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MVS

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MVS Update

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Editor

Jaime Kaminski

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CA1 TMC information

During a recent change to our DASD layout it became necessary to move our CA1TMC (Tape Management Catalog) to another volume. However, when we tried to restart CA1, we received an 'IEFTMS70 5yy-132' message. This message is issued in situations where either the TMC or the audit file has been incorrectly defined. Although I have worked with CA1 for quite a while, I was unable to understand what was wrong and had to resort to calling CA for assistance. The answer they provided was sufficiently unexpected that it seemed worthwhile pointing out the problem to a wider audience. Apparently the message was issued because CA1 could not handle 31-bit UCBs. We had hit the problem by defining our new DASD with 31-bit UCBs rather than our old DASD's 24-bit ones.

Having assisted various people trying to understand the operation of CA1 and its use of the TMC, there would appear to be a lot of confusion and concern over this particular file. A lot of this confusion seems to stem from the TMC being an immovable file that requires 'specialist' procedures to move it.

The TMC is defined as a 340-byte fixed record, unblocked, immovable file. If you were to treat it as such and run a utility like SELCOPY to print out the file it would consist of the following structure.

The first three records are control records responsible for defining the volume ranges in use at your site (up to 52 independent ranges) and for controlling DSNB usage. Following these control records will be one volume record for every volume defined in your volume range. At the end of these records and marked by X'FF' in the first byte are the DSNB records.

If you look in the macro library supplied by CA, you will find two macros – TMMTMREC and TMMDSNB. TMMTMREC defines the layout of the control records and the volume records, while TMMDSNB defines the DSNB layout. Reading these Assembler macros should enable you to be able to see a correlation between the GRW (Generalized Report Writer) and the records. In other words you should be able to see the TMC as nothing more than a strangely unblocked ordinary file. Using the macro layouts for information will also permit the use of routines other than the GRW (or EARL) to read the TMC. So if you

lack the skills in either of these tools, you can simply use the maps as your guides for SELCOPY, etc. So if the TMC is just an ordinary file with an ordinary fixed format, why all the problem when it comes to moving it? There are two primary reasons:

- How the I/O is actually done to the TMC. When CA1 is started (or restarted) it takes note of where the TMC is. After this, access to the file is done through an SVC rather than through conventional I/O. This SVC works out the location of a TMC record and then 'jumps' to the disk location. This calculation of the offset from the start of the file is done by using the control records to work out where the currently in-use volume serial number is located and then adding in the number of control records. To make this clearer, assume you have two volume ranges, VOLSERs 10-19 and 30-49, and we are about to process volume 40. The jump is calculated as follows. Start of file +3records (for the control records) +10 records - VOLSER range 10-19 is 10 records -+11records (Volser 40-30 (start of second range) is 11 records). Hence, a jump to record 24 is required. As standard I/O is not done and is merely a relative record calculation, and because the start of file pointer is only obtained at start-up, should the file be moveable and, for example, defragged elsewhere, the I/O would continue to where the file was previously and not its new location.
- The TMC must not be in use while the move takes place. Any use of tape which could update the TMC while the move is happening would result in potential TMC damage. It is imperative therefore that tape activity is stopped while the move takes place.

Assuming you protect the TMC accordingly, you do not literaly need to follow the procedures as stated in the manual (not that I am recommending you do not follow them, of course). Utilities such as DFDSS are perfectly capable of safely moving a TMC as long as tape activity is controlled during the copy process and that CA1 is restarted after the move. It is the last point that matters most. You really can treat the TMC as an ordinary file and read and process it accordingly as long as CA1 is restarted to reset itself after the move and no tape activity has updated the TMC during your change activity.

Systems Programmer (UK)

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COBOL II scope delimiters

INTRODUCTION

Have you ever needed to get into a compound 'IF' phrase and wished you could end one 'IF' so you could begin another – without using a full stop? This is where scope delimiters can be very helpful.

USING COBOL II SCOPE DELIMITERS

Consider this example of a Y2K fix that would cause serious problems if you were stuck with COBOL I and could not use scope delimiters. In the example, dates are in the format YYMMDD; if there is a 'start' date, there will be a corresponding 'end' date; if there is no 'start' date, there will be no 'end' date:

IF INPUT-START-DATE > ZERO	
MOVE INPUT-START-DATE	TO WS-START-DATE-6
MOVE INPUT-END-DATE	TO WS-END-DATE-6
IF WS-START-DATE-6 < 500000	
MOVE 2Ø	TO WS-START-DATE-CC
ELSE	
MOVE 1Ø	TO WS-START-DATE-CC

Now we are in trouble. We want to repeat the last four lines of code for the end date without leaving the 'IF' sentence, since we only do all this if START-DATE > zero. We could repeat the 'IF' again, but that is wasteful. In COBOL II and above it can look like this:

```
IF INPUT-START-DATE > ZERO
    MOVE INPUT-START-DATE
                               TO WS-START-DATE-6
    MOVE INPUT-END-DATE
                               TO WS-END-DATE-6
    IF WS-START-DATE-6 < 500000
                               TO WS-START-DATE-CC
        MOVE 2Ø
    FLSE
                               TO WS-START-DATE-CC
        MOVE 10
        END-IF
        IF WS-END-DATE-6 < 500000
             MOVE 2Ø TO WS-END-DATE-CC
        ELSE
             MOVE 1Ø
                      TO WS-END-DATE-CC
        END-IF
    END-IF.
```

The last two 'END-IF's are superfluous, we could have put a full stop after the last 'WS-END-DATE-CC'. Some programmers include them for style and readability.

'END-IF' terminates the scope of the immediately preceding 'IF' without terminating the sentence. Just like 'ELSE' matches up with the immediately preceding 'IF'. Both require the same degree of care during use.

OTHER SCOPE DELIMITERS

There are a number of scope delimiters in COBOL II. A selection are shown below:

- END-ADD
- END-CALL
- END-DIVIDE
- END-EVALUATE
- END-MULTIPLY
- END-READ
- END-SUBTRACT
- END-WRITE.

We will consider some of these COBOL II scope delimiters in future issues of *MVS Update*.

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JES2 recovery

THE PROBLEM

Recently an unusual situation occurred at our site which had not happened for many years. One of the datasets in a JES2 procedure concatenation had become corrupted. This meant that the batch jobs using that concatenation were failing to find their JCL. Although we were able to repair the dataset, it needed to be redefined and therefore JES2 was unaware of the new 'correct' library. In common with most sites, an IPL to restore matters was not exactly the preferred option! So it was time to try something that I had not used since the days of MVS/ESA Version 3 Release 3 (currently we are at OS/390 Version 1 Release 3). Because it worked, the procedure may be of use should you encounter the problem at your site.

A SOLUTION

The command £P JES2, ABEND was issued. This forces JES2 to stop and issue a variety of messages, ending with a WTOR. By replying END to the WTOR, and then issuing a S JES2, WARM, NOREQ it was possible to get JES2 to restart as a 'hot' start. This allowed JES2 to find its PROCLIB concatenation again, and (the best bit) nothing stopped or failed in the process.

Should you need to try this, an extract from the log to show all the messages obtained is shown below. This is followed by a description for what a 'hot' start 'officially' means in JES2 terms. I have included the log extract below, because the messages provided can look decidedly worrying if you are not ready for them!

```
£P JES2,ABEND
*£HASPØ95 JES2 CATASTROPHIC ERROR. CODE = £PJ2
£HASPØ88 JES2 ABEND ANALYSIS 842
£HASPØ88 FMID = HJE66Ø3 LOAD MODULE = HASJES2Ø
£HASPØ88 SUBSYS = JES2 OS 1.3.Ø
£HASPØ88 DATE = 99.217 TIME = 19.27.11
£HASPØ88 DESC = OPERATOR ISSUED £PJES2, ABEND
£HASPØ88 MODULE MODULE OFFSET SERVICE ROUTINE
£HASPØ88 NAME BASE + OF CALL LEVEL CALLED
```

£HASPØ88 £HASPØ88 HASPCOMM ØØØ7D828 + ØØEB56 OW21918 *ERROR £PJ2 $\pounds HASPØ88 PCE = COMM (ØCA361BØ)$ \pounds HASPØ88 RØ = ØØØ8C15Ø ØØBF4738 ØØØØ522E ØØØ83F44 £HASPØ88 R4 = ØØØ83Ø2C ØCA36B2C ØØØØØØØ4 ØCA36B3Ø £HASPØ88 R8 = ØØØ8C15Ø ØØØ8D15Ø ØØØØØØØØ ØØØØ6ØØØ £HASPØ88 R12 = ØØØ7D878 ØCA361BØ ØCAØ5CCØ ØØØD37F8 £HASPØ88 -*£HASP198 REPLY TO £HASPØ98 WITH ONE OF THE FOLLOWING: 843 FND - STANDARD ABNORMAL END END.NOHOTSTART - ABBREVIATED ABNORMAL END (HOT-START IS AT RISK) SNAP - RE-DISPLAY £HASPØ88 DUMP - REQUEST SYSTEM DUMP (WITH AN OPTIONAL TITLE) *71 £HASPØ98 ENTER TERMINATION OPTION *72 £HASP426 SPECIFY OPTIONS - JES2 OS 1.3.0 GSVXØ19I JES2 is not active R 72, WARM, NOREQ IEE6ØØI REPLY TO 72 IS;WARM,NOREQ IEF196I IEF285I SYS1.PARMLIB KEPT IEF196I IEF285I VOL SER NOS= SY1ØA1. IXZØØØ1I CONNECTION TO JESXCF COMPONENT ESTABLISHED, 973 GROUP MVSJESP1 MEMBER MVSJESP1£IP01 £HASP537 THE STATEMENTS IN THE INITIALIZATION DECK REQUIRE A 974 CHECKPOINT SIZE OF 431 4K RECORDS IEF196I IEF237I 108A ALLOCATED TO SYS00001 £HASP478 INITIAL CHECKPOINT READ IS FROM CKPT1 976 (SYS1.HASPCKPT ON JESØØ1) LAST WRITTEN THURSDAY, 5 AUG 1999 AT 19:27:08 (LOCAL TIME) *£HASP493 JES2 MEMBER-IP01 HOT START IS IN PROGRESS £HASP492 JES2 MEMBER-IP01 HOT START HAS COMPLETED

HOT START

The following is an extract from the JES2 commands manual that provides a definition of a hot start:

Hot start: a hot start is a warm start of an abnormally terminated JES2 member without an intervening IPL. JES2 performs a hot start when a particular JES2 member has stopped but other systems have continued to function and have not experienced problems. When JES2 hot starts, all address spaces continue to execute as if JES2 had never terminated. Jobs that were processing on output devices are re-queued as if a £I command had been issued. Jobs on input devices must be resubmitted and lines must be restarted. Hot starts have no effect on other members in a MAS configuration.

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Converting a Julian date to Gregorian

INTRODUCTION

In *MVS Update*, Issue 156, September 1999, there was a simple utility to execute job steps based on the day of the week. The utility worked by extracting the system date in Julian format, converting this date into Gregorian format and then calculating the weekday using Zeller's Congruence.

However, there is no need to convert a Julian date to Gregorian format before finding the day of the week. The code below does it directly:

DAYOWEEK	CSECT		
	SAVE	(14,2),,*	
	LR	R2,R15	
	USING	DAYOWEEK,R2	
	LA	R14,SAVEAREA	
	ST	R13,4(,R14)	
	ST	R14,8(,R13)	
	LR	R13,R14	
	TIME	LINKAGE - SVC	GET CURRENT DATE
	ST	R1,WORK1+4	STORE DATE IN ØCYYDDDF FORMAT
	DP	WORK1,=PL3'1000'	SPLIT YEAR AND DAY
	ZAP	WORK2,WORK1(5)	PICK UP YEAR
	CVB	R15,WORK2	YEAR IN BINARY
	М	R14,=F'5'	MULTIPLY BY 5
	BCTR	R15,Ø	SUBTRACT 1
		R15,2	DIVIDE BY 4
		WORK2,WORK1+5(3)	PICK UP DAY
		R14,WORK2	DAY IN BINARY
		R15,R14	ADD TO "YEAR"
		R14,R14	ZEROISE EVEN REGISTER
		R14,=F'7'	GET DAY OF WEEK
			PUT RESULT IN REGISTER 15
			RESTORE REGISTER 13
	RETUR	N (14,2),RC=(15)	Ø=SUN, 1=MON, 2=TUE, 6=SAT
*			
SAVEAREA			
WORK1			WORK AREA
WORK2			WORK AREA
	LTORG		
	YREGS		
	END	DAYOWEEK	

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PDF line commands

INTRODUCTION

In edit mode, you can enter line commands in the left margin (the line number area). You probably already know the most common ones such as 'C'opy, 'A'fter, 'B'efore, 'M'ove, 'D'elete, 'I'nsert, 'R'epeat, etc, and their block equivalents. There are others such as: 'O'verlay, ')' shift right, '('shift left, as well as 'TS' text split, and 'TF' text flow.

LINE COMMANDS

Here are a few of the less common tools available:

- 'X' will temporarily hide (or e'X'clude) a line. 'XX' is the group equivalent. This is handy to get some intervening lines out of the way so that you can get two groups of lines closer together, perhaps on the same screen. The line or lines are replaced with a single line of dashes. You can enter the following commands in the left margin of that line (If 'n' is not specified, it automatically defaults to '1').
- 'Fn' will redisplay the first 'n' lines of excluded text.
- 'Ln' redisplays the last 'n' lines.
- 'Sn' redisplay 'n' lines with the leftmost indentation in a block of excluded lines.
- 'UC' will convert all the characters on a line to upper case. 'UCC' is the group command. 'LC' and 'LCC' are the 'lower-case' equivalents.

Information about the features that involve a 'mask' can be found in the information in HELP (PF1).

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INTRODUCTION

24 September 1999 saw the general availability of OS/390 Version 2 Release 8. The emphasis of the new release is on better integration and availability. Central to the upgrade are security and systems management features, including the ability to manage virtual private network encryption keys dynamically through the Internet Key Exchange, enhanced management and administration of digital certificates used by both server applications and end-users, higher availability of TCP/IP in a Parallel Sysplex, ISPF customization, and the capability to print from ERP and Internet-related applications.

NEW SECURITY FEATURES

The following new security features are particularly important and will enhance the role of OS/390 in supporting e-commerce:

- IPSec VPN provides a secure pathway between OS/390 and other IPSec VPN-capable systems, routers, and firewalls through encryption using the System/390 hardware CMOS Cryptographic Coprocessor.
- The exchange of encryption keys between the end-points of IPSec VPN can now be automated and dynamically managed through Internet Key Exchange (IKE), an IPSec protocol for cryptographic key and security management.
- The centralized management of digital certificates belonging to server applications and their related private encryption keys, is another new feature. This will allow customers and application developers to provide common secure management of these certificates as well as the chain of trust needed to verify user certificates presented to these applications.
- There is SSL client authentication to the TN3270 server, allowing TCP/IP clients to access customer applications traditionally only accessible from a 3270 screen. It also secures against unauthorized access to SNA applications from TCP/IP users.

LDAP DIRECECTORY ACCESS

The LDAP (Lightweight Directory Access Protocol) server has been enhanced to support LDAP Version 3 protocol, enabling OS/390 LDAP Server to interoperate with other LDAP Version 3 clients and servers. LDAP on OS/390 includes Java support, LDAP access to RACF information, and LDAP client authentication using RACF. It also supports the Secure Sockets Layer for encrypted privacy of communication and it supports multiple LDAP servers on multiple systems in a Parallel Sysplex.

VIRTUAL IP ADDRESSING

The SecureWay Communications Server for OS/390 provides Virtual IP Addressing (VIPA) Takeover, allowing real IP addresses to be associated with a pseudo address, assigned to an end-user in the System/390 server. If a connection fails, traffic is routed to an alternative connection associated with the same VIPA.

PRINTER CONSOLIDATION

Another new feature is that Infoprint Server will use the Internet printing protocol to process print jobs over the Internet securely. Infoprint will use datastream transforms to translate data from one printer format to another to allow printing from popular PC and workstation applications as well as many ERP applications. It can convert PCL, Postscript, and PDF files and is being positioned to consolidate enterprise print serving functions around the System/390.

PARALLEL SYSPLEX CONFIGURATION

Parallel Sysplex configuration can be complex but the use of the Parallel Sysplex configuration tool allows a Parallel Sysplex configuration to be created interactively.

ISPF CUSTOMIZATION

One of the problems of installing new releases of OS/390 has been that ISPF customization needs to be redone every time. Normally this involves using the ISRCONFG Assembler resource, and manually reassembling and relinking to make changes. With the introduction of the 'ISPF configuration utility' this problem is alleviated. Upon

entering TSO ISPCCONF, it is possible to convert the old ISRCONFG Assembler source into a keyword file that includes all the options or all those changed from the defaults.

OTHER ENHANCEMENTS

Other enhancements to OS/390 Version 2 Release 8 include:

- Allowing SNA users to use Triple DES encryption. Service policy enhancements are said to improve the capability to monitor and manage network performance to service level agreements.
- The addition of a dynamic update feature to the policy agent allows service policies to be implemented without impacting network availability.
- A new resource reservation protocol support allows users to invoke reservation services, reserve bandwidth, and classify reservations through an RSVP API.
- A new SNMP SLA subagent that enables network administrators to retrieve data and determine whether the current set of SLA policy definitions are performing as desired.
- Improvements in Workload Manager (WLM) that allow it to prioritize workloads at the request level.

ANALYSIS

As with other recent releases of OS/390, the emphasis on security implies the operating system is being aggressively promoted for electronic commerce applications. The emphasis on printer consolidation, and Parallel Sysplex configuration again re-emphasizes the importance of enterprise server centralization.

OS/390 VERSION 2 RELEASE 9

In Version 2 Release 9, IBM plans to make further improvements to native file and print serving for Windows clients, text search support for XML documents and unicode, and additional Unix system services functions. We can also look forward to a new naming convention for OS/390 in the new millennium.

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Selectively blocking commands

THE PROBLEM

It is sometimes necessary to block the issue of certain MVS commands. You may for instance have a job that you do not want CANCELLed or even FORCEed without permission from the systems programming group. There are also commands that can have an impact on system performance (eg doing an LLA refresh on a production system). Although this can be controlled through RACF profiles, it is a somewhat cumbersome process and a mistake could easily end up with the wrong groups being blocked or allowed to issue the command(s).

A SOLUTION

A much easier way would be to have a simple facility available that will allow systems programmers to block commands by putting them into and removing them from a PARMLIB member.

The following facility does exactly that. There are two programs – an MVS command exit, and a routine to manage blocked commands. The MVS command exit intercepts commands of the format 'T NOCMD=xx' with xx the suffix of a PARMLIB member NOCMDxx.

This command is passed to the second routine that manages the blocked commands. This routine looks for the member NOCMDxx in the PARMLIB concatenation and, if the member is found, a CSA table is GETMAINed and the commands are copied from the member into the table. The address of the CSA table is anchored off a subsystem entry by the name of NCMD. (This entry is dynamically added when the first 'T NOCMD=xx' command is entered.) The command exit does a check every time that a command is entered. If the name of the subsystem exists, the command is compared with all the entries in the CSA table. If a match is found, the command is not executed and a message is given. To disable the command blocking, simply issue a 'T NOCMD=NO' command (which of course can also be put into the block table).

Note that the commands are added into the PARMLIB member without blanks, eg to block a LLA refresh the PARMLIB member should have 'FLLA', on a line of its own. The comparison with the entered command is based on the length of the command in the PARMLIB member. For example, CANCELABC will block the command 'CANCEL ABCD' but it will not block the command 'CANCEL A'. The program is not written to accept wild cards (eg CANCEL AB*D) but it can easily be modified to accept this format.

The MVS command exit also contains samples on how to intercept other commands and will be very easy to convert to suit your purposes. It removes the blanks from the passed commands and then calls the associated entry points for each matched command. The field CMDWAREA contains the de-blanked command whilst CMDX@ contains the address of the original MVS command buffer. It is installed by adding it into MPF refer to the *OS/390 Initialization Guide* for a description on how to do this very easy task.

MVS COMMAND EXIT

```
* MVS Command EXIT.
* This module will recognize commands by comparing them to entries in
* a table. It will than do a BAS to an entry point specified for the
* command. If the processing is successful, it will either suppress the
* command from further MVS processing or pass it to MVS, depending on
* a field specified per command in the command table.
* An ESTAE routine is set up to intercept any ABENDs that may result in
* calling the subroutines. The ESTAE routine issues an SVC dump.
* To add another command intercept:
* 1) Do a "F COMMTAB" in column 1. This is the command table and an
    entry can be added The format is:
        DCAL2(L'XX).Length of command textDCAL4(ENTRYPT).Address of processing routineDCAL1(flag).Flag to control MVS processingDCC'MYCOMMND'.Command text (spaces removed)
*
*
*
*
   ΧХ
* 2) The code to process command MYCOMMND should be coded at ENTRYPT.
* 3) Assemble and link.
* 4) Refresh LLA
* 5) Do a "T MPF=xx"
* Module type : Reentrant, called in supervisor state key \emptyset
* Addressing mode: AMODE 31, RMODE any
* Register usage : R12= Base register
                    R13= Pointer to general savearea and workareas
MVSCOMEX CSECT
MVSCOMEX AMODE 31
```

MVSCOMEX RMODE ANY BAKR R14.Ø .Save caller's status LR R12.R15 USING MVSCOMEX.R12 Main driver routine 1 R4.0(R1).Get CMDX Address USING CMDX.R4 STORAGE R3.STORSIZE LA STORAGE OBTAIN, LENGTH=(3), LOC=ANY, SP=229 R3.STORSIZE LA I R R2.R1 .Point to getmained area R9.R9 XR MVCL R2.R8 .Propagate binary zeros USING STORAREA.R1 ST R13.SAVEAREA+4 .Backchain DROP R1 LR R13.R1 .Addressability to getmained area USING STORAREA,R13 R4.CMDX@ .Preserve adress of command buffer ST BAS R14.FINDCMND .Locate command text, remove blanks CHKPRCS BAS R14.CMNDCOMP .See if we have to process command BAS R14.SEEKSSI .Go see if command blocking active LTR .Is it? R15.R15 CHKOURS .No, go see if we have to intercept ΒZ .Compare command with "block" table BAS R14.TSTBLOCK .Must it be blocked? LTR R15.R15 .No, don't block the command ΒZ CHKOURS 'MVSCOMEX(W): This command has been blocked by Software X WTO Support', ROUTCDE=11 CHKOURS ТМ DOFLAG.YES .Is this an intercepted command? BNO .No. not one of ours RETURN BAS R14, PROCCMND .Process command RETURN L R4,RETCODE .Pick up return code LA R3.STORSIZE LR R2.R13 .Pointer to storage area STORAGE RELEASE,LENGTH=(3),SP=229,ADDR=(2) LR R15.R4 .Reload return code ΡR .Return to our caller * This routine locates the command and removes all blanks FINDCMND BAKR R14.0 R4,CMDXCLIP .Get command buffer address L DROP R4 USING CMDXCLIB.R4 .Access the passed command buffer .MVS buffer of actual command text R1.CMDXCMDL LA ST R1.MVSCOM@ .Preserve this address LH R1,CMDXCMDL .Length of command LA R2.CMDXCMDI .Access start of text POINTCMD LA R3.CMDWAREA .Point to start of command wrk area

XR R5.R5 .Clear length counter DESPACEL EQU * .Remove all spaces from the command Ø(R2).C'' CLT .Is it a space? ΒF BUMPUP .Yes CLI Ø(R2),X'ØØ' .Is it low value? ΒE BUMPUP .Yes CLM R5,3,=AL2(L'CMDWAREA) Too long? BL MOVECHAR .No WTO 'MVSCOMEX(E): -Command exceeds max length', ROUTCDE=13 В UPPERCSE .Only accept these characters MOVECHAR MVC $\emptyset(1,R3),\emptyset(R2)$.Move character into wrk area .Point to next empty space R3.1(R3) LA ΙA R5.1(R5) .Bump up length counter by 1 BUMPUP LA R2.1(R2) .Bump up pointer to cmd buffer вст R1.DESPACEL .Do for each character in the cmd CMDWAREA(L'CMDWAREA), SPACES .Uppercase UPPERCSE OC ST R5.CMDLENG .Length of de-spaced command FINDCMNX PR .Reload our return address This routine gets the address of our subsystem entry SEEKSSI BAKR R14.0 .Start of CVT L R4.16 USING CVT.R4 .Establish addressability R4,CVTJESCT .Get JESCT address 1 USING JESCT,R4 .Establish addressability .Get SSCT chain address R4.JESSSCT 1 USING SSCT.R4 .Set addressability SSCTLOOP EQU .Look for our SSCVT * LTR R4.R4 .End of chain? .Yes - not found ΒZ NOBLOCK CIC SSCTSNAM, SUBSYSN .Our SSCVT? ΒF .Gotcha FOUNDSSI R4.SSCTSCTA .Point to next SSCVT entry 1 В SSCTLOOP .Redo the loop FOUNDSSI EQU * .We have found the subsystem SSCTSUSE,=F'Ø' CLC .Has the CSA table been obtained? .No, we are not blocking commands B7 NOBLOCK .Return address of the CSA table L R15.SSCTSUSE .Store the address ST R15,BLOCKTB@ В SEEKSSIX .Get out R15.R15 NOBLOCK XR .No, we are not blocking commands SEEKSSIX PR * This routine checks if the command matches an entry in the "block" table kept in CSA. TSTBLOCK BAKR R14.Ø .Where the CSA "block" table is L R7,BLOCKTB@ .Number of entries in "block" table L R6,Ø(R7) LTR R6.R6 .Any commands in the table? ALLOWCMD .No commands in "block" table ΒZ

	LA	R7,4(R7)	.Start of first command text
CARDLOOP		*	Longth of a Ublack and output
	LA LR	R9,8Ø R1,R7	.Length of a "block" cmd entry .Point to start of "block" cmd
	XR	R8,R8	.Clear length counter
LENGLOOP		Ø(R1),X'4Ø'	.Look for a space
LENGLOOP	BE	SETLENG	.Command must match this length
	LA	R1,1(R1)	.Point to next character
	LA	R8,1(R8)	.Bump up length counter by 1
	ВСТ	R9,LENGLOOP	.Scan until we find a blank or EOC
SETLENG	L	R3,CMDLENG	.Length of de-spaced command
	С	R3,R8	.Command buffer too short?
	BL	ALLOWCMD	.Yes, do not block the command
	LR	R5,R8	.Length to compare
	LR	R3,R8	.Length to compare
	LR LA	R2,R7	.Point to start of this "block" cmd
	CLCL	R4,CMDWAREA R2,R4	.Point to start of passed command .Does command match a table entry?
	BE	BLOCKCMD	.Yes, block the command
	LA	R7,8Ø(R7)	.Point to the next card
	BCT	R6,CARDLOOP	.Compare with each table entry
ALLOWCMD	XR	R15,R15	.This command not blocked
	В	TSTBLOCX	.Get out
BLOCKCMD		R15,NOMVS	.Do not let this through to MVS
	ST	R15,RETCODE	.Plug value into return code field
TSTBLOCX	PR		
*******	******	*****	*****
*******			**************************************
*	This	routine compares the c	**************************************
*	This :	routine compares the c	command to entries in the table
* ******	This :	routine compares the c	command to entries in the table
* ******	This ****** BAKR LA DROP	routine compares the c ************************************	command to entries in the table ************************************
* ******	This BAKR LA DROP USING	routine compares the c ************************************	command to entries in the table
* ******** CMNDCOMP	This BAKR LA DROP USING LA	routine compares the c ************************************	command to entries in the table
* ******	This BAKR LA DROP USING LA EQU	routine compares the c ************************************	command to entries in the table .Point to start of command table .Addressability to command table .Point to end of command table .Compare to all commands in table
* ******** CMNDCOMP	This BAKR LA DROP USING LA EQU ICM	routine compares the c ************************************	command to entries in the table ************************************
* ******** CMNDCOMP	This BAKR LA DROP USING LA EQU ICM STH	routine compares the c ************************************	command to entries in the table ************************************
* ******** CMNDCOMP	This BAKR LA DROP USING LA EQU ICM STH BCTR	routine compares the c ************************************	command to entries in the table . Point to start of command table . Addressability to command table . Point to end of command table . Compare to all commands in table . Pick up command length . Store into parmlist . Reduce length by 1
* ******** CMNDCOMP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX	routine compares the c ************************************	command to entries in the table .Point to start of command table .Addressability to command table .Point to end of command table .Compare to all commands in table .Pick up command length .Store into parmlist .Reduce length by 1 .See if there is a match
* ******** CMNDCOMP	This BAKR LA DROP USING LA EQU ICM STH BCTR	routine compares the c ************************************	command to entries in the table . Point to start of command table . Addressability to command table . Point to end of command table . Compare to all commands in table . Pick up command length . Store into parmlist . Reduce length by 1
* ******** CMNDCOMP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B	routine compares the c ************************************	command to entries in the table .Point to start of command table .Addressability to command table .Point to end of command table .Compare to all commands in table .Pick up command length .Store into parmlist .Reduce length by 1 .See if there is a match
* CMNDCOMP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B	routine compares the c ************************************	command to entries in the table ************************************
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B CLC XR ICM	routine compares the c ************************************	command to entries in the table ************************************
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B CLC XR ICM AR	routine compares the c ************************************	command to entries in the table ************************************
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B CLC XR ICM AR LA	routine compares the c ************************************	command to entries in the table .Point to start of command table .Addressability to command table .Point to end of command table .Compare to all commands in table .Pick up command length .Store into parmlist .Reduce length by 1 .See if there is a match .Yes, there is .Compare command to table entry .Pick up the length of this entry .Bump up by the text length .Bump up by header length
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B CLC XR ICM AR LA CR	routine compares the c ************************************	command to entries in the table .Point to start of command table .Addressability to command table .Point to end of command table .Compare to all commands in table .Pick up command length .Store into parmlist .Reduce length by 1 .See if there is a match .Yes, there is .Compare command to table entry .Pick up the length of this entry .Bump up by the text length .Bump up by header length .End of table?
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B CLC XR ICM AR LA CR BNL	routine compares the c ************************************	command to entries in the table ************************************
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B CLC XR ICM AR LA CR BNL B	routine compares the c ************************************	command to entries in the table .Point to start of command table .Addressability to command table .Point to end of command table .Compare to all commands in table .Pick up command length .Store into parmlist .Reduce length by 1 .See if there is a match .Yes, there is .Compare command to table entry .Pick up the length of this entry .Bump up by the text length .Bump up by header length .End of table?
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE CLC XR ICM AR LA CR BNL B EQU	routine compares the c ************************************	<pre>command to entries in the table</pre>
* CMNDCOMP COMPLOOP	This BAKR LA DROP USING LA EQU ICM STH BCTR EX BE B CLC XR ICM AR LA CR BNL B	routine compares the c ************************************	command to entries in the table .Point to start of command table .Addressability to command table .Point to end of command table .Compare to all commands in table .Pick up command length .Store into parmlist .Reduce length by 1 .See if there is a match .Yes, there is .Compare command to table entry .Pick up the length of this entry .Bump up by the text length .Bump up by header length .End of table? .Yes, match not found. Get out .Redo for each table entry .Address of matched table entry

0 T DOFLAG.YES .Yes. it must be processed by us CMNDCOMX PR .Return to our caller ****** This routine sets up the ESTAE and calls the correct EP PROCCMND BAKR R14.Ø LA R14.PROCCMNX .Start from here after an abend STM R1,R14,RUBLSRGS .Preserve all our registers RUBLIST,=B'Ø1111111111110' Regs 1-14 to be reloaded MVC .Point to ESTAE storage area LA R1.ESTAESTR MVC Ø(ESTAEL1,R1),ESTAEMAC R13,ESTAEPRM ST .R13 required by recovery routine LA R2,ESTAEPRM .Pass R13 contents as a parameter ESTAE MF=(E,(1)), PARAM=((2))R4.MATCH@ .Address of matching cmd tble ent L .Entry point of routine to call ICM R15,15,COMMEP BASR R14.R15 .Call command processing routine CLC RETCODE,=F'Ø' .Was the processing successful? BNZ PROCCMNX .No, pass the command to MVS XR R15.R15 IC .What MVS should do when we're done R15.COMMFLAG СН R15,=AL2(TODECIDE) .Dynamic decision on what to do? BNE PROCCMNX .No, do as specified in table 1 R15,DECISION .Pick up our decision .This is what we pass back to MVS PROCCMNX ST R15,RETCODE LA R1.ESTAESTR .Point to ESTAE storage area MVC Ø(ESTAEL2,R1),ESTAEDEL ESTAE MF=(E,(1)).Remove ESTAE PR .Return to our caller EXAMPLE: This routine processes the "D TS,L" command ROUTINEA BAKR R14.0 WTL '"D TS.L" Command detected' Add code in here to do processing for "D TS,L"... R15.RetCode ST PR EXAMPLE: This routine processes the "CANCEL" command ROUTINEB BAKR R14,Ø WTL '"CANCEL" Command detected' ST R15.RetCode * Add code in here to do processing for "CANCEL"... Set field DECISION depending on whether the command is acceptable * and should also be passed to MVS. PR This routine handles the "T NOCMD=' command * TNOCMD EOU BAKR R14,Ø .Preserve our registers R2.CMDWAREA LA .Point to (compressed) command

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.Plug the address into parmarea ST R2.CALLPRM LA R1,CALLPRM .Address of call parms ST R1.CALLPRM@ .Plug into pointer R1.CALLPRM@ .Point to parameter area ΙA .Call the routine to do the work I TNK FP=TFI NOCMD TNOCMDX ST R15.RETCODE PR Constants follow COMMTAB DS ØF .Add new commands in here COMADESC DC AL2(L'VERBA) .Length of command text DC AL4(ROUTINEA) .Address of processing routine DC .Flag: pass command to MVS AL1(TOMVS) .Command text without blanks VERBA DC C'DTS.L' COMBDESC DC AL2(L'VERBB) .Length of command text .Address of processing routine DC AL4(ROUTINEB) DC AL1(TODECIDE) .Flag: do not pass command to MVS DC C'CANCEL' .Command text without blanks VFRRR COMCDESC DC AL2(L'VERBC) .Length of command text DC AL4(TNOCMD) .Address of processing routine DC AL1(NOMVS) .Flag: do not pass command to MVS VFRBC DC C'TNOCMD=' .Command text without blanks * If the called entry point for a specific command returns a non-zero return code, the command is always also passed to MVS, as we failed. If the entry point returns a zero return code (in R15), a decision * * as to what should be done is based on the third field for the entry in the table, which must always be one of the following. "TODECIDE", means that whether to pass the command to MVS * or not is to be decided. A further field called "Decision" is then * scanned. If this field contains the value $\emptyset \emptyset$, the command is also ENDOFTAB EQU * TOMVS EQU X'ØØ' .Pass the command to MVS as well NOMVS FOU X'Ø4' .Do not pass the command to MVS TODECIDE EOU X'Ø1' .Decision to be made on passing COMMTABL EQU *-COMMTAB CL(L'CMDWAREA)' ' SPACES DC ESTAEMAC ESTAE RECOVER, PARAM=ESTAEMAC, ASYNCH=NO, MF=L ESTAEL1 EOU *-ESTAEMAC ESTAEDEL ESTAE Ø,MF=L ESTAEL2 EQU *-ESTAEDEL SUBSYSN DC CL4'NCMD' .Name of subsystem with address of LTORG * This routine is the ESTAE error recovery routine DS ØF * RECOVER EQU

IR R12,R15 .Load our current address DROP R12 USING RECOVER.R12 RØ,=H'12' .Make sure SDWAREA is available СН BNE SDWAVAIL .SDWA is not available NOTAVAIL XR R15,R15 RR R14 .Go Back to MVS (percolate) SDWAVAIL EOU * .Remember SDWA address USING SDWA.R1 LR R3.R1 .Preserve R1 L R1,SDWAPARM .Address of passed parm L R13,Ø(1) .Reload pointer to workarea .Pointer to SDWA ST RØ.SDWASTOR ST R14.ESTAER14 .Our return address MVC SDMPAREA(SDUMPLEN), SDUMPMAC LA R1.SDMPAREA SDUMP MF=(E,(1))R2.RUBLSRGS+52 .Retry address 1 R1.R3 LR .Restore register 1 R3.RUBLIST LA SETRP RC=4,RETADDR=(2),RETREGS=YES,RUB=(3),DUMP=YES R14.ESTAER14 .Reload our return address L ΒR R14 .Back to MVS SDUMPMAC SDUMPX HDR='MVSCOMEX Command exit routine ABEND'. Х SDATA=(RGN,SUM),MF=L SDUMPLEN EQU *-SDUMPMAC LTORG DSECTs follow STORAREA DSECT SAVEAREA DS 18F .General savearea CMDWAREA DS CL2ØØ .Command moved in here CMDTXTLN DS .Length of (compressed) command н RETCODE DS .SUBTASK's TCB F SDMPAREA DS CL(SDUMPLEN) .SDUMP macro area BLOCKTB@ DS F .Address of NCMD subsystem F DECISION DS .(dynamic) ret code to pass to MVS CMDLENG DS F .Length of command without blanks F .Address of matching cmdtble entry MATCH@ DS RUBSTART DS ØF DS Н RUBLIST DS Н RUBLSRGS DS 15F .Local registers used by ESTAE ESTAEPRM DS F .Parameter passed to ESTAE routine SDWASTOR DS F .Address of passed SDWA .Back-to-MVS address for ESTAE ESTAER14 DS F .ESTAE macro area ESTAESTR DS CL(ESTAEL1) * The folowing two fields are passed to the called subroutine. CALLPRM@ DS .Address of routine specific parms F CMDX@ DS F .Address of CMDX as received *

MVSCOM@ DOFLAG YES NO CALLPRM STOPSIZE		F C X'80' X'00' ØF	.Address (in MVS) of command text .Do we process command?
CALLPRM STORSIZE COMMDSCT COMMLENG COMMEP COMMFLAG COMMTEXT RØ R1 R2 R3 R4 R5 R6 R7 R8 R9 R1Ø R11 R12 R13 R14 R12 R13 R14 R15	EQU DSECT DS DS EQU DS EQU EQU EQU EQU EQU EQU EQU EQU EQU EQU	*-STORAREA CL2 AL4 C *-COMMDSCT ØC Ø 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 101	.Command table DSECT .Length of text to scan for ."entrypoint" to branch to .MVS to process/ ignore instruction .Length of fixed part .Variable length text to scan for
	IHASD CVT D IEFJE IEFJS END	SECT=YES SCT	

BLOCKCMD SOURCE

This routine blocks certain commands from being entered.

```
BLOCKCMD CSECT
BLOCKCMD AMODE 31
BLOCKCMD RMODE 24
                  .Save caller's status
      BAKR R14,Ø
      LR R12,R15
     USING BLOCKCMD.12
*
     Main driver routine
R4,Ø(R1)
R4,Ø(R4)
                       .Ptr to compressed command text addr
     L
LOAD
      L
                      .Point to compressed command text
                       .Ptr to command exit info area
      L
         R5,4(R1)
      USING CMDX,R5
                       .Addressability to command info
```

0700405	LA A	R3,STORSIZE R3,BUFFSIZE	.Our requirement .Add length of the PARMLIB buffer
STORAGE	LR	GE UBTAIN, SP=229, LENG R2, R1	GTH=(3),LOC=BELOW,KEY=8 .Point to getmained area
	LA	R3,STORSIZE	. Torne to geemanned area
	XR	R9,R9	
	MVCL	R2,R8	.Propagate binary zeroes
		STORAREA,R1	
	ST DROP	R13,SAVEAREA+4	.Back chain
	LR	R13,R1	
		STORAREA, R13	.Addressability to getmained area
	MVC	CART, CMDXCART	.Pick up the CART
	MVC	FROMCNID,CMDXC4ID	.Where the command originated (Id)
	MVC	FROMSYS,CMDXISYN	.Where the command originated (Sys)
	L	R5,CMDXCLIP	.Get command buffer address
	DROP	R5	
000FTCFV		CMDXCLIB,R5	Access the passed command buffer
GOGETSFX		R14,GETSUFIX R15,R15	.Locate the data set name in the text .Did we get the suffix?
	LTR BNZ	RETURN	.No, get out
	CLC	SUFFIX,=C'NO'	.Must we turn blocking off?
	BNE	GOREADIT	.No
	B	GOGETSSI	.Go get SSI address
GOREADIT	BAS	R14, READPMEM	.Go read the PARMLIB member
	LTR	R15,R15	.Did we get the suffix?
	BNZ	RETURN	.No, get out
GOGETSSI	BAS	R14,GETSSI@	.Go get the subsystem address
	LTR	R15,R15	.Did we get the address?
	BNZ	EXITMSGS	.No, get out
GOGETCSA		R14,GETCSA@	.Go get the CSA table address
		R15,R15	.Did we get the address?
	BNZ BAS	EXITMSGS R14,COPYCRDS	No, get out
EXITMSGS		SUFFIX,=C'NO'	.Go copy cards into CSA table .Did we turn blocking off?
LATINJUJ	BNE	CHEMPTY1	.No
	L	R8,FROMCNID	.Where command came from
	LA	R9, FROMSYS	.Where command came from
	WTO		and blocking has been turned off', X
		CONSID=(8),SYSNAME=	(9),CART=CART,ROUTCDE=11
	В	RETURN	
CHEMPTY1	ТМ	EMPTY,YES	.Is the PARMLIB member empty?
	BO	RETURN	.Yes, get out
	L	R8, FROMCNID	.Where command came from
	LA	R9, FROMSYS	.Where command came from
	WTO		and blocking has been activated', X (9),CART=CART,ROUTCDE=11
RETURN	EQU	*	.Pick up return code
	LUU	R4,RETCODE	.Pick up return code
	LA	R3,STORSIZE	.Size of area to free
	A	R3,BUFFSIZE	.Add length of the PARMLIB buffer
	LR	R2,R13	.Address of area to free
	STORA		3),ADDR=(R2),SP=229,KEY=8

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LR R15.R4 .Copy return code TOCALLER PR .=>Caller ØD .Align ٦C EJECT This routine locates the member suffix in the text GETSUEIX BAKR R14.0 USING CMDXCLIP,CMDXCLIB .Map to command length and text CLC CMDXCMDL,=H'9' .Length must be at least 9 bytes ΒL INVLSNTX .Invalid syntax 9(R4).X'4Ø' .Must have a blank in col 9 CLI BNF INVLSNTX ... of deblanked command buffer CLI 8(R4).X'4Ø' .Must not have a blank in col 8 ΒE ... of deblanked command buffer INVLSNTX .Must not have a blank in col 7 CLI 7(R4).X'4Ø' BNF SYNTAXOK ... of deblanked command buffer INVLSNTX L .Where command came from R8, FROMCNID .Where command came from LA R9.FROMSYS 'BLOCKCMD(E): -Invalid command syntax, must be "T NOCMD=X WTO XX',CONSID=(8),SYSNAME=(9),CART=CART LA R15.8 ST R15.RETCODE .Set non-zero return code GETSUFXX .Get out R MEMBER.=C'NOCMD' SYNTAXOK MVC .First part of the .Pick up the suffix from cmd buffer MVC SUFFIX,7(R4) MVT LASTBYTE,X'4Ø' .Move a blank into the last byte XR R15.R15 .Clear return code GETSUFXX PR This routine locates the member in the PARMLIB concatenation READPMEM BAKR R14.Ø ΝI EMPTY.NO .Make sure 'PARMLIB empty' flag off R2,WORKBUFF LA .Where we want the member content USING PRM READ BUFFER, R2 .Addressability to buffer .Total size of the buffer L R1.BUFFSIZE ST R1, PRM READ BUFF SIZE LA R1.IEFPRMLA .Point to the macro in storage IEFPRMLA(IEFPRMLL), IEFMACRO MVC IEFPRMLB REQUEST=ALLOCATE.CALLERNAME=CALLER.READ=YES. Х READBUF=(R2), MEMNAME=MEMNAME, MF=(E,(1)), Х ALLOCDDNAME=PARMDDØ1,RETCODE=RETCODE,RSNCODE=RSNCODE LTR R15.R15 .Was the member read in? ΒZ FREEPARM .Yes, go free it READERR L R8, FROMCNID .Where command came from R9.FROMSYS .Where command came from LA MVC ALOCWTOA(ALOCWTOL), ALOCWTOM BAS R14.CODEPRNT .Make RC and REASON printable MVC ALOCWTOA+72(4), PRTRC MVC ALOCWTOA+87(4), PRTRSN LA R1.ALOCWTOA

WTO MF=(E,(1)),CONSID=(8),SYSNAME=(9),CART=CART LA R15.8 ST R15.RETCODE .Set non-zero return code R READPMEX .Get out FREEPARM IEFPRMLB REQUEST=FREE.CALLERNAME=CALLER. χ DDNAME=PARMDDØ1,RETCODE=RETCODE,RSNCODE=RSNCODE LTR .Was the de-allocation successful? R15.R15 BN7 FRFFFRR .No. WORKBUFF+4(4),=F'Ø' .Did we get an empty PARMLIB member? CHKNUM CLC BNF READPMEX .No NOCARDS R8.FROMCNID .Where command came from L .Where command came from R9.FROMSYS LA WTO 'BLOCKCMD(W): -No commands will be blocked as PARMLIB meX mber is empty',CONSID=(8),SYSNAME=(9),CART=CART 0 I EMPTY,YES .Se 'empty' flag R15,R15 .Proceeding with empty table XR ST R15.RETCODE .Set non-zero return code .Get out В READPMEX FREEERR L R8, FROMCNID .Where command came from R9.FROMSYS .Where command came from LA MVC FREEWTOA(FREEWTOL), FREEWTOM .Make RC and REASON printable BAS R14.CODEPRNT MVC FREEWTOA+58(4), PRTRC MVC. FREEWTOA+73(4), PRTRSN LA R1.FREEWTOA WTO MF=(E,(1)),CONSID=(8),SYSNAME=(9),CART=CART ΙA R15.8 ST R15.RETCODE .Set non-zero return code READPMEX PR This routine returns the address of our subsystem. GETSSI@ BAKR R14.0 BAS R14.SEEKSSI .Get our subsystem's address LTR R15,R15 .Did we get the subsystem name? ΒZ GETSSI@X .Yes. get out SUFFIX,=C'NO' CLC .Are we turning blocking off? BNF CHKEMPT1 .No, go see if member is empty .Yes, we don't need the SSI address LA R15.8 В GETSSI@X .Get out CHKEMPT1 TM .Is PARMLIB member empty anyway? EMPTY.YES .No. there are entries in it BNO ADSUBSYS R15.8 .Stop further processing? LA ST R15.RETCODE .Plug the return code GETSSI@X .Get out В ADSUBSYS EQU .First time. add subsystem MVC IEFSSIA(IEFSSIL), IEFSSIM R1.IEFSSIA LA IEFSSI REQUEST=ADD,SUBNAME=SUBSYSN, Х RETCODE=RETCODE,RSNCODE=RSNCODE,MF=(E,(1)) LTR R15.R15 .Did we get the subsystem name? BNZ SSIERROR .No. error. Display and get out

BAS R14.SEEKSSI .Get our subsystem's address LTR R15.R15 .Did we get it? B7 GETSSI@X .Yes. we have the subsystem .Should never occur ABEND ØØØ1.DUMP SSIERROR BAS .No. go make RC & REASON printable R14.CODEPRNT .Where command came from 1 R8, FROMCNID LA R9.FROMSYS .Where command came from MVC SUBSWTOA(SUBSWTOL), SUBSWTOM MVC SUBSWTOA+42(4), PRTRC MVC SUBSWTOA+57(4).PRTRSN LA R1.SUBSWTOA WTO MF=(E,(1)),CONSID=(8),SYSNAME=(9),CART=CART R15.8 LA GETSSI@X PR This routine gets the address of our subsystem entry SEEKSSI BAKR R14.0 R4.16 .Start of CVT USING CVT,R4 .Establish addressability L R4.CVTJESCT .Get JESCT address USING JESCT.R4 .Establish addressability R4.JESSSCT .Get SSCT chain address 1 USING SSCT.R4 .Set addressability SSCTLOOP EQU .Look for our SSCVT * .End of chain? LTR R4.R4 ΒZ NOSSCT .Yes - not found .Our SSCVT? CLC SSCTSNAM, SUBSYSN ΒE FOUNDSSI .Got it L R4,SSCTSCTA .Point to next SSCVT entry .Redo the loop SSCTLOOP R NOSSCT R15.8 .Our subsystem not found LA SEEKSSIX .Get out В FOUNDSSI ST R4.OURSSI@ .Keep the address XR R15.R15 .Clear the return code SEEKSSIX PR This routine gets the address of the CSA table GETCSA@ BAKR R14.0 .Get address of CSA table L R4,OURSSI@ .Get address of subsystem entry USING SSCT,R4 .Addressability to our SSI CLC SSCTSUSE.=F'Ø' .Has the CSA table been obtained? .No, go do getmain ΒE OBTAIN XR R15,R15 .Clear return code MVC CSATAB@.SSCTSUSE .Plug the address GETCSA@X .Get out В OBTAIN FOU CLC SUFFIX.=C'NO' .Are we turning blocking off? BNE CHKEMPT2 .No LA R15.8 .We don't need the CSA address В GETCSA@X .Get out CHKEMPT2 TM EMPTY.YES .Is the PARMLIB member anyway empty?

BNO PREPCSA .No, we will have to obtain CSA LA R15.8 .Yes. do not GETMAIN CSA ST R15.RETCODE .Plug the return code В GETCSA@X .Get out PREPCSA L R3.BUFFSIZE .Length of storage to obtain .Add 4 bytes (first word= length) R3,4(R3) LA SL R3,=AL4(PRM RECORD TEXT-PRM READ HEADER) STORAGE OBTAIN, SP=228, LENGTH=(3), LOC=ANY, KEY=Ø LTR R15.R15 .Did we get the storage? BNZ GETCSA@X .No. get out .Plug the address into subsystem ST R1.SSCTSUSE ST R1.CSATAB@ .Plug the address 4(13.R1).=C'Uninitialized' MVC. ST R3.Ø(R1) .Put the length in the front GETCSA@X PR .Return to our caller This routine removes all blanks from the PARMLIB cards and * * copies them into the CSA table. (The size of the CSA table * is the same as that of the work bufer the data was read into. the fact that the IEFPRMLB macro succeeded reading the data * into the buffer means that the data will also fit into the * CSA table. We cannot have a changing CSA table size as storage * allocated in CSA can only be freed by the task that allocated * it in the first place.) COPYCRDS BAKR R14,Ø CLC SUFFIX.=C'NO' .Are we turning blocking off? BNF CHKEMPT3 .No XR .Yes. set # entries in tab to \emptyset R1.R1 L R2,CSATAB@ R1,Ø(R2) ST .Make card counter Ø R15.8 .Yes. set RC to "don't proceed" LA В COPYCRDX .Get out CHKEMPT3 TM EMPTY,YES .Is the PARMLIB member empty? BNO SETCNTR2 .No. go move the cards into the tbl XR R1,R1 R2.CSATAB@ 1 .Make card counter Ø SETCNTR1 ST R1.0(R2)COPYCRDX .Get out R SETCNTR2 LA R2.WORKBUFF .Where the PARMLIB cards are R4,15,PRM_RECORDS_READ_COUNT ICM 1 .Address of CSA table R5.CSATAB@ ST R4.Ø(R5) .Number of cards in first word .Where the cards will be moved LA R5,4(R5) LA R2,PRM_RECORD_TEXT .Start of the first record CARDLOOP LR R3.R5 MVC Ø(80,R3),=80X'40' .Move spaces into the card LA R6,8Ø .Number of bytes in a card MOVELOOP EQU * CLI .Blank in the card? Ø(R2),X'4Ø' ΒE .Yes, don't move this byte BUMPUP .Move the byte MVC Ø(1.R3).Ø(R2)

	LA	R3,1(R3) .Bu	ump up "to" pointer
BUMPUP	LA		ump up "from" pointer
	BCT	R6,MOVELOOP .Mc	ove the entire card
	LA	R5,8Ø(R5) .Wh	here next card should start
	BCT	R4,CARDLOOP .Do	o for each card
COPYCRDX			

*		routine makes the RC and	•

CODEPRNT		-	
	L	R1,RETCODE	
		-	
	UNPK OI	DOUBLE(4),DOUBLE+5(3) DOUBLE+3,X'FØ'	
	MVC	PRTRC,DOUBLE	
	L	R2,RSNCODE	
		R2,3,DOUBLE	
	STCM	R2,3,DOUBLE+2	
	NC		urn off right half of bytes
	NC		urn off left half of bytes
	TR		ake left half printable
	TR		ake right half printable
	MVC	DOUBLE+4(1),DOUBLE+1 Sw	wap bytes 2 and 3
	MVC	DOUBLE+1(1),DOUBLE+2	
	MVC	DOUBLE+2(1),DOUBLE+4 RS	SN code now printable
	MVC	PRTRSN,DOUBLE	
	CLC	RSNCODE+2(2),=X'ØØØA' (•
	BNE		o, other error
	L		here command came from
	LA	-	here command came from
	WTO		y commands in PARMLIB member. reduX
			assemble with larger table', X
CODEPRNX	חח	CONSID=(8),SYSNAME=(9),	,CARI=CARI,RUUICDE=II
		*****	*****
*		ants follow	
*******			*******
BUFFSIZE	DS	ØF	
	DC	AL4(8000)	
	DS	ØF	
PARMDDØ1	DC	C'PARMDDØ1' .DI	D-NAME used by IEFPRMLB
CALLER	DC	C'COMMAND '.Ca	aller name used for IEFPRMLB
IEFMACRO	IEFPR	MLB MF=(L,IEFPRMLB)	
IEFPRMLL	EQU	*-IEFMACRO	
		I MF=(L,IEFSSIWA)	
IEFSSIL		*-IEFSSIM	
ALOCWTOM	WTO		uring allocation/ reading of parmlX
		<pre>ib member,RC=XXXX(DEC);</pre>	
ALOCWTOL			ength of the message
FREEWTOM	WFO		uring free of PARMLIB member, RC=xX
		<pre>xxx(DEC), RSN=xxxx',MF=</pre>	
FREEWTOL	EQU	*-FREEWTOM .Le	ength of the message

SUBSWTOM	WTO	'BLOCKCMD(E): -Erron	<pre>r during IEFSSI, RC=xxxx(dec), RSN=xxX</pre>
		xx',MF=L	
SUBSWTOL LFTHALVE		*-SUBSWTOM ØCL24Ø	.Length of the message
LFINALVE	DC		,15X'00',X'F2',15X'00',X'F3'
	DC		0',X'F5',15X'ØØ',X'F6',15X'ØØ',X'F7'
	DC		0',X'C4',15X'ØØ',X'C5',15X'ØØ',X'C6'
VALCHARS		ØCL256	,,,
	DC	80X'01',X'00',10X'01	L',X'ØØ',16X'Ø1',2X'ØØ'
	DC	13X'Ø1',2X'ØØ'	
	DC	68X'Ø1',9X'ØØ',7X'Ø1	L',9X'ØØ',8X'Ø1'
	DC	8X'ØØ',6X'Ø1',1ØX'ØØ	
RGTHALVE		X'FØF1F2F3F4F5F6F7F8	
SUBSYSN	DC	C'NCMD'	.Name of our subsystem entry
والمروان والمروان والمروان والمروان	LTORG		****
*		s follow	**********
	2020.	• • • • • • • •	*****
STORAREA			
SAVEAREA		18F	.General savearea
PARMSTRT		F	.Start address of passed parms
FROMCNID		CL4	.Console id command came from
FROMSYS	DS	CL8	.Name of system command came from
RETCODE	DS	F	.Return code
RSNCODE	DS	F	.Reason code
PRTRC	DS	F	.Return code (printable)
PRTRSN	DS	F	.Reason code (printable)
OURSSI@	DS	F	.Address of our subsystem
CSATAB@	DS	F	.Address of the CSA table
MEMNAME MEMBER	DS DS	ØCL8 CL5	This contains "NOCMD"
SUFFIX	DS	CL2	.This contains "NOCMD" .Suffix of PARMLIB member to obtain
LASTBYTE		C	.This will contain a blank
EMPTY	DS	C	.Flag to indicate PARMLIB empty
DOUBLE	DS	D	.Double word work area
CART	DS	CL8	
ALOCWTOA	DS	CL(ALOCWTOL)	.Work area for WTO message
FREEWTOA	DS	CL(FREEWTOL)	.Work area for WTO message
SUBSWTOA		CL(SUBSWTOL)	.Work area for WTO message
IEFPRMLA		CL(IEFPRMLL)	.Work area for IEFPRMLB macro
IEFSSIA		CL(IEFSSIL)	.Work area for IEFSSI macro
WORKBUFF		*	.Parmlib buffer area
STORSIZE	-	*-STORAREA	.Length of area to allocate
RØ R1	EQU EQU	Ø 1	
R2	EQU	2	
R3	EQU	3	
R4	EQU	4	
R5	EQU	5	
R6	EQU	6	
R7	EQU	7	
R8	EQU	8	
R9	EQU	9	

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	END						
	IEFJS	QRY					
	IEFJS						
	IEFJS						
	IEFJE						
	CVT	DSECT=YES					
	IEFZP	MAP PRM_READ_BUFFER=	-YES				
	IEFZB	4D2					
	IEFZB	4DØ					
	IEZVX	1Ø1	.DSECT	for	command	exit	fields
YES	EQU	X'8Ø'					
NO	EQU	X'ØØ'					
R15	EQU	15					
R14	EQU	14					
R13	EQU	13					
R12	EQU	12					
R11	EQU	11					
R1Ø	EQU	1Ø					

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OS/390 Unix Assembler callable services

With the introduction of OpenEdition, or OS/390 Unix System Services as it is now called, IBM has provided a rich and varied set of program callable services that are available to both Assembler language and TSO/E REXX programmers. The main objective of this article is to introduce the reader to the OS/390 Unix callable services that are available to Assembler language programs and to also make reference to REXX callable services where applicable.

The first section of this article provides an overview of OS/390 Unix process management, which I have found to be an important learning stage before attempting to write Unix-style Assembler programs or REXX EXECs.

PROCESS MANAGEMENT

In OS/390 Unix, a child process is created when the parent process issues a fork service call. The creating process is called a parent process and the newly created process is called a child process. A parent process can have many child processes, the number of which is controlled by the MAXPROCUSER statement in the BPXPRMxx PARMLIB member. Each process is given a unique identifier called the Process ID (PID). Each process also knows the ID of its parent through the Parent Process ID (PPID). All processes are related to each other through the PIDs and the PPIDs. The originator of all processes is called the INIT process (PID = 1). In addition to having a PID, each process belongs to a process group. A process group is a collection of one or more processes. Each process group has a unique group ID. Process identifiers are classified as follows:

- PID a process ID. A unique identifier assigned to a process while it runs.
- PGID each process in a process group shares a process group ID (PGID), which is the same as the PID of the first process in the process group.
- PPID a process that creates a new process is called a parent process.

Process management can be broken down into the following areas:

- Processes
- Dubbing
- Threads
- Interprocess communication
- Signals.

A PROCESS

A process exists in an MVS address space and is identified by a TCB and related control blocks. In addition to the TCB, the OS/390 Unix kernel address space maintains a number of control blocks that represent a process. The following BPXPRMxx PARMLIB member statements control OS/390 Unix processes:

- MAXPROCSYS(nnnn) specifies the maximum number of OS/390 Unix processes that the system allows.
- MAXPROCUSER(nnnnn) specifies the maximum number of processes that a single OS/390 Unix user-id can have currently active, regardless of how the processes were created.

To control processes, the following Assembler callable services are available.

BPX1FRK (fork service)

BPX1FRK creates a new process. The fork service replicates the current process into a child process, which then runs in a new address space. The new address space contains a single task (thread) and a single RB (Request Block) structure. The fork service is supported from programs in PSW key 8 only. An additional requirement is that the storage protection key value in the TCBPKF field of the TCB must be 8. All key 8 virtual storage is copied from the parent to the child address space. The child process has a unique PID that does not match any active process group ID. The child has its own copy of the parent's open directory stream. The fork service can be requested from either an MVS or kernel address space. Control is given to the child process at the instruction following the fork service call and not at the program's main entry point. In most implementations, the parent process will continue and the child process will pass control using the EXEC service to a child-specific program. The BPXPRMxx PARMLIB member statement FORKCOPY(COW|COPY) specifies how user storage is to be copied from the parent process to the child process during a fork system call.

BPX1SPN (spawn service)

BPX1SPN spawns a process. The spawn service starts a new process, but the child process is started with another program in the Hierarchical File System (HFS). After the spawn service returns to the parent process, the two processes continue as independent processes. The main benefit of the spawn callable service is that it can create a new process in a separate address space or in the same address space, depending on the setting of the environment variable '_BPX_SHARES=YES|NO'. If an application is multi-threaded, you must use spawn instead of the fork callable service. There are some exceptions where, despite setting '_BPX_SHAREAS=YES', a non-local spawn (child process starts in another address space) is done. A non-local spawn is done in any of the following cases:

- The program that is spawned has the sticky bit on
- The program that is spawned is a SETUID or SETGID program
- The address space has exhausted its private storage.

A new setting of '_BPX_SHAREAS=MUST' was added in OS/390 Version 2 Release 6. It allows the application to force a local spawn or no spawn at all. The new process, which is called the child process, inherits the following attributes from the parent process (process that calls spawn):

- Session membership.
- Real user-id.
- Real group-id.
- Supplementary group-ids.
- Priority.
- Region size.
- Time limit.
- Accounting data.
- Working directory.
- Root directory.
- File creation mask.
- Signal mask.
- Security information, unless the '_BPX_USERID' environment variable specifies otherwise.
- TASKLIB STEPLIB or JOBLIB DD dataset allocations, unless the STEPLIB environment variable specifies otherwise. This causes the child's address space to have the same MVS program search order as the calling process.

- Accounting information.
- JOBNAME of the parent is propagated to the child and appended with a numeric value in the range 1-9, if the JOBNAME is 7 characters or less. If the JOBNAME is 8 characters, the JOBNAME is propagated as is.

The executable program to be run receives control with the following attributes:

- Problem program state
- PSW key 8
- AMODE=31
- Primary ASC mode.

BPX1ATX (attach_exec) service

BPX1ATX attach an OS/390 Unix program. The 'attach_exec' service attaches a task to run an OS/390 Unix executable program in a newly created child process of the caller. The new process is created in the same address space as the caller, and is a subtask of the caller's task. The child process has a unique PID that does not match any active process group ID. The child process has a parent process ID of the process that called attach_exec. The child process is terminated when its parent terminates. The executable file receives control with the following attributes:

- Problem program state
- TCB key of the caller
- AMODE=31
- Primary ASC mode.

The equivalent function is provided by the BPX1SPN SPAWN Service with '_BPX_SHAREAS=YES'. The OMVS command uses the attach_exec system call to run the shell in the TSO/E address space.

BPX1ATM (attach_execmvs) service

BPX1ATM ATTACHes an MVS program. The attach_exec MVS service creates a new child process in the same address space and passes control to a new program in the normal MVS search order (job pack queue, STEPLIB, LPALIB, LINKLIB). The new process that is created is run as a subtask in the address space.

BPX1EXC (EXEC) service

BPX1EXC run a program. The EXEC service does not start a new process, but replaces the program in the current process with another program as indicated on the EXEC call. A successful EXEC call will never return control to the calling program, but control is passed to the main entry point of the new program that is specified on the EXEC call. The executable program receives control with the following attributes:

- Problem program state
- PSW key 8
- AMODE=31
- Primary ASC mode.

The new process inherits the following attributes from the calling process:

- PID
- PPID
- The time left until an alarm signal is generated
- File mode creation mask
- Process signal mask
- Pending signals
- Time accounting information.

BPX1EXM (execmvs) service

BPX1EXM run an MVS program. The execmvs service runs an MVS executable program that is in the LPA or LNKLST concatenation. If

it is invoked from an address space that contains multiple processes, the program can come from a STEPLIB. The call can invoke both unauthorized and authorized programs:

- Unauthorized programs receive control in problem program state, with PSW key 8.
- Authorized programs receive control in problem program state, with PSW key 8 and APF authorization.

Additional Assembler process callable services

The following additional Assembler callable services are available:

•	chpriority (BPX1CHP)	—	change the scheduling priority of a process
•	getpid (BPX1GPI)	_	get the process ID
•	getppid (BPX1GPP)	_	get the parent process ID
•	getpgid (BPX1GEP)	_	get the process group ID
•	getpgrp (BPX1GPG)	_	get the process group ID
•	getpriority (BPX1GPY)	—	get the scheduling priority of a process
•	nice (BPX1NIC)	—	change the nice value of a process
•	setpgid (BPX1SPG)	—	set a process group ID for job control
•	setpriority (BPX1SPY)	—	set the scheduling priority of a process
•	_pid_affinity (BPX1PAF)) —	add or delete an entry in a process's affinity list.

REXX PROCESS MANAGEMENT

There is no FORK or EXEC SYSCALL commands available in REXX. Instead IBM recommend using SPAWN. The following REXX process management calls are available:

Spawn

Spawn invokes the spawn callable service to create a new process, called the child process. The program to be run is executed from the HFS. The child process has a unique PID that does not match any active process group ID. The child parent ID is set to the PID of the process that called spawn. To control whether the spawned child process runs in a separate address space or in the same address space, you can specify the '_BPX_SHAREAS' environment variable. If '_BPX_SHAREAS=NO' is specified, the child process to be created will run in a separate address space from the parent process. If '_BPX_SHAREAS=YES' is specified, the child process to be created is to run in the same address space as the parent. The following attributes are inherited by the child process from the parent process:

- Session membership
- Real user-id
- Real group ID
- Supplementary group IDs
- Priority
- Working directory
- Root directory
- File creation mask
- The process group ID of the parent is inherited by the child
- Signals set ignored in the parent are set to be ignored in the child
- The signal mask is inherited from the parent.

Example

The following example will spawn a new process. The new process is spawned to run/bin/ls. File descriptors greater than or equal to three are not available to the new process. fd.0 to fd.2 are used in remapping the file descriptors from the parent. The current environment is passed to the new process.

```
address syscall
/* Initialize the file descriptor map. The file descriptor map is set so */
/* that the file descriptor for the new file being created is remapped
                                                                          */
/* to file descriptor 1 for the new process. File descriptors 0 and 2
                                                                          */
will not be opened in the new process (-1).
                                                                          */
fd.Ø=-1
fd.2=-1
/* Create a new HFS file. The file descriptor is returned in RETVAL
                                                                          */
'creat /tmp/remdirl 755'
fd.1=retval
/* Initialize the parameter stem. The first parameter is set to the
                                                                          */
pathname for the file being spawned. Additional parameters are
                                                                          */
/* are set in the format the program expects. The parameters specif
                                                                          */
/*
     Display extended attributes for regular files
                                                                          */
/*
      Display permissions, links, owner, group, size, time, name
                                                                         */
/*
     Enables the audit bits to be displayed
                                                                          */
/*
     The current environment is propagated ( environment)
                                                                         */
parm.1='/bin/ls'
parm.2='-1WE'
parm.3='/u/rems'
parm.0=3
/* spawn a new process */
'spawn /bin/ls 3 fd. parm. __environment' /* PID(process ID) is returned
in RETAIL */
pid=retval
```

spawnp

Spawnp invokes the spawn callable service and creates a new process, called a child process. Spawnp functions identically to the spawn function except that it uses the PATH environment variable to resolve relative filenames.

forkexecm

Forkexecm invokes the fork and execmvs callable services to fork and EXEC a program to be executed from the MVS LINKLIB, LPALIB, or STEPLIB library. The call can invoke both authorized and unauthorized MVS programs. Authorized programs receive control in problem program state with PSW key 8 and APF authorization. Unauthorized programs receive control in problem state, with PSW key 8.

Example

The following example invokes the MVS program unixprog and passes a parameter to the program. On input, the MVS program receives a single entry parameter list pointed to by Register 1. The high order bit of the sole parameter entry is set to 1.

"forkexecm unixprog 'This is the parm info'

Dubbing

The first attempt to use a OS/390 Unix services dubs the MVS address space as an OS/390 Unix process. From a Unix perspective, this means OS/390 Unix assigns a PID (Process ID) to the process. Address spaces created by fork are automatically dubbed when they are created. Dubbing also adds the UID/GID assignment to an address space as follows:

Real UID

At process creation, the real UID identifies the user who has created the address space.

Effective UID

Each process has an effective UID. The effective UID is used to determine owner access privileges of a process. This is normally the same as the real UID but can be changed when a program is executed that has a special flag. A program with this special flag set is said to be a set-user-id program. This changes the effective UID of the process to the UID of the owner of the program, to allow additional permissions to the user while the set-user-id program is executed.

Real GID

At process creation, the real GID identifies the current connect group of the user for which the process was created.

Effective GID

Each process has an effective GID. The effective GID is used to determine group access privileges of a process. This is normally the same as the real GID but can be changed when a program is executed which has a special flag. A program with this special flag is said to be

a set-group-id program. This changes the effective GID of the process to the GID of the owner of the program, to allow additional permissions to the user while the set-group-id program is executed. Undub is the inverse of dub. Normally a task (dubbed a thread) is undubbed when it ends. An address space (dubbed a process) is undubbed when the last thread ends.

THE REXX SYSCALL ENVIRONMENT

IBM has provided two additional host command environments in REXX as follows:

- SYSCALL
- SH

To run a REXX program with SYSCALL commands from TSO/E or MVS batch, the syscalls ('ON') function at the beginning of the REXX program ensures that the address space is dubbed an OS/390 Unix process. For a REXX program that utilizes SYSCALL commands from the shell, SH is the initial host environment. The SYSCALL environment is automatically initialized as well, so there is no need to begin the REXX program with a syscalls('ON') call. The SYSCALL environment sets up the REXX pre-defined variables and blocks all signals. The syscalls('ON') function sets the following return code values:

- 0 Successful completion.
- 4 The signal process mask was not set.
- 7 The process was dubbed, but the SYSCALL environment was not established.
- 8 The process could not be dubbed.

The following example shows how you can use the syscalls ('ON') function at the beginning of a REXX program to establish the SYSCALL environment and get the address space dubbed as an OS/390 Unix process:

```
If syscalls('ON') >3 Then Do
    say 'Unable to establish the SYSCALL environment'
    Return
End
```

Or as an alternative:

To end the REXX SYSCALL environment the syscalls ('OFF') function ends the current task and OS/390 Unix process. The following rules apply:

- If the REXX program was run from TSO/E or batch, the task is undubbed, but the REXX program continues running.
- If the REXX program was run from the shell or a program, the REXX program is ended.

Threads

In POSIX, a thread is an entity that allows multiple simultaneous execution paths within a process. The OS/390 Unix design implementation for creating a thread is to attach a TCB within a process (address space). Threads allow multiple tasks to run in a single process within an address space. It allows for concurrent and asynchronous processing without the additional overhead associated with creating a new address space. Each thread of a process can run on an individual processor in a multi-processor environment. When using Assembler as opposed to the C language, it easier to do an ATTACH and let the tasks be dubbed as threads. Threads are created as follows:

- The pthread_create service
- The fork or EXEC service
- Most OS/390 Unix service requests from an undubbed MVS task.

The first routine that is given control in the new task when a thread is created with the pthread_create service is the pthread_create pthread-creating task initialization routine. Each thread that is created with pthread_create runs as a MVS subtask of the initial pthread_creating task (IPT). The IPT is the task that issued the first pthread_create call within the address space. Threads created by pthread_create are represented by a eight-character thread ID. There are three thread types.

Heavy-weight threads

A heavy-weight thread has been defined as a task that is attached when needed. A heavy-weight thread is created by issuing the pthread_create system call and specifying PTATWEIGHT= PTATHEAVY in the BPXYPTAT mapping macro. When a heavy-weight thread terminates, the task (TCB) that supports it is terminated and all EOT (end_of_task) resource managers are called to clean up after it.

Medium-weight threads

Medium-weight threads reuse MVS tasks. A medium-weight thread is created by issuing the pthread_create system call and specifying PTATWEIGHT=PTATMEDIUM in the BPXYPTAT mapping macro.

When a medium-weight thread is created, it is dispatched using an MVS task that is maintained in a pool. When a medium-thread terminates using pthread_exit, the MVS task is recycled in the pool without going through the MVS EOT (end_of_task) resource managers. OS/390 Unix reuses the task. An example of an OS/390 Unix application that uses medium-weight threading is the OS/390 Internet Connection Server.

Light-weight threads

Light-weight threads have not yet been implemented in LE/390 and OS/390 Unix. Thread manipulation is available using the following OS/390 Unix Assembler callable services:

- pthread_create (BPX1PTC) create a thread
- pthread_cancel (BPX1PTB) cancel a thread
- pthread_detach BPX1PTD) detach a thread
- pthread_exit_and_get (BPX1PTX) exit and GET a new thread
- pthread_join (BPX1PTJ) wait on a thread
- pthread_kill (BPX1PTK) send a signal to a thread
- pthread_quiesce (BPX1PTQ) quiesce threads in a process
- pthread_self (BPX1PTS) query the thread ID
- pthread_security_np (BPX1TLS) create/delete thread-level security environment for callers thread. An installation has the

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following two ways of allowing an application to use this service:

- Define the BPX.SERVER FACILITY class profile. For an application to access this service, it must be given read access to this profile.
- Assign the user-id a UID of 0 so that it operates as a superuser.
- pthread_setintr (BPX1PSI) examine and change the interrupt state
- pthread_setintrtype (BPX1PST) examine and change the interrupttype
- pthread_tag_np (BPX1PTT) set, query, or both set and query the callers thread tag data.
- pthread_testintr(BPX1PTI)-cause a cancellation point to occur.

Please refer to the *Unix System Services Programming: Assembler Callable Services Reference* for a detailed description of the thread callable services.

OS/390 Unix threads are controlled by the following BPXPRMxx PARMLIB member statements:

- MAXTHREADTASKS(nnnn)-specifies the maximum number of MVS tasks that a single process can have concurrently active for pthread_created threads.
- MAXTHREADS(nnnn) specifies the maximum number of pthread_created threads, including running, queued, and exited but undetached, that a single process can have concurrently active.

Signals

In OS/390 Unix applications, the basis for error handling is the generation, delivery, and handling of signals. Each process has a signal mask that defines the set of signals currently blocked from delivery and that is inherited by a child from its parent. Applications can be coded to generate and send signals, and to handle and respond to signals delivered to it. During the time between the generation of a signal and the delivery of a signal (when the actual signal is

performed), the signal is said to be pending. It is valid for the process to block the signal. If a signal that is blocked is generated for a process and the action for that signal is either the default action or to catch the signal, the signal remains pending for the process until the process either unblocks the signal or changes the action to ignore the signal.

The signal mask, which is passed as a parameter using the Assembler callable services, is structured as a 64-bit mask (8 bytes) of signals that are to be blocked during execution of the signal-catching function. The leftmost bit represents signal number 1, and the rightmost bit represents signal number 64. Bits that are set to 1 represent signals that are blocked. In the REXX callable services, the signal mask is defined as a string of 64 characters with values 0 or 1, representing the 64 bits in a signal mask.

Example

Re-sets the action for signals 5-64 to their default state:

X'ØFFFFFFFFFFFFF

Ignore signals 1-4

X'FØØØØØØØØØØØØØØ

The following OS/390 Unix SIGNAL callable services are available for the Assembler language programmer.

BPX1SIA (sigaction service)

BPX1SIA examines or changes a signal action. The sigaction service examines, changes or both examines and changes the action that is associated with a specific signal for all threads in the process.

BPX1SA2 (_sigactionset service)

BPX1SA2 examines or changes a set of signal actions. The _sigactionset service examines, changes or both examines and changes the actions that are associated with a set of signals.

BPX1SIP (sigpending service)

BPX1SIP examines pending signals. The sigpending service returns the union of signals that are pending on the thread and the set of signals that are pending on the process.

BPX1SIP (sigprocmask service)

BPX1SIP examines or changes a process's signal mask. The sigprocmask service examines, changes or both examines and changes the actions that are associated with a set of signals.

BPX1SSU (sigsuspend service)

BPX1SSU changes the signal mask and suspends the thread until a signal is delivered. The sigsuspend service replaces a thread's current signal mask with a new signal mask. It then suspends the caller's thread until delivery of a signal whose action is either to process a signal-catching service or to end the thread.

BPX1SWT (sigwait service)

BPX1SWT waits for a signal. The sigwait service waits for an asynchronous signal. If a signal that is specified in the signal set is sent to the invoker of sigwait, the value of that signal is returned to the invoker and the sigwait service ends. The following signals are supported:

- SIGHUP (1) Hang-up detected on controlling terminal
- SIGINT (2) Interactive attention
- SIGABRT (3) Abnormal termination
- SIGILL (4) Illegal or invalid hardware instruction
- SIGPOLL (5) Pollable event
- SIGURG (6) High bandwidth data is available at a socket
- SIGSTOP (7) Stop executing
- SIGFPE (8) Erroneous arithmetic operation (hardware and software)
- SIGKILL (9) An unconditional terminating signal
- SIGBUS (10)Bus error
- SIGSEGV (11) Invalid access to memory (hardware and software)
- SIGSYS (12)Bad system call

- SIGPIPE (13) Write on a pipe with no readers
- SIGALRM (14)Asynchronous time-out signal generated as a result of an alarm()
- SIGTERM (15)Termination
- SIGUSR1 (16)Reserved as an application-defined signal 1 (Software only)
- SIGUSR2 (17)Reserved as an application-defined signal 2 (Software only)
- SIGABND (18)Abend
- SIGCONT (19)Continue if stopped
- SIGCHLD (20)Child process terminated or stopped
- SIGTTIN (21) A background process is attempting a read
- SIGTTOU (22) A background process is attempting a write
- SIGIO (23)Completion of input or output
- SIGQUIT (24)Interactive termination
- SIGTSTP (25)Interactive stop
- SIGTRAP (26) Trap used by ptrace call
- SIGIOERR (27)I/O error. Serious software error such as a system read or write.
- SIGWINCH (28)Change size of window
- SIGXCPU (29)CPU time limit exceeded
- SIGXFSZ (30)File size limit exceeded
- SIGVTALARM (31) Virtual timer expired
- SIGPROF (32)Profiling timer expired
- SIGDCE (38)Exclusive use by DCE.

For a full description of signals that are available please refer to the mapping macro *BPXYSIGH* – *Signal Constants* in the *Unix System Services Programming: Assembler Callable Services Reference*.

The signals SIGSTOP and SIGKILL cannot be blocked or ignored, they are delivered to the program no matter what the signal mask specifies. Please note that the use of the system linkage stack with PC and BAKR instructions prevents signals from being delivered.

REXX SIGNAL SERVICES

The following REXX signal services are available:

- Sigaction this invokes the sigaction callable service to examine or change, or both, the action associated with a specific signal for all the threads in the process.
- Sigpending this invokes the sigpending callable service to return the union of the set of signals that are pending on the thread and the set of signals pending on the process. Pending signals at the process level are moved to the thread that called the sigpending callable service.
- Sigprocmask this invokes the sigprocmask callable service to examine or change the calling thread's signalmask.
- Sigsuspend this invokes the sigsuspend callable service to replace a thread's current signal mask with a new signal; then it suspends the caller's thread until delivery of a signal whose action is either to process a signal-catching service or to end the thread.
- Sleep this invokes the sleep callable service to suspend running of the calling thread (process) until the number of seconds has elapsed (*sleep* number), or until a signal is delivered to the calling thread to invoke a signal-catching function or to end the thread.
- Alarm this invokes the alarm callable service to generate a SIGALRM signal after the number of seconds specified has elapsed (*alarm* seconds)
- Kill this invokes the kill callable service to send a signal to a process or process group (*kill* pid signal).
- Pause this invokes the pause callable service to suspend execution of the calling thread until delivery of a signal that either executes a signal-catching function or ends the thread.

Establishing and deleting the signal interface routine in REXX

The syscalls ('SIGON') function establishes the Signal Interface Routine (SIR). For a REXX program run from the shell or a program, the SIR is established by default. After the SIR has been established, the sigaction syscall command can be issued to catch the signals that are to be processed and the sigprocmask syscall command can be used to unblock the signals. The syscalls ('SIGOFF') function deletes the signal interface routine.

The following example enables the sigalarm signal. Sigaction is used to set the action for the sigalrm to be caught. Sigprocmask is used to unblock sigalrm. The alarm is set by using the alarm service and the process waits for the completion of the child or the alarm:

```
/*REXX */
address syscall
/* Insert instructions to spawn a new process */
. . . . . . .
. . . . . . .
pid=retval
/* Child process ID from spawn
                                                                         */
call syscalls ('SIGON')
/* Establish the signal interface routine
                                                                         */
'sigaction' sigalrm sig_cat 0 'old_ handler old_flag'
/* catch a SIGALRM signal
                                                                         */
'sigprocmask' sig_unblock sigaddset(sigsetempty(),sigalrm) 'mask'
/* use sigaddset and sigsetempty to create a signal
                                                                         */
/* mask with the sigalrm bit
                                                                         */
'alarm 20'
/* set alarm to expire in 20 seconds
                                                                         */
'waitpid (pid) st. 0'
/* Wait for process termination or alarm
                                                                         */
alarm ind=retval
/* Status
                                                                         */
/* Check if alarm went off
                                                                         */
. . . . . . .
. . . . . . .
call syscalls ('SIGOFF') /* Turn off signals
                                                                         */
/* Determine process status using stem variable st. returned from
waitpid */
. . . . . . . . .
. . . . . . . . .
Exit
                                                                         */
/* All done
sigsetempty: return copies(0,64)
sigaddset: return overlay(1,arg(1),arg(2))
```

REXX statements for defining signal sets

REXX statements for defining signal sets are shown below. The C function is shown first followed by the equivalent REXX statement with its parameters and returns.

sigsetempty()	sigsetempty: return copies(0,64) Parameters: none Returns: signal set
sigsfillset()	sigfillset: return copies(1,64) Parameters: none Returns: signal set
sigsaddset()	sigaddset: return overlay(1,arg(1),arg(2)) Parameters: signal set, signal number Returns: signal set
sigsdelset()	sigdelset: return overlay(0,arg(1),arg (2)) Parameters: signal set, signal number Returns: signal set
sigismember()	sigismember: return substr(arg(1),arg(2),1) Parameters: signal set, signal number Returns: 0 (not member) or 1 (is member).

Refer to the REXX example in the section entitled *Establishing and Deleting the Signal Interface Routine in REXX*, for the sigsetempty and sigaddset REXX statements.

DETERMINING THE OS/390 CALLABLE SERVICE LEVEL

The Aassembler language programmer can determine the OS/390 Unix release-level by interrogating the CVT feature flags. At the time of writing the following values are defined:

•	CVTH6603	EQU	X'04'	HBB6603 (OS/390 Release 3) functions are present
•	CVTH6605	EQU	X'40'	HBB6605 (OS/390 Release 5) functions are present
•	CVTH6606	EQU	X'20'	HBB6606 (OS/390 Release 6) functions are present.

Assembler callable services syntax

To code a callable service, a CALL macro followed by the name of the callable service and parameter list is required. The required syntax is shown below:

CALL Servic	e_Name,(PARM_1, PARM_2,
	Return_value, Return_code, Reason_code)
CALL	CALL is the Assembler macro that passes control to the specified program and passes a parameter list.
Service_Name	Call service module name in the form BPX1xxx, where: 'xxx=' is a three-character symbol that is unique to the service. An example would be BPX1CHM=chmod service.
PARM parameter	rs PARM_1, PARM_2, etc, are placeholders for variables that may be part of a service syntax.
Return_value	Indicate the success failure of the callable service.
	If the callable service fails, a -1 is returned. For most successful calls to OS/390 Unix Services, the return value is set to 0. The BPX1GGI and BPX1GGN callable services return zeroes instead of -1 when the service fails. The fork (BPX1FRK) callable service returns a positive return value to indicate successful invocation.
Return_code	The return_code parameter is referred to as the errno in the POSIX and X/OPEN C interface. The Return Code is returned only if the service fails. All the Return codes and descriptions can be found in <i>OS/390 Unix System Services Messages and codes</i> .
Reason_code	The Reason_code parameter usually accompanies the Return_code value when the callable service fails. It further defines the return_code. All the Reason codes and their descriptions can be found in the <i>OS/390 Unix System Services Messages and</i> <i>Codes</i> .

Linkage conventions

The following linkage conventions are used when invoking the Callable Services:

- R1 Parameter list address. The last word in the list must have a 1 in the high order bit (Sign bit)
- R13 Savearea address
- R14 Return address
- R15 Entry point address of the service stub that is being called.

Additional notes

- 1 R2-R13 are restored on return from a callable service. R0, R1, R14, R15 are not restored.
- 2 The caller must be running with 31-bit addressing mode (AMODE=31).

Mapping macros

Many of the Assembler callable services provide mapping macros to map the parameter options. Many can be expanded with or without a DSECT statement. Please refer to the *Unix System Services Programming: Assembler Callable Services Reference* for the complete list of mapping macros that are available.

OS/390 UNIX ASSEMBLER CALLABLE SERVICES

The following list of OS/390 Unix Assembler Callable Services is a subset of what is currently available. Please refer to the *Unix System Services Programming: Assembler Callable Services Reference* for the complete list.

- access (BPX1ACC) determine if a file can be accessed.
- asyncio (BPX1AIO) asynchronous I/O for sockets.
- auth_check_resource_np(BPX1ACK) determine a user's access to protected MVS resource. The authorization required to invoke this service is one of the following:
 - Read access to the BPX.SERVER FACILITY class profile

- A UID of 0 when the BPX.SERVER Facility class profile is not defined.
- chattr (BPX1CHR) change the attributes of a file or directory.
- chaudit (BPX1CHA) change audit flags for a file by path.
- chdir (BPX1CHD) change the working directory.
- chmod (BPX1CHM) change the mode of a file or directory.
- chown (BPX1CHO) change the owner or group of a file directory.
- chroot (BPX1CRT) change the root directory.
- close (BPX1CLO) close a file.
- fchattr (BPX1FCR) change the attributes of a file or directory by descriptor.
- fchaudit (BPX1FCA) change audit flags for a file by descriptor. fchdir (BPX1FCD) – change the working directory.
- fchmod (BPX1FCM) change the mode of a file or directory by descriptor.
- fchown (BPX1FCO) change the owner and group of a file or directory by descriptor.
- fstat (BPX1FST) get status information about a file by descriptor.
- getcwd (BPX1GCW) get the pathname of the working directory.
- getegid (BPX1GEG) get the effective group ID.
- geteuid (BPX1GEU) get the effective user-id.
- getgid (BPX1GID) get the real group ID.
- getgroups (BPX1GGR) get a list of supplementary group IDs.
- getgroupsbyname (BPX1GUG) get a list of supplementary group IDs by user name.
- getpwnam (BPX1GPN) access the user database by user name.
- getpwuid (BPX1GPU) access the user database by user-id.

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- getuid (BPX1GUI) get the real user-id.
- getpwd (BPX1GWD) get the pathname of the working directory.
- lchown (BPX1LCO) change the owner or group of a file, directory, or symbolic link.
- link (BPX1LNK) create a link to a file.
- loadhfs (BPX1LOD) load a program into storage by HFS pathname.
- lstat (BPX1LST) get status information about a file or symbolic link by pathname.
- mkdir (BPX1MKD) make a directory.
- mount (BPX1MNT) make a file system available.
- mknod (BPX1MKN) make a directory, a FIFO, a character special, or a regular file.
- open (BPX1OPN) open a file.
- opendir (BPX1OPD) open a directory.
- openstat (BPX2OPN) opens a file, creates a file descriptor for it, and obtains its status.
- pipe (BPX1PIP) create an unnamed pipe.
- quiesce (BPX1QSE) quiesce a file system.
- read (BPX1RED) read from a file.
- read_extlink (BPX1RDX) read an external symbolic link.
- readdir (BPX1RDD) read an entry from a directory.
- readdir2 (BPX1RD2) read multiple entries from a directory.
- realpath (BPX1RPH) resolve a pathname.
- rename (BPX1REN) rename a file or directory.
- resource (BPX1RMG) obtain system-wide resource management data from the kernel address space.
- set _dub_default (BPX1SDD) Get the dub default service.

- setegid (BPX1SEG) set the effective GID.
- seteuid (BPX1SEU) set the effective user-id.
- setgid (BPX1SGI) set the GID.
- setgroups (BPX1SGR) set the supplementary group IDs list.
- setitimer (BPX1STR) set the value of the interval timer.
- setpgid (BPX1SPG) set a process GID for job control.
- setregid (BPX1SRG) set the real and/or effective GIDs.
- setuid (BPX1SUI) set user-ids.
- stat (BPX1STA) get status information about a file by pathname.
- symlink (BPX1SYM) create a symbolic link to a pathname.
- truncate (BPX1TRU) change the size of a file.
- ttyname (BPX2TYN) (X/Open version) get the name of a terminal.
- unmask (BPX1UMK) set the file mode creation mask.
- uname (BPX1UNA) obtain the name of the current operating system.
- unlink (BPX1UNL) remove a directory entry.
- unquiesce (BPX1UQS) unquiesce a file system.
- utime (BPX1UTI) set file access and modification times
- _wlm (BPX1WLM) WLM interface service.
- write (BPX1WRT) write to a file.

There are also a number of callable services that deal with:

- Socket processing
- Semaphores
- Memory mapping
- Message queue processing.

Callable Service Request Table (CSRTABLE)

When link-editing an Assembler module that contains OS/390 Unix Callable Services, a service stub must be included as part of the load module. The library containing the service stubs is specified in the SYSLIB concatenation in the linkage editor step. An alternative to this method is to use the system control offsets to callable services. The Callable Service Request Table (CSRTABLE), whose address is contained in the CVT provides the addresses of all the callable services. To locate the required callable service, an offset into this table is required. The full list of offsets is available in the Unix System Services Programming: Assembler Callable Services Reference.

When using the offsets, the registers must be set as follows:

- Register 1 To contain the address of the parameter list. Bit 0 of the last address in the list must be set on.
- Register 14 To contain the return address in the invoking module.
- Register 15 To contain the address of the callable service code.

The following example locates and executes the chown service:

L	R15,16	Address of the CVT
USING	CVT,R15	Obtain addressability to The CVT
L	R15,CVTCSRT	Pointer to the Callable Service Request Table
L	R15,24(,R15)	CSR slot
L	R15,64(,R15)	Chown slot entry = 64
BALR	R14,R15	Execute the callable service

OS/390 UNIX ASSEMBLER CALLABLE SERVICES EXAMPLES

<pre>************************************</pre>	THE ROOT DIRECTORY *
OPEN_READ_CLOSE_DIR EQU * MVC CALLAREA(CALLLEN),CALLL MVC DIR_LEN,=AL4(L'ROOT) MVC DIR_NAME(L'ROOT),ROOT OPEN_DIR EQU * CALL BPX10PD, (DIR_LEN,	
DIR_NAME, RETVAL, RETCODE, RSNCODE),	DIRECTORY NAME X RETURN VALUE:-1 OR FD X RETURN CODE X REASON CODE X

VL, Х MF=(E,CALLAREA) TEST RETVAL ICM R15.B'1111'.RETVAL ΒI OPEN DIR ABEND BRANCH IF NEGATIVE (-1 = FAILURE)STORE THE DIRECTORY STCM R15.B'1111'.DIR DESCP READ_DIR EQU * LA R15.DIR READ BUFFER DIR READ BUFFER ADDRESS STCM R15,B'1111',DIRREAD BUFFER@ STORE AWAY MVC DIR READ BUFFER LENGTH,=AL4(DIR READ BUFFER LEN) PRIMARY ALET, PRIMARY ALET PRIMARY ADDRESS SPACE ХC INVOKE_DIR_READ EQU * CALL BPX1RDD, Х READDIR (DIR DESCP. DIRECTORY FILE DESCRIPTOR χ DIRREAD_BUFFER@, BUFFER χ PRIMARY_ALET, BUFFER ALET χ DIR_READ_BUFFER_LENGTH, BUFFER SIZE χ RETVAL. RET VALUE: Ø. -1. ENTRIES READ χ RFTCODF. RETURN CODE χ RSNCODE). REASON CODE Х χ VL. MF=(E,CALLAREA) ICM R6,B'1111',RETVAL TEST RETVAL BL READ DIR ABEND -1= FAILURE ΒZ CLOSE DIR ALL DIRECTORY ENTRIES RETURNED R3,DIR_READ_BUFFER DIR READ BUFFER ADDRESS LA USING DIRE,R3 ADDRESSABILITY TO DIRE PROCESS DIR ENTRY EQU * * THE END OF A DIRECTORY IS INDICATED IN ONE OF TWO WAYS: * * 1 A RETURN VALUE OF Ø ENTRIES IS RETURNED. * * 2 SOME FILE SYSTEMS MAY RETURN A NULL NAME ENTRY AS THE LAST ENTRY * IN THE CALLERS BUFFER. A NULL ENTRY HAS AN ENTRY_LENGTH OF 4 AND A NAME_LENGTH OF Ø. CLC DIRENTINFO(L'DIRENTLEN+L'DIRENTNAML).=X'ØØØ4ØØØØ' ΒE CLOSE DIR ALL DIRECTORY ENTRIES RETURNED LR R4.R3 R4-> DIRE LA R4,L'DIRENTLEN+L'DIRENTNAML(Ø,Ø) START OF NAME SLR R5.R5 ZEROISE ICM R5,B'ØØ11',DIRENTNAML LOAD NAME LENGTH ALR R4.R5 $R4 \rightarrow END OF NAME + 1$ USING DIRENTPFSDATA, R4 ADDRESS PHYSICAL FILE SYSTEM * SPECIFIC DATA ICM R5.B'ØØ11'.DIRENTLEN ENTRY LENGTH ALR R3.R5 R3-> NEXT DIRE IN BUFFER BCT R6.PROCESS DIR ENTRY PROCESS ALL DIRECTORY ENTRIES READ NEXT DIRE INVOKE_DIR_READ В CLOSE_DIR EQU * CALL BPX1CLD. CLOSEDIR Х

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(DIR_DESCP, DIRECTORY FILE DESCRIPTOR χ RETVAL, RETURN VALUE: Ø OR -1 χ RETCODE. RETURN CODE Х RSNCODE). REASON CODE χ χ VL. MF=(E,CALLAREA) R15,B'1111',RETVAL ICM CLOSE DIR OKAY? CLOSE_DIR_ABEND BRANCH IF NEGATIVE (-1 = FAILURE)ΒL RETURN TO CALLER BR R2 OPEN_DIR_ABEND EQU * ABEND ØØ1.DUMP OPEN DIRECTORY ABEND CLOSE DIR ABEND EQU * ABEND ØØ2.DUMP CLOSE DIRECTORY ABEND READ DIR ABEND EQU * ABEND ØØ3.DUMP READ DIRECTORY ABEND TITLE 'STORAGE ITEMS' CALLL CALL ,(,,,,,,),MF=L CALLEN FOU *-CALLE IFNGTH WORKAREA DSECT SAVEAREA DS SAVEAREA CL72 PREVSA EQU SAVEAREA+4,4 @ OF PREVIOUS SAVEAREA C'/' ROOT DC ROOT DIRECTORY RETVAL DS AL4 RETURN VALUE RETCODE DS AL4 RETURN CODE RSNCODE DS AL4 REASON CODE DIR_DESCP DS AL4 DIRECTORY DESCRIPTOR PRIMARY ALET DS AL4 PRIMARY ALET FILE MODE FOR MKDIR SMODE DS XL(S MODE**#**LENGTH) DIRREAD BUFFER@ DS AL4 DIRECTORY BUFFER ADDRESS DIR OPEN BUFFER DS ØX DIRECTORY BUFFER FOR OPEN DIR_LEN DS AL4 DIRECTORY LENGTH DIR NAME DS CL1023 DIRECTORY NAME DIR OPEN LEN EQU *-DIR OPEN BUFFER OPEN DIR BUFFER LENGTH DIR_READ_BUFFER_AREA DS ØX DIRECTORY READ BUFFER DIR_READ_BUFFER_LENGTH DS XL4 DIRECTORY BUFFER LENGTH DIR READ BUFFER DS CL4Ø DIRECTORY BUFFER FOR READ DIR_READ_BUFFER_LEN EQU *-DIR_READ_BUFFER OPEN DIR BUFFER LENGTH CALLAREA DS CL(CALLLEN) PARM LIST AREA WORKALEN EQU *-WORKAREA WORK AREA LENGTH TITLE 'MAPPING OF A DIRECTORY ENTRY' BPXYDIRE DSECT=YES,LIST=YES * THIS EXAMPLE WILL CHANGE THE MODE OF A DIRECTORY * * SERVICES : CHMOD MOD_A_DIR_EQU * MVC CALLAREA(CALLLEN), CALLL MAKE RENT CHMOD_BUF(L'NEW_DIR), NEW_DIR MODIFY THE DIRECTORY MVC MVC CHMOD_LEN,=AL4(L'NEW_DIR) DIRECTORY NAME LENGTH LA R4.SMODE FILE MODE AREA

USING S_MODE,R4 ADDRESSABILITY XC S MODE.S MODE CLEAR MODE FLAGS ***** * OWNER= READ/WRITE/SEARCH* * GROUP= READ/WRITE/WRITE * * OTHER= READ/SEARCH * * MODE = 775 ***** MVI S_MODE2,S_IRUSR 775(RWX - RWX - R - X)MVI S_MODE3, S_IWUSR+S_IXUSR+S_IRWXG+S_IROTH+S_IXOTH CALL BPX1CHM. Х CHANGE FILE MODES (CHMOD LEN, PATHNAME LENGTH χ χ CHMOD BUF. PATHNAME S_MODE, BPXYMODE AND BPXYFTYP Х RETVAL. RETURN VALUE: Ø OR -1 χ RETCODE. RETURN CODE Х RSNCODE), REASON CODE Х VL. χ MF=(E,CALLAREA) R4,B'1111',RETVAL **USER INFO RETURNED?** ICM ΒL CHMOD ABEND BRANCH IF NEGATIVE (-1 = FAILURE)RETURN TO CALLER BR R2 CHMOD ABEND EQU * CHMOD ABEND ABEND ØØ4.DUMP TITLE 'LITERAL POOL' I TORG TITLE 'STORAGE ITEMS' CALLL CALL ,(,,,,,,),MF=L CALLLEN EQU *-CALLL LENGTH NEW DIR DC CL18'/u/remØØ1/testdir/' mkdir WORKAREA DSECT SAVEAREA DS CL72 SAVEAREA PREVSA EQU SAVEAREA+4,4 @ OF PREVIOUS SAVEAREA RETVAL DS AL4 RETURN VALUE RETCODE DS AL4 **RETURN CODE** RSNCODE DS AL4 REASON CODE XL(S_MODE#LENGTH) FILE MODE FOR CHMOD SMODE DS CHMOD BUFFER DS ØX CHMOD BUFFER DS AL4 CHMOD LENGTH CHMOD_LEN CHMOD BUF DS CL1ØØ DIRECTORY NAME MAX 100 CHARS CALLAREA DS CL(CALLLEN) PARM LIST AREA WORKALEN EQU *-WORKAREA WORK AREA LENGTH TITLE 'MAPPING OF A DIRECTORY ENTRY' BPXYDIRE DSECT=YES,LIST=YES TITLE 'FILE TYPE DEFINITIONS' BPXYFTYP DSECT=MEANINGLESS.LIST=YES TITLE 'MODE CONSTANTS FOR SYSCALL' BPXYMODE DSECT=YES.LIST=YES TITLE 'SYSCALL CONSTANTS' BPXYCONS DSECT=MEANINGLESS,LIST=YES

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Selecting messages from the log

INTRODUCTION

The following program was designed to extract messages from an MVS log or from a log-archiving file. You can select a specific period (a start time and an end time), the message IDs desired (up to nine can be specified), and the task that originated them. All these possibilities are optional.

The main difference between this program and a regular sort is that you can retrieve both the lines that contain the selected messages, and also the continuation lines. In a log there are some messages that span across several lines. I call these continuation lines, and you may wish to retrieve them also, as an option. Furthermore, there are message IDs that appear preceded by some symbols (for example, DFH messages appear with a plus sign before them). For a regular sort, you would need to know this and adjust your sort specifications. This program takes care of that, by automatically adjusting the comparison whenever the message IDs in the log begin with a '+', '-', or '\$'. The result of your request is placed in a file, with the same format as a regular log.

This application consists of a REXX EXEC with an associated panel that generates a job. This job copies the actual log to a temporary file, or uses as input an already existing file that you specify. Then it runs an Assembler program that makes the selections and produces the output file. At the beginning of this EXEC there are a few variables that you will need to set, namely the LOADLIB containing the Assembler program module, and the default output filename created. The input panel is shown in Figure 1.

In the example, you want to retrieve messages starting by HASP373, IEF403, and IEF404 that occurred between 23:00 and 1:00 hours, and from no job or STC in particular, since that field is blank. Note that message IDs need not be full message IDs, just the beginning letters. For example, you could specify IEF, and you would get all messages beginning with those letters. The same is true for the job / STC field. This field refers to the job identifier as it appears in log column 38. Also note that we are not asking for continuation lines, and that we

want to retrieve the messages from the 'real' log, not from some backup archive, since we leave the 'File' field empty.

You can also bypass the entry panel and the EXEC and work directly with a job. To do so, have a look at the job generated by the EXEC. All the DDnames and parameters needed by the Assembler program are explained in the comment on top of it. This way, you can use this program for automated tasks without the need of manual intervention.

Figure 1: The input panel

LOGMESS SOURCE

```
/* LOGMESS is a routine to extract selected messages from the log
                                                                  */
/*
      or from a log-like file. Its components are:
                                                                   */
/*
                                                                   */
/* LOGMESS - EXEC to generate a job to run the Assembler program.*/
/* LOGMESSP - An ISPF panel to input data for this EXEC.
                                                                  */
/* LOGMESSB - The Assembler program that does the work.
                                                                   */
                                                    ----*/
/*==
loadlib = "MY.LOADLIB" /* where logmessb module is */
tempfile = userid()".TEMPFIL" /* temporary file name */
jobnome = userid()".JOBTEMP" /* temporary job file name */
                                      /* output file name */
outfile = userid()".LOGMESS.D"right(date("S"),6)".T"time("S")
```

```
do k = 1 to 9
   interpret "M"k"=___
end
jobid="
c = "N"
                                          /* default no continuation */
do forever
   ADDRESS ISPEXEC "ADDPOP ROW(1) COLUMN(1)"
   ADDRESS ISPEXEC "DISPLAY PANEL(LOGMESSP)"
   if rc<>Ø then exit
   ADDRESS ISPEXEC "REMPOP"
   if yy="" & mm="" & dd="" & hi="" & hf="" & mi="" & mf="" then leave
   if yy="" | mm="" | dd="" | hi="" | hf="" then do
      ERRO = "Please specify start and end date/time fully"
      iterate
   end
   if mi="" then mi = "\emptyset\emptyset"
   if mf="" then mf = "\emptyset\emptyset"
   if ¬(datatype(yy,"W")&datatype(mm,"W")&,
        datatype(dd,"W")&datatype(hi,"W")&,
        datatype(hf,"W")&datatype(mi,"W")&,
        datatype(mf,"W")) then do
      ERRO = "Invalid non numeric value in date or time"
      iterate
   end
   if mi > 59 | mf > 59 then do
      ERRO = "Invalid minutes specified"
      iterate
   end
   if hi > 24 | hf > 24 then do
      ERRO = "Invalid hour specified"
      iterate
   end
   if mm > 12 | dd > 31 then do
      ERRO = "Invalid month or day specified"
      iterate
   end
   leave
end
continuation = c
datinij = ""
datendj = ""
if yy<>"" then do
   yy = right(yy,2,"\emptyset")
   mm = right(mm, 2, "\emptyset")
   dd = right(dd, 2, "\emptyset")
   datini = mm"/"dd"/"yy
   datinij = data_jul(yy||mm||dd)
   if hf < hi then datendj = datinij+1
   else datendj = datinij
   hi = right(hi, 2, "\emptyset")":"right(mi, 2, "\emptyset")
   hf = right(hf,2,"Ø")":"right(mf,2,"Ø")
end
```

```
if hi = "" then hi = "*"
if hf = "" then hf = "*"
beg_time = strip(datinij hi)
end_time = strip(datendj hf)
x = \emptyset
do k = 1 to 9
  interpret "mesg=M"k
   mesg = space(translate(mesg," ","_"),Ø)
   if mesq <>"" then do
      zz= left(mesg,1)
      if zz="+"|zz="-"|zz='$' then mesg=substr(mesg,2)
     x = x+1
     msg.x = mesg
   end
end
jobid = space(translate(jobid," ","_"),Ø)
if jobid = "" then jobid = "*"
xx = msg(off)
delete jobnome
"free dd (jobe)"
"alloc da('"jobnome"') dd(jobe),
   new blksize(8000) lrecl(80) recfm(f,b),
   dsorg(ps) space(1 1) tracks delete "
if rc<>Ø then do
  say "Error "rc" allocating "jobnome
   exit
end
dropbuf
queue "//"userid()"1 JOB CLASS=X,MSGCLASS=X,"
                MSGLEVEL=(1,1),REGION=2000K"
queue "//
queue "//*"
queue "//STEPØ
                 EXEC PGM=IEFBR14"
if logfile = "" then do
  queue "//FICTEMP DD DISP=(NEW,CATLG,DELETE),UNIT=SYSDA,"
                  DSN="tempfile","
  queue "//
  queue "//
                  RECFM=FB,LRECL=133,BLKSIZE=133ØØ,"
 queue "//
                  DSORG=PS,SPACE=(TRK,(90,90))"
end
queue "//SYSPRINT DD SYSOUT=*"
queue "//*"
if logfile = "" then do
  gueue "//STEP1 EXEC PGM=SDSF.PARM='++60,228'"
  queue "//ISFOUT DD DUMMY"
  queue "//ISFIN DD *"
  queue "LOG"
  queue "PRINT ODSN '"tempfile"' * SHR"
  queue "PRINT 1 999999"
  queue "PRINT CLOSE"
  queue "/*"
  queue "//*"
end
queue "//STEP2 EXEC PGM=LOGMESSB"
if logfile = "" then ,
```

```
queue "//FICTEMP DD DISP=(OLD,DELETE),DSN="tempfile
else,
  gueue "//FICTEMP DD DISP=SHR,DSN="logfile
gueue "//STEPLIB DD DISP=SHR,DSN="loadlib
gueue "//SYSPRINT DD SYSOUT=*"
queue "//SAIDA DD DISP=(NEW,CATLG,DELETE),"
queue "// SPACE=(TRK,(30,30),RLSE),"
                  RECFM=FB.LRECL=133.UNIT=SYSDA."
queue "//
queue "//
                  DSN="outfile
queue "//PARMLINE DD *"
queue continuation
queue beg_time
queue end time
queue jobid
do j = 1 to x
  queue msg.j
end
queue "/*"
queue ""
"execio * diskw jobe (finis"
"submit '"jobnome"'"
say "Job submitted; See output in "outfile
"free dd(jobe)"
exit
/*===
                                                             ____* /
/*
          data_jul converts dates from YYMMDD to YYDDD
                                                                     */
____* /
data_jul: procedure
parse arg date_in
aa = left(date_in,2)
mm = substr(date_in,3,2)
dd = right(date in.2)
if aa//4 = \emptyset then ac = 1
else ac = \emptyset
select
   when mm = 1 then x = \emptyset
   when mm = 2 then x = 31
  when mm = 3 then x = 59 + ac
   when mm = 4 then x = 90 + ac
   when mm = 5 then x = 120 + ac
   when mm = 6 then x = 151 + ac
   when mm = 7 then x = 181 + ac
   when mm = 8 then x = 212 + ac
   when mm = 9 then x = 243 + ac
  when mm = 10 then x = 273 + ac
   when mm = 11 then x = 304 + ac
   when mm = 12 then x = 334 + ac
   otherwise nop
end
j = x + dd
return aa||right(j,3,'Ø')
```

LOGMESSP SOURCE

```
)ATTR
 TYPE(INPUT) COLOR(RED)
                              JUST(LEFT) CAPS(ON)
 $ TYPE(OUTPUT) COLOR(WHITE) SKIP(ON)
                                          INTENS(HIGH)
 ? TYPE(TEXT) COLOR(PINK)
                              SKIP(ON)
                                          INTENS(HIGH)
 % TYPE(TEXT)
                COLOR(YELLOW) SKIP(ON)
                                          INTENS(HIGH)
 # TYPE(TEXT) COLOR(WHITE) SKIP(ON)
                                          INTENS(HIGH)
 + TYPE(TEXT) COLOR(GREEN) SKIP(ON)
                                          INTENS(LOW)
)BODY WINDOW(71,19)
+
? File (empty for LOG):_LOGFILE
+
? Date and hour beginning / end (optional):
       Date beginning (yy/mm/dd)..:_YY%/_MM%/_DD+
%
%
       Hour beginning (ØØ to 23)..:_HI%:_MI+
%
       Hour end
                      (ØØ to 23)..: HF%: MF+
+
? Job/Stc/Tsu identifier (optional).:_JOBID
                                              +
+
? Messages to search..:_M1
                                    +_M2
                                                 +_M3
                                                                +
                       _M4
                                    +_M5
%
                                                 +_M6
                                                                +
%
                       _M7
                                    +_M8
                                                 +_M9
                                                                +
+
? Continuation lines (Y, N)..:_C
+
#
       $ERRO
                                                PF3/15 - Cancel
     ENTER - Execute
)INIT
\&END = PFK(END)
&ZWINTTL = 'Log Messages'
)END
```

+

LOGMESSB SOURCE

*:		*
*		*
*	LOGMESSB - This program extracts messages from a log-type file,	*
*	including continuation lines. The input file must be	*
*	identical to the produced by an SDSF PRINT command (with CCs	*
*	at column one). The following DDnames are used in this program:	*
*		*
*	FICTEMP - The input file containing a log print.	*
*	SAIDA - The output file, similar to the input, containing	*
*	only the desired messages.	*
*	SYSPRINT – Standard job output.	*
*	PARMLINE – Parameter file that controls message selection.	*
*	It's format is as follows:	*
*	Line1.:C Continuation lines (Y or N).	*
*	Line2.:YYDDD HH:MM Start date/hour or * for all.	*
*	Line3.:YYDDD HH:MM End date/hour or * for all.	*

* Line4.: Job/Stc/Tsu as it appears at log column 38 or * for all. * * Line5 onwards: message identifier as it appears at log column 57. * * Up to 9 identifiers can be specified, one per line, or none. * * Do not include + - \$ symbols that appear at the beginning * * * of some messages. + * *----&PROGRAM SETC 'LOGMESSB' This program's name Jobid column in log file &JOBPOS SETC '38' &NUMPOS SETC '43' Number column in log file &MSGPOS SETC **'**57' Message column in log file &DATPOS SETC '2Ø' Date column in log file &TI SETC '12' Length of parm table entries & PROGRAM CSECT (minus 4 bytes for length). & PROGRAM AMODE 31 & PROGRAM RMODE 24 SAVE (14,12) Start stuff LR R12,R15 USING & PROGRAM, R12 R13.SAVEA+4 ST R11,SAVEA LA ST R11,8(R13) LR R13.R11 B OPENPRT DC CL16' & PROGRAM 1.1' DC CL8'&SYSDATE' Open files OPENPRT DS ØН OPEN (SYSPRINT, OUTPUT) LTR R15.R15 BNZ EXIT OPENSAI DS ØН OPEN (SAIDA, OUTPUT) LTR R15,R15 ΒZ OPENTEMP MVC XMSGTYPE,=C'OPEN ' MVC XMSGDSN,=CL44'SAIDA. Program terminated.' PUT SYSPRINT, XMSGLINE В EXIT OPENTEMP DS ØН OPEN (FICTEMP, INPUT) LTR R15,R15 ΒZ OPENPARM XMSGTYPE,=C'OPEN ' MVC MVC XMSGDSN,=CL44'TEMPFILE. Program terminated.' PUT SYSPRINT.XMSGLINE EXIT В OPENPARM DS ØН OPEN (PARMLINE, INPUT) LTR R15,R15

ΒZ GETCONTN MVC XMSGTYPE,=C'OPEN ' MVC XMSGDSN,=CL44'PARMLINE. Program terminated.' PUT SYSPRINT.XMSGLINE B FXIT -+ Read parameter file and process it. GETCONTN EOU * Get continuation lines option GET PARMLINE, CONTINUE GETBEGTI EOU Get beginning time * L R6,=F'-1' Assume no begtime specified. R7.=F'-1' Assume no endtime specified. L R2.=F'-1' L Assume no job/stc specified. R5 is table elements counter XR R5.R5 LA R4,MSGTAB Load message table address in R4 Get beginning date and hour GET PARMLINE, BEGTIME CLI BEGTIME,C'*' No begtime? ΒF GETENDTI No, jump ahead. R6.=F'1Ø' Otherwise, leng is 11 (Ex 1Ø). GETENDTI EQU PARMLINE, ENDTIME Same thing for endtime. GET CLI ENDTIME.C'*' No endtime? ΒF GETJOBID No, jump ahead. R7.=F'1Ø' Otherwise, leng is 11 (Ex 10). L GETJOBID EQU * Get job selection GET PARMLINE, JOBID All jobs? CLI JOBID.C'*' ΒF GETPARM Yes, jump SR No, find jobname length R9.R9 LA R3,JOBID Copy address to R3 Call find space subroutine BAL R1Ø, FINDSPC R9,=H'1' Length ready for executed CLC SH LR R2,R9 Keep it in R2 GETPARM EOU * GET PARMLINE,Ø(Ø,R4) Get parms (msg ids) SR Parm table entry R9,R9 LR R3.R4 Copy initial pointer to R3 Call find space subroutine BAL R1Ø.FINDSPC SH R9,=H'1' Length ready for executed CLC ST R9,&TL+Ø(Ø,R4) Store length after parm LA R4.&TL+4(Ø.R4)Next tab entry (TL + 4)Increment counter LA R5,1(Ø,R5) СН R5.=H'9' Limit of 9 attained? ΒE PARMEND Yes, ignore others and go ahead В GETPARM Otherwise, get another. PARMEND EOU * CLOSE PARMLINE Store number of table elements ST R5,MSGNUM XR R8.R8 Clear lines read counter MVI LMESSAGE,C'Ø' Clear message flag

* Read	a line	from the log file an	d process it. *
*======= READFILE	GET CLI BE LA CLI BNE	READFILE R8,1(Ø,R8) CONTINUE,C'N' READFIL1 &DATPOS+WFICTEMP(23)	Read log line Header line with CC on column 1 Yes, ignore it Increment line counter No continuation lines? Yes, jump ahead ,=23C' ' Date, time or job/stc? Yes, read next Else, jump to test start date.
* READFIL1	CLI BZ CLC BNE	* LMESSAGE,C'Ø' SEEKNUM &DATPOS+WFICTEMP(26) SEEKNUM SAIDA,WFICTEMP FINDNUMB	Message flag on? No, jump ahead ,=26C' ' Date, time or job/stc? Yes, jump Else, write line (type 1 line) Go look for continuation number.
			" in numpos (line type 2). * numtable. If found, write line. *
SEEKNUM	CLC BE CLC	BEGTEST &JOBPOS+WFICTEMP(5), BEGTEST &NUMPOS+WFICTEMP(3),	Continuation number process Are there elements in numtab? No, jump ahead. =5C' ' Jobpos first 5 bytes blank? No, jump ahead =3C' ' Numpos 3 bytes blank? Yes, jump ahead Load numbers table address.
SEEKNUM1	EQU CLC BE CLI BE CLI BE	* Ø(2,R9),=C'' SEEKNEXT Ø(R9),X'ØØ' BEGTEST Ø(R9),X'FF' BEGTEST	Entry erased (2 spaces), loop to next. Logical end of table(lowvalue), exit search. Physical end of table(highvalue), exit search.
SEEKNUM2		* RØ,R8 RØ,=F'1ØØ' RØ,4(Ø,R9) SEEKNUM3 R11,CLEARTAB SEEKNEXT	Compare current line (R8) minus 100 with line number in table. If greater, call routine to clear table entry.
SEEKNUM3	EQU CLC BNE PUT	* 1(3,R9),&NUMPOS+WFIC SEEKNEXT SAIDA,WFICTEMP	TEMP Compare number No match, loop to next entry. Match, write line (type 2)

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SEEKNEXT	CLC BNE BAL B EQU LA B	WFICTEMP(2),=C' E' READFILE R11,CLEARTAB READFILE * R9,8(Ø,R9) SEEKNUM1	End multiple line? No, read next Yes, clear table entry and read next. Next entry (8 bytes more)
* Test	for st	art / end hour and dat	ze. *
BEGTEST	CLC BE MVI LTR BM	* WFICTEMP(2),=C' S' READFILE LMESSAGE,C'Ø' R6,R6 ENDTEST R6,COMPARE2 READFILE R6,=F'-1' JOBTEST	Start date hour test Sequence (continuation) line? Yes, read next Reset message flag Start specified? If R6 negative, no, jump Else, compare Lower: start not reached Start reached: destroy test Else, accept line.
ENDTEST	EQU	* R7,R7 JOBTEST R7,COMPARE3 EXIT	End date hour test End specified? If R7 negative, no, jump Else, compare High: end reached, exit program.
* Test *	for jo	bnames / stcnames and	for message ids. *
JOBTEST		* R2,R2 MSGTEST R2,COMPARE4 READFILE	Job/stc/tsu id specified? No, jump ahead (R2 negative) Yes, execute compare No match, read another line.
MSGTEST		* R11,MSGNUM R11,R11 WRITLINØ R4,MSGTAB	Number of msg ids specified. If zero, write all lines, otherwise, load msg table address for compares.
MSGTEST1	EQU L CLI BE CLI BE CLI BE EX BNE B	* R5,&TL+Ø(Ø,R4) &MSGPOS+WFICTEMP,C'+' MSGTEST2 &MSGPOS+WFICTEMP,C'-' MSGTEST2	See if message match parms Load parm length (EX ready) Message starts by '+'? Yes, jump to correct execute Message starts by '-'? Yes, jump to correct execute Message starts by '\$'? Yes, jump to correct execute Compare message Not equal, loop for next Otherwise, junp to write line
MSGTEST2	EQU EX	* R5,COMPARE1	Compare one byte ahead for messages starting with + - \$.

MSGTEST3	BE EQU LA BCT B	WRITLINØ * R4,&TL+4(Ø,R4) R11,MSGTEST1 READFILE	Next table address (4+TL) Loop for number of entries.
WRITLINØ		* SAIDA,WFICTEMP CONTINUE,C'N' READFILE	Write line (type Ø line). No continuation lines? No, read next line.
* See i	f a li	ne ends with a three-	digit number. *
FINDNUMB	EQU	*	
FINDNUM	LA EQU	R1Ø,WFICTEMP+132 *	Point to end of line
TINDHOIT	CLI BNE BCT	Ø(R1Ø),X'4Ø' FINDNUM1 R1Ø,FINDNUM	Look for three digits preceded by a space. Loop to beginning of line.
FINDNUM1	CLI BL CLI BH S CLI BH S CLI BH S CLI BH S CLI BE EQU	<pre>FINDNUMF * Ø(R1Ø),C'Ø' FINDNUMF Ø(R1Ø),C'9' FINDNUMF R1Ø,=F'1' Ø(R1Ø),C'0' FINDNUMF Ø(R1Ø),C'9' FINDNUMF R1Ø,=F'1' Ø(R1Ø),C'9' FINDNUMF R1Ø,=F'1' Ø(R1Ø),C'9' SEEKFREE * LMESSAGE,C'1' READFILE</pre>	Empty line? 3 digits found, move ahead. Not 3 digits, set message flag and read next line.
			e to store the number. * ver 100 lines gone by). *
SEEKFREE	LA	* R9,NUMTABLE * Ø(2,R9),=C' MOVENUM Ø(R9),X'ØØ' MOVENUM	Seek for a free entry in the table R9 points beginning of table Look for a free entry in the table (either with spaces or low-values).

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	LR S C BL BAL	RØ,R8 RØ,=F'1ØØ' RØ,4(Ø,R9) SEEKFRE1 R11,CLEARTAB *	Compare current line (R8) minus 100 with line number in table. If greater, clear table entry.
SEEKFRE1	EQU LA CLI BE EQU MVC ST LH AH STH B	<pre>^ R9,8(Ø,R9) Ø(R9),X'FF' READFILE SEEKLOOP * Ø(4,R9),Ø(R1Ø) R8,4(Ø,R9) RØ,NUMTOTAL RØ,=H'1' RØ,NUMTOTAL READFILE</pre>	Increment tab pointer End of table? Yes, abandon the search. Else loop. Found empty entry, move number (with initial space). Store line number. Increment numtotal.
* Close	files	and exit.	*
EXIT *	CLOSE CLOSE L LM XR BR	* SAIDA FICTEMP SYSPRINT R13,SAVEA+4 R14,R12,12(R13) R15,R15 R14 , executed instruction	*
*			*
	EQU CLI BE LA LA B	* Ø(R3),X'4Ø' FINDSPCF R9,1(Ø,R9) R3,1(Ø,R3) FINDSPC *	Find space subroutine: counts the number of chars in a string up to the first space. Returns the number in R9. R3 points current position
FINDSPCF	BR	Â R1Ø	Return
CLEARTAB		*	Clear table entry with spaces
	MVC LH SH STH	Ø(8,R9),=8C' ' RØ,NUMTOTAL RØ,=H'1' RØ,NUMTOTAL	Decrement numtotal.
* COMPAREØ COMPARE1 COMPARE2 COMPARE3	CLC CLC		

COMPARE4 *	CLC	&JOBPOS+WFICTEMP(Ø),	JOBID Compare jobnames	
	LTORG			
SAVEA	DS	18F	Save register area	
CONTINUE		C	Continuation lines option	
DEOTINE	DS	ØF		
BEGTIME	DS	CL12	Beg date and hour (yyddd hh:mm)	
ENDTIME JOBID	DS DS	CL12 CL8	End date and hour. Job selection	
MSGNUM		F	Number of entries in parm table	
MSGTAB	DS	36F	9 entries max (12+4 bytes each)	
WFICTEMP		CL133	Log line read area	
LMESSAGE		C	Message flag (zero or one)	
NUMTOTAL		Н'Ø'	Total entries in numtable	
NUMTABLE	DC	6ØF'Ø'	Numbers table: 30 entries.	
ENDTABLE	DC	X'FF'	End of table mark.	
XMSGLINE	DS	ØCL8Ø	Error messages line	
	DC	C'===>>>>> Error in	•	
XMSGTYPE		C'''		
	DC	C' dataset '		
XMSGDSN *	DC	CL50' '		
PARMLINE	DCB	DSORG=PS,RECFM=FB,MA LRECL=8Ø, EODAD=PARMEND, DDNAME=PARMLINE	CRF=(GM),	X X X
*				
FICTEMP	DCB	DSORG=PS,RECFM=FB,MA LRECL=133, EODAD=EXIT, DDNAME=FICTEMP	CRF=(GM),	X X X
*				
SAIDA	DCB	DSORG=PS,RECFM=FB,MA LRECL=133, DDNAME=SAIDA	CRF=(PM),	X X
*				
SYSPRINT	DCB	DSORG=PS,RECFM=FB,MA LRECL=8Ø, DDNAME=SYSPRINT	CRF=(PM),	X X
*		DUMARE SISENINI		
	YREGS END			

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BMC has launched Version 2.1 of InTune OS/390 application performance tuner, promising an ability to share information between systems via Parallel Sysplex, to analyse the target application, and simplify customization of batch reports. The product analyses the performance of programs running on OS/390, CICS, DB2, IMS, NATURAL, ADABAS, and Datacom, and pinpoints code efficiency problems.

Users can fine-tune applications in production, in test, or under development. Version 2.1 identifies application program delays and presents this information for analysis through an interactive interface for both traditional and Parallel Sysplex environments. The product lets users go inside the code and view precisely where a resource delay is occurring. In-built support for Parallel Sysplex allows organizations to share performance information between various systems within the sysplex, allowing users to choose all or any specific individual systems when invoking requests.

For further information contact:

BMC Software, Inc, 2101 Citywest Blvd, Houston, TX 77042, USA. Tel: (713) 918 8800 Fax: (713) 918 8000 or

BMC Software Ltd, Compass House, 207-215 London Road, Camberley, Surrey, GU15 3E, UK. Tel: (01276) 24622 Fax: (01276) 61201 http://www.bmc.com

* * *

PROIV has announced PROIV Mainframe 4.0, a Web-enabled version of its System/ 390-based application development software. PROIV Mainframe 4.0 has a full Windows-style GUI, and offers the ability to port applications between OS/390, Unix, and NT.

The utility allows developers to create Windows and Java front-ends for mainframe applications, opening up the whole environment to the Internet and intranet using standard Web browsers. The PROIV JavaSuite allows applications to be delivered to Web browser clients without the need to amend application code.

PROIV applications are databaseindependent and may be deployed on mainframes running DB2, DL/I or VSAM. PROIV Mainframe 4.0 supports the following clients; Windows 95/98 and NT, 3270, 5550 (double-byte character support), and Web browsers (Java); it is also Year 2000 compliant.

For further information contact:

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