var syslog.first . . . lsddate lstime .
    if substr(strip(syslog.first),logstart,1) = 'S' then
      do
        first = first + 1
        iterate
      end
    first = first + 1
    if datatype(lsddate,'W') = 1 then
      do
        say 'SYSLOG starts at:' right(date('N',lsddate,'J'),11) lstime
        leave
      end
    end
  end
  last = syslog.0
  do forever
    parse var syslog.last . . . ledate letime .
    if substr(strip(syslog.last),logstart,1) = 'S' then
      do
        last = last - 1
        iterate
      end
    last = last - 1
    if datatype(ledate,'W') = 1 then
      do
        say 'SYSLOG ends at :' right(date('N',ledate,'J'),11) letime
        leave
      end
    end
  end
  /*****
  /* Join multi-line messages and load the SYSLOGX stem
  /*****
  xcount = 0

```

```

msglist = ''
do l=1 to syslog.0
/*****/
/* Strip any leading ls, or spaces */
/*****/
    syslog.l = strip(syslog.l,'L','1')
    syslog.l = strip(syslog.l,'L',' ')
/*****/
/* Determine whether this is a single or multi-record message */
/*****/
    clrtype = substr(syslog.l,logstart,1)
    clrstyp = substr(syslog.l,logstart+1,1)
    if l < syslog.0 then
        do
            nextlr = l + 1
            nlrtype = substr(syslog.nextlr,logstart,1)
        end
    else
        nlrtype = 'N'
/*****/
/* Parse the 7th, 8th, and 9th tokens */
/*****/
    parse var syslog.l . . . . . token7 token8 token9 .
/*****/
/* Check whether the task identifier is present */
/*****/
    taskid = substr(syslog.l,logstart+37,8)
    if taskid = '      ' then taskid = 'N/A'
/*****/
/* Determine how to process each record type */
/*****/
    select
/*****/
/* Single records */
/*****/
        when clrtype = 'N' & nlrtype <> 'S' then
            do
                xcount = xcount + 1
                syslogx.xcount = syslog.l
                startx.xcount = 1
                endx.xcount = 1
/*****/
/* Determine which token to use as the msgid */
/*****/
            select
                when clrstyp = ' ' & taskid = 'N/A' then
                    call msgctr token7
                when clrstyp = 'C' then
                    do
                        syslogx.xcount = syslog.l

```

```

        msgid.xcount = 'COMMAND'
    end
    when clrstyp <> ' ' then
        call msgctr token7
    otherwise call msgctr token8
    end
end
end
/*****
/* L records - Label line of a multiple line message */
/*****
    when clrtype = 'L' then
    do
        xcount = xcount + 1
        syslog.xcount = syslog.l
        msgid.xcount = 'LABEL'
    end
/*****
/* Bypassed records (L and W) */
/*****
    when clrtype = 'W' then
    do
        xcount = xcount + 1
        syslog.xcount = syslog.l
/*****
/* Determine which token to use as the msgid */
/*****
        select
            when clrstyp = ' ' & taskid = 'N/A' then
                call msgctr token7
            when clrstyp <> ' ' then
                call msgctr token7
            otherwise call msgctr token8
        end
    end
/*****
/* Non MLWTO multi-line records */
/*****
    when clrtype = 'N' & nlrtype = 'S' then
    do
        xcount = xcount + 1
        parse var syslog.nextlr . remainder
        syslog.xcount = strip(syslog.l) strip(remainder)
        startx.xcount = 1
        endx.xcount = 1 + 1
/*****
/* Determine which token to use as the msgid */
/*****
        select
            when clrstyp = ' ' & taskid = 'N/A' then
                call msgctr token7

```

```

        when clrstyp <> ' ' then
            call msgctr token7
        otherwise call msgctr token8
    end
end
/*****
/* MLWTO 'M' records */
*****/
    when clrtype = 'M' then
    do
        xcount = xcount + 1
        syslogx.xcount = strip(syslog.l)
        startx.xcount = 1
        endx.xcount = 1
/*****
/* Determine which token to use as the msgid */
*****/
        select
            when clrstyp = ' ' & taskid = 'N/A' then
                call msgctr token7
            when clrstyp <> ' ' then
                call msgctr token7
            otherwise call msgctr token8
        end
    end
/*****
/* MLWTO 'D' records */
*****/
    when clrtype = 'D' then
    do
        parse var syslog.l . . . remainder
        syslogx.xcount = syslogx.xcount strip(remainder)
    end
/*****
/* MLWTO 'E' records */
*****/
    when clrtype = 'E' then
    do
        parse var syslog.l . . . remainder
        syslogx.xcount = syslogx.xcount strip(remainder)
        endx.xcount = 1
    end
/*****
/* Eliminate the 'S' continuation records */
*****/
    when clrtype = 'S' then nop
/*****
/* Eliminate the 'O' records */
*****/
    when clrtype = 'O' then

```

```

do
  xcount = xcount + 1
  syslog.xcount = syslog.1
  msgid.xcount = 'LOG'
end
/*****
/* Eliminate the 'X' records */
*****/
  when clrtype = 'X' then
do
  xcount = xcount + 1
  syslog.xcount = syslog.1
  call msgctr token8
end
/*****
/* Eliminate the '+' records */
*****/
  when clrtype = '+' then nop
/*****
/* Unexpected conditions */
*****/
  otherwise
do
  EXITRC = 2
  say '====>' syslog.1
end
end
end
/*****
/* Report number of expanded records */
*****/
say
exptime = format((time('e') - readtime),,1)
say commafy(xcount) 'expanded records generated in' exptime 'seconds'
say
/*****
/* Load SYSLOGX if SYSLOGX DD is found */
*****/
if listdsi("SYSLOGX" "FILE") = Ø then
do
  say 'SYSLOGX DD found, writing' commafy(xcount) 'expanded SYSLOG',
  'records to' sysdsname
  ERC = tsotrap(1 "EXECIO * DISKW SYSLOGX (STEM SYSLOGX. FINIS)")
  if ERC = 1 then
do
  say 'LRECL='syslrecl 'is too small for longest record, some',
  'records were truncated'
end
end
end
/*****

```

```

/* Search the expanded records */
/*****/
  ocount = 0
  do i=1 to xcount
/*****/
/* Look for the tokens */
/*****/
  parms = origparm
  do j=1 to tcount
/*****/
/* Determine whether this is an equal or not equal token pair */
/*****/
  if pos('=',word(parms,1)) <> 0 then
    do
      parse var parms token '=' value parms
      select
/*****/
/* Select all records for the specified julian date */
/*****/
      when token = 'DATE' then
        do
          msgdate = substr(syslogx.i,logstart+19,5)
          if value = 'TODAY' then
            value = date('J')
          if value = 'YESTERDAY' then
            value = yestjul
          if msgdate <> value then
            leave j
          end
/*****/
/* Select all records for the specified time */
/*****/
      when token = 'TIME' then
        do
          msgtime = substr(syslogx.i,logstart+25,11)
          if abbrev(msgtime,value) = 0 then
            leave j
          end
/*****/
/* Select all records for the specified Message ID */
/*****/
      when token = 'MSGID' then
        do
          if abbrev(msgid.i,value) = 0 then
            leave j
          end
/*****/
/* Select all records with the jobnum, console ID, etc */
/*****/
      when token = 'TASKID' then

```



```

do
    taskid = substr(syslogx.i,logstart+37,11)
    if abbrev(taskid,value) = 0 then
        leave j
    end
/*****
/* Select all records for a specific LPAR (can be any substring) */
*****/
    when token = 'LPAR' then
do
    lparid = substr(syslogx.i,logstart+10,4)
    if pos(value,lparid) = 0 then
        leave j
    end
/*****
/* Select all command records containing the text string */
*****/
    when token = 'CMD' then
do
    cmdrec = substr(syslogx.i,logstart,1)
    cmdind = substr(syslogx.i,logstart+1,1)
    if cmdrec <> 'N' then
        leave j
    if cmdind <> 'C' then
        leave j
    cmdlen = length(syslogx.i) - (logstart+56)
    cmdtxt = substr(syslogx.i,logstart+56,cmdlen)
    if pos(value,cmdtxt) = 0 then
        leave j
    end
/*****
/* Select all records containing the text string */
*****/
    when token = 'STRING' then
do
    if pos(value,syslogx.i) = 0 then
        leave j
    end
/*****
/* Select all records starting with the text string in the nth word */
*****/
    when abbrev(word(syslogx.i,token),value) = 0 then
        leave j
    otherwise nop
end
end
if pos('-',word(parms,1)) <> 0 then
do
    parse var parms token '-' value parms
    if abbrev(word(syslogx.i,token),value) = 1 then

```

```

        leave j
    end
/*****
/* Make sure the message has the number of tokens in the pair */
*****/
    if words(syslogx.i) < token then
        leave j
    else
        if j < tcount then
            iterate
        else
/*****
/* Put the matches in the output stem */
*****/
            do
                ortype = substr(syslogx.i,logstart,1)
                select
/*****
/* Write 'N' records to the output stem (expanded) */
*****/
                    when ortype = 'N' then
                        do
                            ocount = ocount + 1
                            output.ocount = syslogx.i
                            fcount = fcount + 1
                        end
/*****
/* Write 'M' records to the output stem (not expanded) */
*****/
                    when ortype = 'M' then
                        do
                            do k=startx.i to endx.i
                                ocount = ocount + 1
                                output.ocount = syslogx.k
                            end
                            fcount = fcount + 1
                        end
                    otherwise say '====> Unexpected ORTYPE' ortype syslogx.i
                end
            end
        end
    end
end
/*****
/* Load OUTPUT */
*****/
EXITRC = tsotrap(1 "EXECIO * DISKW OUTPUT (STEM OUTPUT. FINIS)")
/*****
/* Drop the SYSLOGX. and OUTPUT. stems */
*****/
drop syslogx. output.

```

```

/*****
/* Summary
/*****
select
  when fcount = 0 then
    do
      EXITRC = 4
      say fcount 'records matched search criteria, RC='EXITRC
    end
  when fcount = 1 then
    say fcount 'record matched search criteria,',
      'output is in OUTPUT DD'
  otherwise
    say commafy(fcount) 'records matched search criteria,',
      'output is in OUTPUT DD'
end
say
say commafy(words(msglist)) 'unique Message ID's found'
say
say
/*****
/* Message summary
/*****
say 'MSGID Counts:'
say
msgtotal = 0
do i=1 to words(msglist)
  interpret,
  'sortin.i = left(translate(word(msglist,i),'-'','_'),12)',
    'right('word(msglist,i)',7)'
  parse var sortin.i msgid msgcnt
  msgtotal = msgtotal + msgcnt
end
/*****
/* Sort the MSGLIST descending by count
/*****
call stemsort(20 'SORT FIELDS=(14,7,CH,D)')
/*****
/* Add percentages
/*****
do i=1 to sortout.0
  parse var sortout.i msgid msgcnt .
  msgid = left(msgid,12)
  msgcnt = strip(msgcnt,'T','00'x)
  msgpct = msgcnt/msgtotal*100
  msgcnt = right(commafy(msgcnt),9)
  if msgpct < 1 then
    sortin.i = msgid msgcnt ' <1%'
  else
    sortin.i = msgid msgcnt ' ' right(trunc(msgpct),3)%'

```

```

    msgcnt.i = sortin.i
end
/*****
/* Sort the MSGLIST ascending by MSGID */
/*****
call stemsort(30 'SORT FIELDS=(1,12,CH,A)')
/*****
/* Print the sorted MSGLIST in MSGCNT and MSGID sequence */
/*****
say
say 'Sorted Descending by Frequency          Sorted Alpha by MSGID'
say
do i=1 to sortout.0
    say msgcnt.i '      ' strip(sortout.i,'T','00'x)
end
say
say 'MSGID Total:' commafy(msgtotal)
say
/*****
/* Shutdown */
/*****
shutdown: nop
/*****
/* Put unique shutdown logic before the call to stdexit */
/***** @REFRESH BEGIN STOP      2002/08/03 08:42:33 *****/
/* Shutdown message and terminate */
/*****
call stdexit time('e')
/***** @REFRESH END STOP      2002/08/03 08:42:33 *****/
/* Non-refreshable subroutines */
/* MSGCTR - Maintain counters for each unique MSGID */
/*****
/* MSGCTR - Maintain counters for each unique MSGID */
/*****
msgctr: arg msgid
/*****
/* Deal with imbedded '-<>', leading '+', '*', '@' and trailing ':' */
/*****
msgid = strip(translate(msgid,'_ ','-<>'))
if left(msgid,1) = '+' then
    msgid = strip(strip(msgid,'L','+'))
if left(msgid,1) = '*' then
    msgid = strip(strip(msgid,'L','*'))
if left(msgid,1) = '@' then
    msgid = strip(strip(msgid,'L','@'))
if left(reverse(msgid),1) = ':' then
    msgid = strip(strip(msgid,'T',':'))
/*****
/* Filter things that are not message IDs, but set a bogus msgid.x */
/*****

```

```

trash = 'NO'
select
  when verify(msgid,'0123456789+', 'M') = 0 then trash = 'YES'
  when verify(msgid,'().,/*=?_','M') <> 0 then trash = 'YES'
  when left(msgid,1) = '#' then trash = 'YES'
  when left(msgid,1) = '-' then trash = 'YES'
  when left(msgid,2) = '+-' then trash = 'YES'
  when left(msgid,1) = '$' & left(msgid,5) <> '$HASP' then
    trash = 'YES'
  when pos("'",msgid) <> 0 then trash = 'YES'
  when datatype(left(msgid,1)) = 'NUM' then trash = 'YES'
  when length(msgid) < 6 then trash = 'YES'
  when length(msgid) > 11 then trash = 'YES'
  otherwise nop
end
if trash = 'YES' then
  do
    msgid.xcount = '$TRASH$'
    return
  end
/*****
/* If the message is not already on the MSGLIST, add it          */
/*****
    if pos(msgid,msglist) = 0 then msglist = msglist msgid
/*****
/* Load the MSGID stem                                          */
/*****
    msgid.xcount = msgid
/*****
/* Set the MSGID VAR to 0, if this is the first occurrence      */
/*****
    select
      when symbol(msgid) = 'LIT' then interpret msgid '=' 0
      when symbol(msgid) = 'BAD' then return
      otherwise nop
    end
/*****
/* Increment the counter                                          */
/*****
    interpret msgid '=' msgid '+' 1
    return
/***** @REFRESH BEGIN SUBBOX 2004/03/10 01:25:03 *****/
/* 26 Internal Subroutines provided in LOGLYZER                  */
/* Last Subroutine REFRESH was 4 Apr 2005 02:31:14              */
/*                                                                */
/* RCEXIT - Exit on non-zero return codes                       */
/* TRAP - Issue a common trap error message using rcexit       */
/* ERRMSG - Build common error message with failing line number */
/* STENTRY - Standard Entry logic                               */
/* STDEXIT - Standard Exit logic                                */

```

```

/* MSG      - Determine whether to SAY or ISPEXEC SETMSG the message */
/* DDCHECK  - Determine whether a required DD is allocated          */
/* DDLIST   - Returns number of DD's and populates DDLIST variable  */
/* DDSNS    - Returns number of DSNs in a DD and populates DDSNS   */
/* QDSN     - Make sure there are only one set of quotes           */
/* UNIQDSN  - Create a unique dataset name                         */
/* TSOTRAP  - Capture the output from a TSO command in a stem      */
/* TSOQUIET - Trap all output from a TSO command and ignore failures */
/* SAYDD    - Print messages to the requested DD                   */
/* JOBINFO  - Get job related data from control blocks             */
/* COMMAFY  - Add commas to a number                               */
/* STEM     - Set a value into a stem and increment the number     */
/* STEMSORT - Sort a stem variable (caller must strip nulls)      */
/* VIODD    - EXECIO a stem into a sequential dataset              */
/* G2J      - Convert a Gregorian date to a Julian data            */
/* PTR      - Pointer to a storage location                         */
/* STG      - Return the data from a storage location               */
/* MODTRACE - Module Trace                                         */
/***** @REFRESH END   SUBBOX   2004/03/10 01:25:03 *****/
/***** @REFRESH BEGIN RCEXIT  2005/04/03 12:58:07 *****/
/* RCEXIT   - Exit on non-zero return codes                        */
/*-----*/
/* EXITRC   - Return code to exit with (if non zero)              */
/* ZEDLMSG  - Message text for it with for non zero EXITRCs      */
/*****
rcexit: parse arg EXITRC zedlmsg
        EXITRC = abs(EXITRC)
        if EXITRC <> 0 then
            do
                trace 'o'
/*****
/* If execution environment is ISPF then VPUT ZISPFRC            */
/*****
        if execenv = 'TSO' | execenv = 'ISPF' then
            do
                if ispfenv = 'YES' then
                    do
                        zisprc = EXITRC
/*****
/* Does not call ISPWRAP to avoid obscuring error message modules */
/*****
                address ISPEXEC "VPUT (ZISPFRC)"
            end
        end
/*****
/* If a message is provided, wrap it in date, time, and EXITRC    */
/*****
        if zedlmsg <> '' then
            do
                zedlmsg = time('L') execname zedlmsg 'RC='EXITRC

```

```

        call msg zedlmsg
    end
/*****
/* Write the contents of the Parentage Stack */
/*****
        stacktitle = 'Parentage Stack Trace ('queued()' entries):'
/*****
/* Write to MSGDD if background and MSGDD exists */
/*****
        if tsoenv = 'BACK' then
            do
                if subword(zedlmsg,9,1) = msgdd then
                    do
                        say zedlmsg
                        signal shutdown
                    end
                else
                    do
                        call saydd msgdd 1 zedlmsg
                        call saydd msgdd 1 stacktitle
                    end
                end
            end
        else
/*****
/* Write to the ISPF Log if foreground */
/*****
            do
                zerrlm = zedlmsg
                address ISPEXEC "LOG MSG(ISRZ003)"
                zerrlm = center(' 'stacktitle' ',78,'-')
                address ISPEXEC "LOG MSG(ISRZ003)"
            end
/*****
/* Unload the Parentage Stack */
/*****
            do queued()
                pull stackinfo
                if tsoenv = 'BACK' then
                    do
                        call saydd msgdd 0 stackinfo
                    end
                else
                    do
                        zerrlm = stackinfo
                        address ISPEXEC "LOG MSG(ISRZ003)"
                    end
                end
            end
/*****
/* Print the VARDUMP values (if present) */
/*****

```

```

if vardump <> 'NONE' then
do
  if tsoenv = 'BACK' then
do
  say
  say 'Selected variable values:'
  say
do vd=1 to words(vardump)
  interpret "say word(vardump,vd) '='",
    word(vardump,vd)
end
  say
end
end
else
do
  zerrlm = 'Selected variable values:'
  address ISPEXEC "LOG MSG(ISRZ003)"
do vd=1 to words(vardump)
  interpret "zerrlm = word(vardump,vd) '='",
    word(vardump,vd)
  address ISPEXEC "LOG MSG(ISRZ003)"
end
end
/*****
/* Put a terminator in the ISPF Log for the Parentage Stack */
/*****
  if tsoenv = 'FORE' then
do
  zerrlm = center(' 'stacktitle' ',78,'-')
  address ISPEXEC "LOG MSG(ISRZ003)"
end
/*****
/* Signal SHUTDOWN. SHUTDOWN label MUST exist in the program */
/*****
  signal shutdown
end
else
  return
/***** @REFRESH END RCEXIT 2005/04/03 12:58:07 *****/
/***** @REFRESH BEGIN TRAP 2004/12/13 14:00:48 *****/
/* TRAP - Issue a common trap error message using rcexit */
/*-----*/
/* PARM - N/A */
/*****
trap: trace 'off'
  trapttype = condition('C')
  if trapttype = 'SYNTAX' then
    msg = errortext(RC)
  else

```



```

    msg = condition('D')
    trapline = strip(sourceline(sigl))
    msg = traptypes 'TRAP:' msg', Line:' sigl ""'trapline'"
    if trap = 'YES' & tsoenv = 'BACK' then
    do
        trap = 'NO'
        traplinemsg = msg
        say traplinemsg
        signal on syntax name trap
        signal on failure name trap
        signal on novalue name trap
        say
        say center(' Trace of failing instruction ',78,'-')
        trace 'i'
        interpret trapline
    end
    if trap = 'NO' & tsoenv = 'BACK' then
    do
        say center(' Trace of failing instruction ',78,'-')
        say
    end
    if tsoenv = 'FORE' then
        call rcexit 666 msg
    else
        call rcexit 666 traplinemsg
/***** @REFRESH END TRAP 2004/12/13 14:00:48 *****/
/***** @REFRESH BEGIN ERRMSG 2002/08/10 16:53:04 *****/
/* ERRMSG - Build common error message with failing line number */
/*-----*/
/* ERRLINE - The failing line number passed by caller from SIGL */
/* TEXT - Error message text passed by caller */
/*****
errmsg: nop
    parse arg errline text
    return 'Error on statement' errline',' text
/***** @REFRESH END ERRMSG 2002/08/10 16:53:04 *****/
/***** @REFRESH BEGIN STDENTRY 2005/01/30 07:49:49 *****/
/* STDENTRY - Standard Entry logic */
/*-----*/
/* MSGDD - Optional MSGDD used only in background */
/*****
stdentry: module = 'STDENTRY'
    if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
    parse arg sparms
    push trace() time('L') module 'From:' sigl 'Parms:' sparms
    arg msgdd
    parse upper source . . execname . . execdsn . . execenv .
/*****
/* Start-up values */
/*****

```

```

EXITRC = Ø
MAXRC = Ø
trap = 'YES'
ispfenv = 'NO'
popup = 'NO'
lockpop = 'NO'
headoff = 'NO'
hcreator = 'NO'
keepstack = 'NO'
lpar = mvsvar('SYSNAME')
jobname = mvsvar('SYMDEF','JOBNAME')
zedlmsg = 'Default shutdown message'
lower = xrange('a','z')
upper = xrange('A','Z')
/*****
/* Determine environment */
/*****
    if substr(execenv,1,3) <> 'TSO' & execenv <> 'ISPF' then
        tsoenv = 'NONE'
    else
        do
            tsoenv = sysvar('SYSENV')
            signal off failure
            "ISPQRY"
            ISPRC = RC
            if ISPRC = Ø then
                do
                    ispfenv = 'YES'
/*****
/* Check if HEADING ISPF table exists already, if so set HEADOFF=YES */
/*****
                call ispwrap "VGET (ZSCREEN)"
                if tsoenv = 'BACK' then
                    htable = jobinfo(1)||jobinfo(2)
                else
                    htable = userid()||zscreen
                TBCRC = ispwrap(8 "TBCREATE" htable "KEYS(HEAD)")
                if TBCRC = Ø then
                    do
                        headoff = 'NO'
                        hcreator = 'YES'
                    end
                else
                    do
                        headoff = 'YES'
                    end
                end
                signal on failure name trap
            end
/*****

```

```

/* MODTRACE must occur after the setting of ISPFENV */
/*****/
    call modtrace 'START' sig1
/*****/
/* Start-up message (if batch) */
/*****/
    startmsg = execname 'started' date() time() 'on' lpar
    if tsoenv = 'BACK' & sysvar('SYSNEST') = 'NO' &,
        headoff = 'NO' then
        do
            jobinfo = jobinfo()
            parse var jobinfo jobtype jobnum .
            say jobname center(' 'startmsg' ',61,'-') jobtype jobnum
            say
            if ISPRC = -3 then
                do
                    call saydd msgdd 1 'ISPF ISPQRY module not found,',
                        'ISPQRY is usually in the LINKLST'
                    call rcexit 20 'ISPF ISPQRY module is missing'
                end
/*****/
/* If MSGDD is provided, write the STARTMSG and SYSEXEC DSN to MSGDD */
/*****/
    if msgdd <> '' then
        do
            call ddcheck msgdd
            call saydd msgdd 1 startmsg
            call ddcheck 'SYSEXEC'
            call saydd msgdd 0 execname 'loaded from' sysdsname
/*****/
/* If there are PARMS, write them to the MSGDD */
/*****/
            if parms <> '' then
                call saydd msgdd 0 'Parms:' parms
/*****/
/* If there is a STEPLIB, write the STEPLIB DSN MSGDD */
/*****/
            if listdsi('STEPLIB' 'FILE') = 0 then
                do
                    steplibs = dddsns('STEPLIB')
                    call saydd msgdd 0 'STEPLIB executables loaded',
                        'from' word(ddsns,1)
                    if dddsns('STEPLIB') > 1 then
                        do
                            do stl=2 to steplibs
                                call saydd msgdd 0 copies(' ',31),
                                    word(ddsns,stl)
                            end
                        end
                    end
                end
end

```

```

        end
    end
/*****
/* If foreground, save ZFKA and turn off the FKA display          */
/*****
else
do
    fkaset = 'OFF'
    call ispwrap "VGET (ZFKA) PROFILE"
    if zfka <> 'OFF' & tsoenv = 'FORE' then
do
    fkaset = zfka
    fkacmd = 'FKA OFF'
    call ispwrap "CONTROL DISPLAY SAVE"
    call ispwrap "DISPLAY PANEL(ISPBLANK) COMMAND(FKACMD)"
    call ispwrap "CONTROL DISPLAY RESTORE"
end
end
/*****
    pull tracelvl . module . sigl . sparms
    call modtrace 'STOP' sigl
    interpret 'trace' tracelvl
    return
/***** @REFRESH END    STDENTRY 2005/01/30 07:49:49 *****/
/***** @REFRESH BEGIN STDEXIT 2004/08/02 06:06:40 *****/
/* STDEXIT - Standard Exit logic */
/*-----*/
/* ENDTIME - Elapsed time */
/* Note: Caller must set KEEPSTACK if the stack is valid */
/*****
stdexit: module = 'STDEXIT'
    if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
    parse arg sparms
    push trace() time('L') module 'From:' sigl 'Parms:' sparms
    call modtrace 'START' sigl
    arg endtime
    endmsg = execname 'ended' date() time() format(endtime,,1)
/*****
/* if MAXRC is greater than EXITRC then set EXITRC to MAXRC */
/*****
EXITRC = max(EXITRC,MAXRC)
endmsg = endmsg 'on' lpar 'RC='EXITRC
if tsoenv = 'BACK' & sysvar('SYSNEST') = 'NO' &,
    headoff = 'NO' then
do
    say
    say jobname center(' 'endmsg' ',61,'-') jobtype jobnum
/*****
/* Make sure this isn't a MSGDD missing error then log to MSGDD */
/*****

```

```

        if msgdd <> '' & subword(zedlmsg,9,1) <> msgdd then
            do
                call saydd msgdd 1 execname 'ran in' endtime 'seconds'
                call saydd msgdd 0 endmsg
            end
        end
/*****
/* If foreground, reset the FKA if necessary */
*****/
        else
            do
                if fkaset <> 'OFF' then
                    do
                        fkafix = 'FKA'
                        call ispwrap "CONTROL DISPLAY SAVE"
                        call ispwrap "DISPLAY PANEL(ISPBLANK) COMMAND(FKAFIX)"
                        if fkaset = 'SHORT' then
                            call ispwrap "DISPLAY PANEL(ISPBLANK)",
                                "COMMAND(FKAFIX)"
                        end
                        call ispwrap "CONTROL DISPLAY RESTORE"
                    end
                end
/*****
/* Clean up the temporary HEADING table */
*****/
                if ispfenv = 'YES' & hcreator = 'YES' then
                    call ispwrap "TBEND" htable
/*****
/* Remove STDEXIT and MAINLINE Parentage Stack entries, if there */
*****/
                call modtrace 'STOP' sigl
                if queued() > 0 then pull . . module . sigl . sparms
                if queued() > 0 then pull . . module . sigl . sparms
                if tsoenv = 'FORE' & queued() > 0 & keepstack = 'NO' then
                    pull . . module . sigl . sparms
/*****
/* if the Parentage Stack is not empty, display its contents */
*****/
                if queued() > 0 & keepstack = 'NO' then
                    do
                        say queued() 'Leftover Parentage Stack Entries:'
                        say
                        do queued()
                            pull stackundo
                            say stackundo
                        end
                        EXITRC = 1
                    end
                end
/*****
/* Exit */
*****/

```

```

/*****
      exit(EXITRC)
/***** @REFRESH END   STDEXIT   2004/08/02 06:06:40 *****/
/***** @REFRESH BEGIN MSG       2002/09/11 01:35:53 *****/
/* MSG      - Determine whether to SAY or ISPEXEC SETMSG the message */
/*-----*/
/* ZEDLMSG  - The long message variable                               */
/*****
msg: module = 'MSG'
      parse arg zedlmsg
      if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
      parse arg sparms
      push trace() time('L') module 'From:' sigl 'Parms:' sparms
      call modtrace 'START' sigl
/*****
/* If this is background or OMVS use SAY                               */
/*****
      if tsoenv = 'BACK' | execenv = 'OMVS' then
        say zedlmsg
      else
/*****
/* If this is foreground and ISPF is available, use SETMSG           */
/*****
      do
        if ispfenv = 'YES' then
/*****
/* Does not call ISPWRAP to avoid obscuring error message modules   */
/*****
          address ISPEXEC "SETMSG MSG(ISRZ000)"
        else
          say zedlmsg
      end
      pull trancelvl . module . sigl . sparms
      call modtrace 'STOP' sigl
      interpret 'trace' trancelvl
      return
/***** @REFRESH END   MSG       2002/09/11 01:35:53 *****/
/***** @REFRESH BEGIN DDCHECK   2004/11/09 22:48:36 *****/
/* DDCHECK  - Determine whether a required DD is allocated           */
/*-----*/
/* DD      - DDNAME to confirm                                       */
/*****
ddcheck: module = 'DDCHECK'
          if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
          parse arg sparms
          push trace() time('L') module 'From:' sigl 'Parms:' sparms
          call modtrace 'START' sigl
          arg dd
          dderrmsg = 'OK'
          LRC = listdsi(dd "FILE")

```

```

/*****
/* Allow sysreason=3 & 22 to verify SYSOUT and tape DD statements */
/*****
    if LRC <> 0 & sysreason <> 3 & sysreason <> 22 then
        do
            dderrmsg = errmsg(sigl 'Required DD' dd 'is missing')
            call rcexit LRC dderrmsg sysmsglvl2
        end
        pull tracelvl . module . sigl . sparms
        call modtrace 'STOP' sigl
        interpret 'trace' tracelvl
        return
/***** @REFRESH END   DDCHECK   2004/11/09 22:48:36 *****/
/***** @REFRESH BEGIN DDLIST    2002/12/15 04:54:32 *****/
/* DDLIST - Returns number of DDs and populates DDLIST variable */
/*-----*/
/* N/A - None */
/*****
ddlist: module = 'DDLIST'
    if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
    parse arg sparms
    push trace() time('L') module 'From:' sigl 'Parms:' sparms
    call modtrace 'START' sigl
/*****
/* Trap the output from the LISTA STATUS command */
/*****
    call outtrap 'lines.'
    address TSO "LISTALC STATUS"
    call outtrap 'off'
    ddnum = 0
/*****
/* Parse out the DDNAMEs and concatenate into a list */
/*****
    ddlist = ''
    do ddl=1 to lines.0
        if words(lines.ddl) = 2 then
            do
                parse upper var lines.ddl ddname .
                ddlist = ddlist ddname
                ddnum = ddnum + 1
            end
        else
            do
                iterate
            end
        end
    end
/*****
/* Return the number of DDs */
/*****
    pull tracelvl . module . sigl . sparms

```

```

        call modtrace 'STOP' sigl
        interpret 'trace' tracelvl
        return ddnum
/***** @REFRESH END   DDLIST   2002/12/15 04:54:32 *****/
/***** @REFRESH BEGIN DDDSNS   2002/09/11 00:37:36 *****/
/* DDDSNS - Returns number of DSNs in a DD and populates DDDSNS */
/*-----*/
/* TARGDD - DD to return DSNs for */
/*****/
dddsns: module = 'DDDSNS'
        if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
        parse arg sparms
        push trace() time('L') module 'From:' sigl 'Parms:' sparms
        call modtrace 'START' sigl
        arg targdd
        if targdd = '' then call rcexit 77 'DD missing for DDDSNS'
/*****/
/* Trap the output from the LISTA STATUS command */
/*****/
        x = outtrap('lines.')
        address TSO "LISTALC STATUS"
        dsnum = 0
        ddname = '$DDNAME$'
/*****/
/* Parse out the DDNAMEs, locate the target DD and concatenate DSNs */
/*****/
        do ddd=1 to lines.0
            select
                when words(lines.ddd) = 1 & targdd = ddname &,
                    lines.ddd <> 'KEEP' then
                    dddsns = dddsns strip(lines.ddd)
                when words(lines.ddd) = 1 & strip(lines.ddd),
                    <> 'KEEP' then
                    dddsn.ddd = strip(lines.ddd)
                when words(lines.ddd) = 2 then
                    do
                        parse upper var lines.ddd ddname .
                        if targdd = ddname then
                            do
                                fdsn = ddd - 1
                                dddsns = lines.fdsn
                            end
                        end
                    end
                otherwise iterate
            end
        end
        end
/*****/
/* Get the last DD */
/*****/
        ddnum = ddlist()

```



```

        lastdd = word(ddlist,ddnum)
/*****
/* Remove the last DSN from the list if not the last DD or SYSEXEC */
/*****
        if targdd <> 'SYSEXEC' & targdd <> lastdd then
            do
                dsnum = words(dddns) - 1
                ddsns = subword(dddns,1,dsnum)
            end
/*****
/* Return the number of DSNs in the DD */
/*****
        pull tracelvl . module . sigl . sparms
        call modtrace 'STOP' sigl
        interpret 'trace' tracelvl
        return dsnum
/***** @REFRESH END   DDDSNS   2002/09/11 00:37:36 *****/
/***** @REFRESH BEGIN QDSN    2002/09/11 01:15:23 *****/
/* QDSN      - Make sure there are only one set of quotes */
/*-----*/
/* QDSN      - The DSN */
/*****
qdsn: module = 'QDSN'
        if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
        parse arg sparms
        push trace() time('L') module 'From:' sigl 'Parms:' sparms
        call modtrace 'START' sigl
        parse arg qdsn
        qdsn = ""strip(qdsn,"B","")""
        pull tracelvl . module . sigl . sparms
        call modtrace 'STOP' sigl
        interpret 'trace' tracelvl
        return qdsn
/***** @REFRESH END   QDSN      2002/09/11 01:15:23 *****/
/***** @REFRESH BEGIN UNIQDSN 2004/09/01 18:03:04 *****/
/* UNIQDSN   - Create a unique dataset name */
/*-----*/
/* PARM      - N/A */
/*****
uniqdsn: module = 'UNIQDSN'
        if wordpos(module,probe) <> 0 then trace 'r'; else trace 'n'
        parse arg sparms
        push trace() time('L') module 'From:' sigl 'Parms:' sparms
        call modtrace 'START' sigl
/*****
/* Compose a DSN using userid, exec name, job number, date and time */
/*****
        jnum = jobinfo(1) || jobinfo(2)
        udate = 'D'space(translate(date('0'),'','/'),0)
        utime = 'T'left(space(translate(time('L'),'',':'),0),7)

```

```

        uniqdsn = userid()'.execname'.jnum'.update'.utime
        if sysdsn(qdsn(uniqdsn)) = 'OK' then
            do
/*****
/* Wait 1 seconds to ensure a unique dataset (necessary on z990) */
/*****
                RC = syscalls('ON')
                address SYSCALL "SLEEP 1"
                RC = syscalls('OFF')
                uniqdsn = uniqdsn()
            end
/*****
                pull tracelvl . module . sigl . sparms
                call modtrace 'STOP' sigl
                interpret 'trace' tracelvl
                return uniqdsn
/***** @REFRESH END    UNIQDSN    2004/09/01 18:03:04 *****/

```

*Editor's note: this article will be concluded next month.*

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Computer Associates has announced new versions of Unicenter CA-Spool Print Management r11 (CA-Spool), Unicenter CA-View r11 (CA-View), and Unicenter CA-Deliver r11 (CA-Deliver).

CA-Spool provides management of all printing and spooling tasks throughout the enterprise. It features a Web interface and enhancements to critical processing points. Mainframe reports can be transformed to PDF, HTML, and RTF formats.

CA-Deliver is an online report management system that automates distribution, tracking, and printing. CA-View is an automated system that allows immediate online viewing of mainframe output.

For further information contact:  
URL: [www.ca.com](http://www.ca.com).

\*\*\*

Informatica has announced the latest version of PowerCenter for Mainframe, which provides a single data integration platform to help organizations access, transform, and integrate data from a large variety of systems and deliver that information to other transactional systems, real-time business processes, and people.

PowerCenter for Mainframe provides a unified data integration platform that leverages the business-critical aspects of mainframe computing while providing the same time-to-value and cost-of-ownership advantages that PowerCenter brings to non-mainframe environments, the company claims. These include PowerCenter's metadata-driven architecture for defining integration tasks once and deploying them anywhere on or off the mainframe, and Informatica PowerExchange's ability to access real-time, changed-only, and batch data.

For further information contact:  
URL: [www.informatica.com/products/powercenter/default.htm](http://www.informatica.com/products/powercenter/default.htm).

\*\*\*

IBM has announced an expanded suite of new Tivoli software products that automate real-time systems management and monitoring capabilities for eServer zSeries. The Tivoli OMEGAMON software suite allows organizations to proactively detect, isolate, and repair application performance problems anywhere within their IT infrastructure.

This enhanced product suite integrates Candle's technology, creating the Tivoli systems management solution for zSeries customers.

The IBM Tivoli OMEGAMON XE solution includes products that monitor and manage zSeries operating systems and subsystems including z/OS Unix System Services and Parallel Sysplex, z/VM, Unix System Services, CICS, DB2, IMS, zSeries networks, storage, WebSphere Application Server (WAS), WebSphere Integration Brokers, and WebSphere MQ.

For further information contact:  
URL: [www.ibm.com/ondemand](http://www.ibm.com/ondemand).

\*\*\*

StreamFoundry has announced its IT application suite, Central Management System (CMS) for problem, change, and request management.

The product reduces the overall processing requirements of the network, server, and database environments.

For further information contact:  
URL: [www.streamfoundry.com/product.html](http://www.streamfoundry.com/product.html).

