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MVS

November 2003

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MVS Update

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Subscriptions and back-issues

A year's subscription to *MVS Update*, comprising twelve monthly issues, costs £340.00 in the UK; \$505.00 in the USA and Canada; £346.00 in Europe; £352.00 in Australasia and Japan; and £350.00 elsewhere. In all cases the price includes postage. Individual issues, starting with the January 1999 issue, are available separately to subscribers for £29.00 (\$43.50) each including postage.

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Trevor Eddolls

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Printed in England.

NOTIFYEXTENT catalog option

INTRODUCTION

This article describes a new CATALOG function that is available with the DFSMS APAR OW54162 (DF/SMS 1F0 – UW88212/DFSMS 1G0 – UW88213).

This new function implements the new MODIFY CATALOG,NOTIFYEXTENT(xx) parameter, which notifies you about catalogs using more than xx per cent of their maximum number of extents (123).

This enhancement is designed to help reduce system outages resulting from catalogs reaching the maximum number of extents without warning. When this occurs, any catalog operations that need to extend the catalog again will fail, and this can cause outages in online systems or major applications.

The installation may now set a threshold value as a percentage of the maximum extents, and, for any catalog that exceeds that percentage, an immediate action message is sent to the console to warn that the threshold has been exceeded.

The installation can then take preventative action as necessary to prevent an unscheduled outage. The percentage of the maximum is calculated as the allocated extents divided by 123, multiplied by 100 and rounded to the nearest per cent.

IMPLEMENTATION

The enhancement consists of four parts:

- A new form of the MODIFY CATALOG command to establish the threshold for messages:

MODIFY CATALOG, NOTIFYEXTENT(xxx)

where xxx is a percentage number from 0 to 99. A value of zero means that normal monitoring will be suppressed, and is the default.

This setting is retained across catalog restarts, but not IPLs. So you should add this command to your automation package to activate this functionality. It is not possible to add this command in COMMND00 because it is processed before the CATALOG address space is active.

- New message IEC361I is issued whenever either the index or data component of a catalog exceeds the current threshold:

```
IEC361I CATALOG catname HAS REACHED xxx% OF THE MAXIMUM EXTENTS
```

This message is issued as an immediate action (descriptor code 2) message unless the percentage value exceeds 90%, and then it is issued as a critical action message (descriptor code 11). If normal monitoring has been disabled (eg the current threshold is zero), the message will still be issued when any catalog exceeds 90%.

- Message IEC348I, resulting from the MODIFY CATALOG, ALLOCATED command, has been modified to display the current percentage allocation values for all catalogs that are not deleted or closed, and that have been referenced since the last IPL:

```
F CATALOG, ALLOCATED
```

```
IEC351I CATALOG ADDRESS SPACE MODIFY COMMAND ACTIVE
```

```
IEC348I ALLOCATED CATALOGS 263
```

```
*CAS*****
```

*	FLAGS -VOLSER-USER-CATALOG NAME	%	*
*	Y-I-R- DLB\$02 0001 CATALOG. ZOSR14	4	*
*	Y-I-R- SYSBA1 0001 CATALOG. MCAT. BKUP01	1	*
*	YSI-E- DB2K01 0001 CATALOG. KDB2	22	*
*	Y-I-R- SYSBA1 0001 CATALOG. BXCF	5	*
*	Y-I-R- SYSBA1 0001 CATALOG. BTSO	13	*
*	Y-I-R- SYSBA1 0001 CATALOG. BPRODUCT	11	*
*	Y-I-R- SYS\$A1 0001 CATALOG. BACKUP	1	*
*	Y-I-E- SYSPA2 0001 CATALOG. MCAT. PROD01	1	*

```
*****
```

```
* Y/N-ALLOCATED TO CAS, S-SMS, V-VLF, I-ISC, C-CLOSED, D-DELETED, *
```

```
* R-SHARED, A-ATL, E-ECS SHARED, K-LOCKED *
```

```
*CAS*****
```

```
IEC352I CATALOG ADDRESS SPACE MODIFY COMMAND COMPLETED
```

- Message IEC359I, resulting from the MODIFY CATALOG, REPORT command has been modified to display the current

setting of the extent monitoring threshold, or '(NONE)' if normal monitoring is disabled.

```
F CATALOG, REPORT
IEC351I CATALOG ADDRESS SPACE MODIFY COMMAND ACTIVE
IEC359I CATALOG REPORT OUTPUT 423
*CAS*****
* CATALOG COMPONENT LEVEL      = HDZ11G0      *
* CATALOG ADDRESS SPACE ASN   = 001D          *
* SERVICE TASK UPPER LIMIT    = 180           *
* SERVICE TASK LOWER LIMIT    = 60            *
* HI GHEST # SERVICE TASKS   = 38             *
* CURRENT # SERVICE TASKS    = 38             *
* MAXIMUM # OPEN CATALOGS    = 1, 024         *
* ALIAS TABLE AVAILABLE       = YES            *
* ALIAS LEVELS SPECIFIED      = 1              *
* SYS% TO SYS1 CONVERSION    = OFF            *
* CAS MOTHER TASK             = 009A3920     *
* CAS MODIFY TASK             = 009A3700     *
* CAS ANALYSIS TASK          = 009A0E88     *
* CAS ALLOCATION TASK         = 009A33D8     *
* VOLCAT HI -LEVEL QUALIFIER = SYS1          *
* NOTIFY EXTENT               = 10%           *
* DELETE UCAT/VVDS WARNING    = ON             *
* DATASET SYNTAX CHECKING    = ENABLED        *
*CAS*****
```

It should be understood that this support does not provide any protection or warnings for catalogs that have no secondary space allocation.

Systems Programmer (France)

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Records with duplicate fields in a file

I was asked several times by our production people for a solution to a problem they often encountered. They receive from customers files containing identical fields in several records, and what they need is a file in which those duplicate records are eliminated. The field, identified by a length and an offset, can be in any position within the record. For this elimination of duplicates to take place, the file should be sorted by field (however, this is not always the case because there can be higher-level sort fields, like date, customer number, etc).

So I developed a program that reads a file and creates an output file, eliminating contiguous records with equal field contents and writing out just the first record of each duplicate. The program needs, as parameters, the length and offset of the field to consider. The input file can be a VSAM or a fixed-length sequential. The output file is a sequential. Any sort needed for the input file should be done within the JCL, before running the program.

At the end, the program writes out a message to sysprint stating how many records were read and how many were written.

As an example, here is a piece of JCL, examining a field with length 12 and offset 154:

```
//STEP4    EXEC PGM=DUPPLICAT, PARM='12, 154'  
//INFILE   DD DISP=SHR, DSN=input_file  
//OUTFILE  DD DISP=SHR, DSN=output_file  
//SYSPRINT DD SYSOUT=*
```

DUPPLICAT SOURCE CODE

```
*=====*  
* DUPPLICAT - Find duplicate fields in contiguous records of a file *  
* Parameters: length (max 250) and offset of field. *  
* DDnames: Infile, Outfile, Sysprint *  
* This program reads an input file and copies it to an output file. *  
* If more than one contiguous record has identical fields, only the *  
* first record is written to the output file. *
```

```

* The input file should be sorted by field. It can be a sequential      *
* or a VSAM. The output file must be sequential.                      *
*=====*
&PROGRAM SETC  'DUPLICAT'
&PROGRAM AMODE 31
&PROGRAM RMODE 24
&PROGRAM CSECT
    SAVE  (14, 12)
    LR    R12, R15
    USING &PROGRAM, R12
    ST    R13, SAVEA+4
    LA    R11, SAVEA
    ST    R11, 8(R13)
    LR    R13, R11
    B     GETPARMS
    DC    CL16' &PROGRAM 1.2'
    DC    CL8' &SYSDATE'

*=====*
* Separate input parameter by comma into its components, convert      *
* them to binary form and store them in fields PARM1 and PARM2.      *
* If any of the parms do not exist, the program terminates.           *
*=====*
GETPARMS DS   ØH
    LR  R2, R1          Copy parm pointer to R2.
    L   R2, Ø(Ø, R2)    Load parm address
    LH  R3, Ø(R2)       Load parm length in R3
    OPEN (SYSPRINT, OUTPUT) Open sysprint (for error msgs)
    LTR  R3, R3         Any parm entered?
    BZ   ERRMSGØ        No, error
*
    LR  R6, R2
    AR  R6, R3          R6: point after end of parms
    LA  R6, 2(Ø, R6)    Skip 2 bytes of parmlength
    LA  R2, 2(Ø, R2)
    LR  R4, R2          R4: Current char to ckeck
    LA  R11, PARM1      Area to keep parms
    XR  R9, R9          Clear length counter
*
LOOPARMS EQU  *
    CR  R4, R6          End of parms?
    BNL CONVERT          Yes, go convert the last one
    CLI Ø(R4), C', '
    BE   CONVERT          Comma (separator) found?
    LA   R9, 1(Ø, R9)    Yes, go convert parm
    LA   R4, 1(Ø, R4)    Inc index (char counter)
    B    LOOPARMS         Inc pointer (current char)
                        And continue
*
CONVERT EQU  *
    LTR  R9, R9          Any chars in current parm?

```

BZ	CONVERT2	No, skip pack and cvb instructions
S	R9, =F'1'	Sub one for ex
EX	R9, EXPACK	Execute pack
LA	R9, 1(Ø, R9)	Increment again
CVB	R7, PARMPACK	Convert to binary into R7
ST	R7, Ø(R11)	And store it
*		
CONVERT2	EQU *	
CR	R4, R6	End of all parms?
BNL	CHECKPRM	Yes, jump ahead
AR	R2, R9	Add length to base pointer
LA	R2, 1(Ø, R2)	And skip comma
XR	R9, R9	Reset length
LR	R4, R2	R4: Current char
LA	R11, 4(Ø, R11)	Point next storearea
B	LOOPARMS	
*		
CHECKPRM	CLC PARM1, =F' -1'	Parm1 specified?
BE	ERRMSGØ	No, error
CLC	PARM2, =F' -1'	Same for parm2
BE	ERRMSGØ	
CLC	PARM1, =F' 250'	Length greater than 250?
BH	ERRMSG1	
L	R7, PARM1	R7: Length - 1 (for compare)
S	R7, =F' 1'	
=====		
* Check what input file we have. First open it as VSAM. *		
* If Error, assume non-VSAM file. If VSAM, test ACB for ESDS. *		
* If ESDS, modify RPL accordingly. *		
* Output file must be a sequential. *		
=====		
OPENACB1	EQU *	Open ACB for VSAM input file
	OPEN INFLEA	If error, go open sequential
	LTR R15, R15	
	BNZ OPENDCB1	
	TESTCB ACB=INFLEA,	X
	ATRB=ESDS	Check if VSAM ESDS
	BNE OPENDCB2	No, go open output file
*		
ESDSFILE1	EQU *	
	MODCB RPL=INFILER,	X
	OPTCD=ADR	Modify RPL for ESDS
	B OPENDCB2	
*		
OPENDCB1	EQU *	Open sequential input file
	OPEN (INFLED, INPUT)	
	LTR R15, R15	
	BNZ ERRMSG2	
	MVI FILETYP1, C' S'	Set flag sequential (nonvsam)

LA R2, INFILED	Address IHADCB of input file with R2
USING IHADCB, R2	
*	
OPENDCB2 EQU *	Open output file
OPEN (OUTFILE, OUTPUT)	
LTR R15, R15	
BNZ ERMSG3*	
=====	=====
* Read and compare loop	*
*=====	=====**
READFILE EQU *	
XR R8, R8	Record count for input file
XR R9, R9	Record count for output file
L R3, PARM2	R3 is last position of string to
A R3, PARM1	compare. Lrec1 cannot be smaller
*	
READLOOP EQU *	
CLI FILETYP1, C' V'	VSAM file?
BNE READSEQ1	No, go to sequential
*	
READVSA1 EQU *	
GET RPL=INFLER	Read VSAM file
LTR R15, R15	End of file?
BNZ EXIT0	
L R4, VAREA1	Get address of data in R4.
SHOWCB RPL=INFLER,	
AREA=LRECL1,	
LENGTH=4,	
FIELD=RECLEN	
L R5, LRECL1	Get record length in R5
B COMPARE	and jump to compare
*	
READSEQ1 EQU *	
GET INFILED	Read sequential (Locate method)
LR R4, R1	copy address of data to R4.
LH R5, DCBLRECL	Load R5 with record length.
*	
COMPARE EQU *	
LA R8, 1(0, R8)	Increment input record counter
CR R5, R3	Record smaller than last position?
BL ERMSG4	Yes, error.
LR R6, R4	
A R6, PARM2	Add offset to record address
C R8, =F'1'	First record has no previous
BE COMPARE1	string to compare with, so skip
EX R7, EXCOMPAR	Execute compare
BE COMPARE2	If strings equal, do not write
*	
COMPARE1 EQU *	

```

        PUT    OUTFILE, (R4)          Write sequential
        LA     R9, 1(Ø, R9)          Increment output counter
*
COMPARE2 EQU   *
        EX    R7, EXMOVE           Move to string
        B     READLOOP             and read next
*=====
* Send final messages, close files, and exit
*=====
EXIT0   EQU   *
        LR    RØ, R8
        BAL   R1Ø, UNPACK
        MVC   ENDNUM1, OUT1Ø
        LR    RØ, R9
        BAL   R1Ø, UNPACK
        MVC   ENDNUM2, OUT1Ø
        PUT   SYSPRINT, ENDMMSG1
        PUT   SYSPRINT, ENDMMSG2
*
EXIT1   EQU   *
        CLOSE INFILED
        CLOSE INFILEA
        CLOSE OUTFILE
        CLOSE SYSPRINT
        L     R13, SAVEA+4
        LM   R14, R12, 12(R13)
        XR   R15, R15
        BR   R14
*=====
* Subroutines and work areas
*=====
EXCOMPAR EQU   *
        CLC   Ø(Ø, R6), STRING
*
EXMOVE   EQU   *
        MVC   STRING, Ø(R6)
*
EXPACK   EQU   *
        PACK  PARMPACK, Ø(Ø, R2)
*
UNPACK   EQU   *
        CVD   RØ, REGDECIM
        UNPK  OUT12, REGDECIM
        BR   R1Ø
*
ERRMSGØ  EQU   *
        PUT   SYSPRINT, =CL8Ø' > Parameters missing'
        B     EXIT1
ERRMSG1  EQU   *

```

	PUT	SYSPRINT, =CL80' > Parm1 (length) exceeds limit of 250'	
	B	EXIT1	
ERRMSG2	EQU	*	
	PUT	SYSPRINT, =CL80' > Error opening input file'	
	B	EXIT1	
ERRMSG3	EQU	*	
	PUT	SYSPRINT, =CL80' > Error opening output file'	
	B	EXIT1	
ERRMSG4	EQU	*	
	PUT	SYSPRINT, =CL80' Record smaller than compare position'	
	PUT	SYSPRINT, =CL80' at last read record. Program terminated'	
	B	EXIT0	
*			
ENDMSG1	DC	C' Number of records read from INFILE . . :'	
ENDNUM1	DS	CL10	
	DC	CL40' '	
ENDMSG2	DC	C' Number of records written to OUTFILE . . :'	
ENDNUM2	DS	CL10	
	DC	CL40' '	
STRING	DS	0C	
STRING1	DS	CL250	
*			
INFILEA	ACB	DDNAME=INFILE	
INFILER	RPL	ACB=INFILEA, OPTCD=LOC, AREA=VAREA1, ARG=CHAVE1	X X X
*			
INFILED	DCB	DSORG=PS, MACRF=(GL), EODAD=EXIT0, DDNAME=INFILE	X X
*			
OUTFILE	DCB	DSORG=PS, MACRF=(PM), DDNAME=OUTFILE	X
*			
SYSPRINT	DCB	DSORG=PS, MACRF=(PM), LRECL=80, DDNAME=SYSPRINT	X X
*			
SAVEA	DS	18F	
VAREA1	DS	F	
CHAVE1	DS	F	
LRECL1	DS	F	
FILETYP1	DC	C' V'	
PARMPACK	DS	D	
PARM1	DC	F' -1'	
PARM2	DC	F' -1'	
	DS	0D	
REGDECIM	DS	CL9	
	DS	0F	

```
OUT12    DS      OCL12
OUT10    DS      CL10
          DS      CL2
*
LTORG
DCBD   DSORG=PS
YREGS
END
```

A peek into SMF30 data

Many performance problems in an application can be quickly solved once the bottlenecks are identified. Performance monitoring tools help us identify bottlenecks, but what if a site doesn't have those expensive tools? What if a site is planning to get rid of its performance tools to cut down costs? Sometimes it is difficult to justify the investment in performance monitoring tools because many sites using them do not use a lot of the information that they provide.

Here is a small utility program that would interpret the SMF 30 records (address space-level accounting record) and produce a report that could help identify candidates for performance tuning. This program can be customized according to your requirements – such as populating a database for performance analysis or producing customized reports. We have utilized this in the past by populating an Access database. The information in the reports is limited to analysing batch workloads.

SOFTWARE DEPENDENCIES

This program works well with the SMF30 records produced with Version 2 Release 10 of OS/390.

Compile this program after concatenating the SYS.MACLIB provided with your release of OS/390. SMF type 30 layout will be

expanded using the IFASMFR macro and the program is set to produce the report using SMF records as input.

GETTING STARTED

Although SMF data given as is to this program can contain any of the SMF record types, it will utilize only the type 30 records with sub-types 4 and 5. It is therefore preferable to filter out those records using IFASMFDP, and SORT these in order of JOBNAMES/STEP name so that the step level reports generated are more readable. JCL to extract the SMF 30 (subtypes 4 and 5) and subsequently SORT the SMF 30 records is given below:

```
//**** Place Job Card Here
//STEP010 EXEC PGM=IFASMFDP
//SYSPRINT DD SYSOUT=*
//SYSUDUMP DD SYSOUT=*
//DUMPIN DD DISP=SHR, DSN=MY. SMF. DATASET
//DUMPOUT DD DSN=&&OUTSMF, DISP=(NEW, PASS),
//           DCB=(LRECL=32760, RECFM=VBS), SPACE=(CYL, (50, 20), RLSE)
//SYSIN DD *, DCB=BLKSIZE=80
  INDD(DUMPIN, OPTIONS(DUMP))
  OUTDD(DUMPOUT, TYPE(30(4, 5)))
START(0600)
END(2200)
/*
//STP020 EXEC PGM=SORT
//SYSPRINT DD SYSOUT=*
//SYSOUT DD SYSOUT=*
//SORTIN DD DISP=(OLD, DELETE), DSN=&&OUTSMF
//SORTOUT DD DSN=MY. SORTED. SMF. DATASET, DISP=(NEW, CATLG),
//           DCB=(LRECL=32760, RECFM=VBS), SPACE=(CYL, (50, 20), RLSE)
//SYSIN DD *
  SORT FIELDS=(247, 8, CH, A, 215, 8, CH, A, 255, 2, BI, A)
/*
```

As one may want to focus on a particular set of jobs, this utility allows the user to limit the report output by using masks at the job level. The include criterion is specified on the SYSIN card and can contain the exact job name; alternatively, it can include wildcard characters such as '*' to indicate 0 or more character match, or '?' to indicate exactly one character match. Multiple criteria can be given in the SYSIN card. All jobs satisfying the criteria on the SYSIN cards will be included in the generated reports.

The JCL to run this utility program is shown below:

```
//RUN      EXEC PGM=SMF30PG
//STEPLIB DD DISP=SHR,DSN=<Loadlib containing the utility program>
//SYSIN    DD *
JOBNAME=PRD*
//SMFIN   DD DISP=SHR,DSN=MY.SORTED.SMF.DATASET
//STEPSU   DD SYSOUT=*
//JOBSU    DD SYSOUT=*
//STEPTIME DD SYSOUT=*
//JOBTIME  DD SYSOUT=*
//STEPSTOR DD SYSOUT=*
```

The reports produced by the sample JCL show only the jobs that start with PRD (by providing the JOBNAME=PRD* in SYSIN DD card).

REPORTS

This utility produces five reports that can be used to analyse the address space accounting information – DD cards JOBSU, STEPSU, JOBTIME, STEPTIME, and STEPSTOR.

A report specified by the DD card JOBTIME contains the following information at the job level:

- Job name and job number.
- Job start date (format YYYY-MM-DD) and start time (format HH:MM:SS).
- Job end date (format YYYY-MM-DD) and end time (format HH:MM:SS).
- Total CPU time taken by the job (format HH:MM:SS).
- SRB time taken by the job (format HH:MM:SS).
- Total EXCP count for the job.
- WLM class and WLM service name.

A report specified by the DD card JOBSU contains similar information to JOBTIME except that total CPU time and SRB time is expressed in terms of service units. EXCP count is expressed in terms of I/O service units.

A report specified by the DD card STEPTIME contains the following information:

- Job name and job number.
- Step start date (format YYYY-MM-DD) and start time (format HH:MM:SS).
- Step end date (format YYYY-MM-DD) and end time (format HH:MM:SS).
- Step name.
- Program name.
- CPU time taken by the step (format HH:MM:SS).
- SRB time taken by the step (format HH:MM:SS).
- EXCP count for the step.

Report specified by the DD card STEPSU contains similar information to STEPTIME except that CPU time and SRB time are expressed in terms of service units. Additionally I/O service units are also reported.

A report specified by STEPSTOR contains information about memory usage at the step level. The following information is contained in the report:

- Job name and job number.
- Step name.
- Program name.
- Step start date (format YYYY-MM-DD) and start time (format HH:MM:SS).
- Step end date (format YYYY-MM-DD) and end time (format HH:MM:SS).
- User storage below the 16MB line (in KB).
- User storage above the 16MB line (in KB).

- System storage below the 16MB line (in KB).
- System storage above the 16MB line (in KB).

The output format generated by this utility is easily portable to a spreadsheet or a database like Microsoft Access. Once the data is imported to an office tool like Excel or Access, various in-built functions (like ORDERING or SORTING, Data/Time functions) can be used to quickly get to valuable information like identifying the top consumers of CPU time, identifying long-running jobs, charting the trends over a period of time, or capacity planning. With good job naming conventions, the information can also be used for billing or charge-back purposes. Looking at resource usage in terms of service units may also be useful when moving jobs from one system to other. The STEPSTOR report can be used to see whether the program is using too much storage below the line. If yes, it can be re-compiled with DATA (31) option (if COBOL) to relieve storage below the line. In heavily paging environments, one may want to look at long-running programs consuming high storage to see whether there is potential to reduce their storage requirements. On the other hand, if the paging activity is negligible (or UIC stays at 255), one may want to look at exploiting the storage to reduce elapsed/CPU time for the programs showing high EXCP counts. Depending on what those programs do, it may be possible to use various buffering techniques like batch LSR or System Managed Buffering (SMB).

SMF30

```

SMF30PG TITLE 'A SAMPLE SMF30 REPORT WRITER PROGRAM'
SMF30PG CSECT
SMF30PG AMODE 31
SMF30PG RMODE ANY
MACRO
$STRGACQ &TYPE, &LEN=, &LOC=
AI F (' &TYPE' EQ 'R').L0001
L R0, &LEN
AGO .L0002
.L0001 ANOP
LR R0, &LEN
.L0002 ANOP
A R0, =F' 4'

```

```

STORAGE OBTAIN, LENGTH=(0), LOC=&LOC
ST    R0, 0(, R1)
LA    R1, 4(, R1)
MEND
MACRO
$STRGREL &TYPE, &ADDR=
AIF   (' &TYPE' EQ 'R'). L0001
L    R1, &ADDR
AGO  . L0002
.L0001 ANOP
LR    R1, &ADDR
.L0002 ANOP
S    R1, =F' 4'
L    R0, 0(, R1)
STORAGE RELEASE, ADDR=(1), LENGTH=(0)
MEND
*
YREGS
USING SMF30PG, R12
$CONS DS 0H
B   $CONSL(R15)
DC  C' SMF30PG &SYSDATE. &SYSTIME. '
DS  0H
$CONSL EQU *-$CONS
STM  R14, R12, 12(R13)      Store caller information
LR   R12, R15                R12 Base
LR   R2, R1                  Save R1 for later
$STRGACQ A, LEN==A($DYNL), LOC=31 Acq Dyn Stor
ST   R13, 4(, R1)            Save Back pointer
ST   R1, 8(, R13)           Save forward pointer
LR   R13, R1                 My save area
USING $DYN, R13             Let Assembler know
BAS  R14, OPEN#000           Open all the files.
LTR  R15, R15                Mandatory files opened ?
BNZ  RETURN4                N: bail out
LA   R1, SYSIN               point to SYSIN offset
M    R0, =A(FILESIZE)        Get the real offset
A    R1, FILEDATA            Add start
MVI  INCL, INCLALL          Default to include all
CLI  FILEIND-FILEINFO(R1), X'00' Was the file opened ?
BE   SMFP#020                N:
L    R2, FILEREC-FILEINFO(, R1) Get the record address
L    R3, FILEDCB-FILEINFO(, R1) Get the SYSIN DCB Addr
SMFP#005 DS 0H
LR   R0, R2                  Move record ADDR to R0
GET  (3)                    Get a record
CLC  0(08, R2), =C'JOBNAME=' Is it a mask ?
BNE  SMFP#005                N:
CLI  08(R2), C' '            Is there a mask ?
BE   SMFP#005                N:

```

XC	I NCL, I NCL	Indicate mask present
\$STRGACQ	A, LEN==A(MASKL), LOC=31	Acq stor for mask
MVC	MASKNEXT-MASKDSCT(4, R1), MASKLIST	Store next adr
ST	R1, MASKLIST	Store anchor
MVC	MASKNAME-MASKDSCT(8, R1), Ø8(R2)	Store Mask
LA	R4, Ø8(, R2)	Bump past literal
LA	R5, Ø8(, R2)	Bump past literal
LA	R6, 8	Length of jobname
SMFP#006	DS ØH	
CLI	Ø(R4), C' '	is this char blank ?
BE	SMFP#007	Y:
LA	R4, 1(, R4)	bump past this char
BCT	R6, SMFP#006	loop back
SMFP#007	DS ØH	
SR	R4, R5	subtract start address
LTR	R4, R4	is it zero ?
BZ	SMFP#005	y: back to read
ST	R4, MASKNAML-MASKDSCT(R1)	store length
B	SMFP#005	read next record
SMFP#010	DS ØH	
SMFP#015	DS ØH	
SMFP#020	DS ØH	
LA	R1, SMFIN	process smf file
M	RØ, =A(FI LESIZ)	
A	R1, FILEDATA	
CLI	FILEIND-FILEINFO(R1), X' 00'	is the file open ?
BE	RETURN4	n: go back
L	R2, FILEREC-FILEINFO(, R1)	get the record address
L	R3, FILEDCB-FILEINFO(, R1)	get the dcb address
SMFP#025	DS ØH	
LR	RØ, R2	store record address in RØ
GET	(3)	get a record
MVI	STEP, ISSTEP	indicate this is a step rec
CLI	SMF3ØRTY-SMFRCRD3Ø(R2), X' 1E'	is this rec type 3Ø ?
BNE	SMFP#025	
CLC	SMF3ØSTP-SMFRCRD3Ø(2, R2), =H' 04'	is this a subtype 4 ?
BL	SMFP#025	
CLC	SMF3ØSTP-SMFRCRD3Ø(2, R2), =H' 05'	is this a subtype 5 ?
BH	SMFP#025	
CLC	SMF3ØSTP-SMFRCRD3Ø(2, R2), =H' 04'	
BE	SMFP#251	
MVI	STEP, ISJOB	
SMFP#251	DS ØH	
LA	R1, SMF3ØDTE-SMFRCRD3Ø(R2)	get the end date
BAS	R14, CNVD#ØØØ	convert date
MVC	ENDDATE, WORKDATE+2	save date for later
ICM	R1, 15, SMF3ØTME-SMFRCRD3Ø(R2)	get the end time
BAS	R14, CNVT#ØØØ	convert time
MVC	ENDTIME, WORKTIME+2	save time for later
CLC	SMF3ØIOF-SMFRCRD3Ø(4, R2), =F' 0'	Inden section present ?

BE	SMFP#026	
CLC	SMF301LN-SMFRC30(2, R2), =H'0'	I den section present ?
BE	SMFP#026	
CLC	SMF301ON-SMFRC30(2, R2), =H'0'	I den section present ?
BE	SMFP#026	
L	R4, SMF301OF-SMFRC30(R2)	map the i den section
AR	R4, R2	
MVC	JOBNAME, SMF30JBN-SMF30ID(R4)	get jobname
CLI	INCL, INCLALL	is this include all ?
BE	SMFP#254	n:
L	R5, MASKLIST	point to mask list
SMFP#253	DS ØH	
	LTR R5, R5	is one present ?
	BZ SMFP#025	n: get another record
	BAS R14, CHEK#000	check the mask
	LTR R15, R15	check good ?
	BZ SMFP#254	y: process record
	L R5, MASKNEXT-MASKDSCT(, R5)	get next mask
	B SMFP#253	check again
SMFP#254	DS ØH	
	MVC JOBNUM, SMF30JNM-SMF30ID(R4)	get the job number
	MVC STEPNAME, SMF30STM-SMF30ID(R4)	get the step name
	MVC PROGRAM, SMF30PGM-SMF30ID(R4)	get the program name
	LA R1, SMF30STD-SMF30ID(R4)	Get start date
	BAS R14, CNVD#000	Convert date
	MVC STRTDTE, WORKDATE+2	Save for later
	ICM R1, 15, SMF30SIT-SMF30ID(R4)	Get start time
	BAS R14, CNVT#000	Convert Time
	MVC STRTTME, WORKTIME+2	Save time for later
SMFP#026	DS ØH	
	BAS R14, HDR#000	Build header info
	BAS R14, SUDT#000	Build SU details
	BAS R14, TMDT#000	Build Time details
	BAS R14, STDT#000	Build Storage details
	B SMFP#025	Get next record
SMFP#030	DS ØH	
SMFP#035	DS ØH	
SMFP#040	DS ØH	
	B RETURNØ	
RETURNØ	DS ØH	
	LA R5, Ø	Process OK. RC=0
	B RETURN	
RETURN4	DS ØH	
	LA R5, 4	Process not OK. RC=4
RETURN	DS ØH	
	BAS R14, CLOS#000	Close all files.
	LR R1, R13	Save current dyn area
	L R13, 4(, R13)	Restore old dyn area
	\$STRGREL R, ADDR=1	Release curr dyn area
	L R14, 12(, R13)	Restore R14

LR	R15, R5	R15->RC
LM	R0, R12, 20(R13)	Restore other regs
BSM	Ø, R14	goback
*	Check if the pattern mask matches the jobname	
CHEK#000	BAKR R14, Ø	
	LA R2, MASKNAME-MASKDSCT(R5)	Get address of Mask
	LA R3, MASKNAML-MASKDSCT(R5)	Point to length of mask
	LA R4, JOBNAM	Point to current jobname
	ASAXWC PATTERNSTR=(2), PATTERNSTRLEN=(3), STRING=(4), STRINGLEN=A(8), ZEROORMORE=Z, ONECHAR=0, RETCODE=RETCODE, WORKAREA=WORKAREA, MF=(E, MYLIST)	- - - - - - - - -
	L R15, RETCODE	RC=Ø, match, else nomatch
CHEK#999	PR	Go back to caller
*	Open all files.	
OPEN#000	BAKR R14, Ø	
	SLR RØ, RØ	Initialize RØ
	LA R1, FILESIZE	Get the size of 1 file rec
	LA R3, DDINFOS	Get no. of DDs
	MR RØ, R3	Get size required
	LR RØ, R1	
	\$STRGACQ R, LEN=Ø, LOC=31	Acq Dyn Stor
	ST R1, FILEDATA	Store returned addr
	SLR R2, R2	
	LA R3, DDINFO	Point to DDINFO
	LR R4, R1	
	LA R5, DDINFOS	Get no. of DDs
	USING FILEINFO, R4	
OPEN#005	DS ØH	
	LA RØ, MODOPENL	Get length of OPEN MODEL
	\$STRGACQ R, LEN=Ø, LOC=31	ACQ dyn area
	ST R1, FILEOPEN	Store returned addr
	MVC Ø(MODOPENL, R1), MODOPEN	Move OPEN MODEL
	LA RØ, MODCLOS	Get Length of CLOSE MODEL
	\$STRGACQ R, LEN=Ø, LOC=31	ACQ dyn area
	ST R1, FILECLOS	Store returned addr
	MVC Ø(MODCLOS, R1), MODCLOS	Move close model
	LA RØ, MODDCBEL	Get length of DCBE model
	\$STRGACQ R, LEN=Ø, LOC=31	ACQ dyn area
	ST R1, FILEDCBE	Store returned addr
	MVC Ø(MODDCBEL, R1), MODDCBE	Move DCBE model
	USING DCBE, R1	
	MVC DCBEEODA, 16(R3)	Save EOD address
	MVC DCBESYNA, 2Ø(R3)	Save SYNAD address

DROP R1	
CLI 8(R3), C' I'	Is this open for Input y:
BNE OPEN#010	Get input DCB length
LA R0, MODDCBIL	ACQ dyn area
\$STRGACQ R, LEN=0, LOC=24	Store returned addr
ST R1, FILEDCB	Move input DCB model
MVC 0(MODDCBIL, R1), MODDCB1	
B OPEN#015	
OPEN#010 DS 0H	
LA R0, MODDCBOL	Get length of output DCB
\$STRGACQ R, LEN=0, LOC=24	ACQ dyn area
ST R1, FILEDCB	Store returned addr
MVC 0(MODDCBOL, R1), MODDCBO	Move output DCB model
OPEN#015 DS 0H	
USING IHADCB, R1	
MVC DCBDDNAM, 0(R3)	Move DD name
MVC DCBDCBE, FILEDCBE	Move DCBE address
DROP R1	
L R0, 12(, R3)	Move length of record
\$STRGACQ R, LEN=0, LOC=31	ACQ dyn area
ST R1, FILEREC	Store returned addr
MVC FILERECL, 12(R3)	save lrecl
XC FILEIND, FILEIND	initialize indicator
MVC FILEHEDR, 24(R3)	move header address
L R6, FILEDCB	point to file dcb
L R7, FILEOPEN	point to open area
CLI 8(R3), C' I'	is this open for input ?
BNE OPEN#020	n:
OPEN ((6), INPUT), MF=(E, (7)), MODE=31	open file
B OPEN#025	
OPEN#020 DS 0H	
OPEN ((6), OUTPUT), MF=(E, (7)), MODE=31	open file
OPEN#025 DS 0H	
LTR R15, R15	Open success ?
BZ OPEN#035	Y:
CLI 9(R3), C' M'	Was this a mandatory file ?
BNE OPEN#040	n:
LA R2, 4	indicate RC=4
B OPEN#040	
OPEN#035 DS 0H	
MVI FILEIND, X' 01'	Indicate file open
CLI 8(R3), C' I'	Was this input ?
BE OPEN#040	y:
LR R1, R4	point to offset
BAS R14, HDRW#000	Write header record
OPEN#040 DS 0H	
LA R3, DDINFOL(, R3)	Bump to next dd
LA R4, FILESZ(, R4)	Bump to next file
BCT R5, OPEN#005	open them

	LR	R15, R2	R15->RC
OPEN#999	PR		
	DROP	R4	
*	Write header records to output file		
HDRW#000	BAKR	R14, Ø	
	LR	R4, R1	Point to offset
	L	RØ, FILEREC-FILEINFO(R4)	Get record addr
	L	R3, FILEDCB-FILEINFO(R4)	Get DCB addr
	L	R1, FILERECL-FILEINFO(R4)	Get record length
	L	R14, FILEHDR-FILEINFO(R4)	Get header rec addr
	LR	R15, R1	length in R15
	ICM	R15, B'1000', BLANK	Default to ''
	LTR	R14, R14	Is header present ?
	BZ	HDRW#999	n: bail out
	MVCL	RØ, R14	Move record to output
	L	RØ, FILEREC-FILEINFO(R4)	point to record addr
	PUT	(3)	Write output
			goback
HDRW#999	PR		
*	Close all files.		
CLOS#000	BAKR	R14, Ø	
	L	R2, FILEDATA	Point to file information.
	LA	R3, DDINFOS	Get number of DDs
	USING	FILEINFO, R2	
CLOS#005	DS	ØH	
	L	R6, FILEDCB	Get DCB addr
	L	R7, FILECLOS	Get close area
	CLI	FILEIND, X'ØØ'	Is this file open ?
	BE	CLOS#Ø1Ø	n:
	CLOSE	((6)), MF=(E, (7)), MODE=31	Close file
CLOS#010	DS	ØH	
	L	R1, FILEOPEN	Point to open area
	\$STRGREL	R, ADDR=1	Release storage
	L	R1, FILECLOS	Point to close area
	\$STRGREL	R, ADDR=1	Release storage
	L	R1, FILEDCB	Point to dcbe area
	\$STRGREL	R, ADDR=1	Release storage
	L	R1, FILEDCBE	Point to dcbe area
	\$STRGREL	R, ADDR=1	Release storage
	L	R1, FILEREC	Point to rec area
	\$STRGREL	R, ADDR=1	Release storage
	MVI	FILEIND, X'ØØ'	Indicate file closed
	LA	R2, FILESIZ(, R2)	Bump to next
	BCT	R3, CLOS#005	Close them
			go back
CLOS#999	PR		
	DROP	R2	
PRTF#000	BAKR	R14, Ø	Write output to files
	M	RØ, =A(FILESIZE)	R1 points to offset
	A	R1, FILEDATA	add start offset
	CLI	FILEIND-FILEINFO(R1), X'ØØ'	is this file open ?

BE	PRTF#999	n: goback
L	R2, FILEREC-FILEINFO(R1)	Point to record addr
L	R3, FILEDCB-FILEINFO(R1)	Point to DCB
L	R4, FILERECL-FILEINFO(R1)	Get record length
LR	R0, R2	
LR	R1, R4	
LA	R14, PRTLINE	Point to PRTLINE
LA	R15, L' PRTLINE	Get PRTLINE address
ICM	R15, B'1000', BLANK	Default to ''
MVCL	R0, R14	Move record to output
LR	R0, R2	R0-> record
PUT	(3)	Put to record to file/. goback
PRTF#999	PR	
HEDR#000	BAKR R14, Ø	write header part of record
	MVC PRTLINE(L' BLANK), BLANK	initialize prtline
	MVC PRTLINE+1(L' PRTLINE-1), PRTLINE	
	LA R1, PRTLINE	point to prtline
USING	HEDR, R1	let assembler know
MVC	HEDRJNME, JOBNAM	move jobname
MVC	HEDRJNUM, JOBNUM	move jobnumber
MVC	HEDRSTDT, STRTDTE	move start date
MVC	HEDRSTTM, STRTTME	move start time
MVC	HEDRENDT, ENDDATE	move end date
MVC	HEDRENTM, ENDTIME	move end time
CLI	STEP, ISJOB	is this step info ?
BE	HEDR#999	n: goback
MVC	HEDRSTNM, STEPNAM	move stepname
MVC	HEDRPGNM, PROGRAM	move program name
HEDR#999	PR	goback
SUDT#000	BAKR R14, Ø	write SU details
	BAS R14, HEDR#000	write header part of rec
USING	SMFRCD3Ø, R2	let assembler know
CLC	SMF3ØPOF, =F' Ø'	is there info ?
BE	SUDT#999	n: goback
CLC	SMF3ØPLN, =H' Ø'	is there info ?
BE	SUDT#999	n: goback
CLC	SMF3ØPON, =H' Ø'	is there info ?
BE	SUDT#999	n: goback
A	R2, SMF3ØPOF	Add offset to section.
USING	SMF3ØPRF, R2	Let assembler know
LA	R1, PRTLINE	Point to prtline
CLI	STEP, ISJOB	is this job information ?
BE	SUDT#005	y:
LA	R1, HEDRSTEP(, R1)	Bump step header info
LA	R3, STEPSU	indicate stepsu
B	SUDT#010	
SUDT#005	DS ØH	
	LA R1, HEDRJOB(, R1)	Bump past job header info
	LA R3, JOBSU	indicate jobsu

SUDT#010	DS	ØH	
	USING	SUDSECT, R1	Let assembler know
	MVC	SUTOT, EDMASKN	Move mask to su totals
	MVC	SUCPU, EDMASKN	Move mask to su cpu
	MVC	SUSRB, EDMASKN	Move mask to su srb
	MVC	SUIO, EDMASKN	Move mask to su io
	ICM	R14, 15, SMF3ØSRV	get total su
	CVD	R14, DWORD	conver to decimal
	ED	SUTOT, DWORD+3	edit to su totals
	ICM	R14, 15, SMF3ØCSU	get cpu su
	CVD	R14, DWORD	conver to decimal
	ED	SUCPU, DWORD+3	edit to su cpu
	ICM	R14, 15, SMF3ØSRB	get srb su
	CVD	R14, DWORD	conver to decimal
	ED	SUSRB, DWORD+3	edit to su srb
	ICM	R14, 15, SMF3ØIO	get io su
	CVD	R14, DWORD	conver to decimal
	ED	SUIO, DWORD+3	edit to su io
	CLI	STEP, ISTEP	is this step info ?
	BE	SUDT#015	y: goback
	MVC	SUWLM, SMF3ØWLM	move wlm class
	MVC	SUSRVN, SMF3ØSCN	move wlm service name
	MVC	JOBWLMS, SMF3ØWLM	move wlm class
	MVC	JOBWLMS, SMF3ØSCN	move wlm service name
SUDT#015	DS	ØH	
	LR	R1, R3	point to dd offset
	BAS	R14, PRTF#000	write this record
SUDT#999	PR		goback
	DROP	R1, R2	
TMDT#000	BAKR	R14, Ø	write TIME details
	BAS	R14, HEDR#000	write header part of rec
	USING	SMFRCD3Ø, R2	let assembler know
	CLC	SMF3ØCOF, =F' Ø'	is there info ?
	BE	TMDT#999	n: goback
	CLC	SMF3ØCLN, =H' Ø'	is there info ?
	BE	TMDT#999	n: goback
	CLC	SMF3ØCON, =H' Ø'	is there info ?
	BE	TMDT#999	n: goback
	ZAP	DEXCP, =P' Ø'	Initialize excp count
	CLC	SMF3ØUOF, =F' Ø'	is there info ?
	BE	TMDT#003	n: goback
	CLC	SMF3ØULN, =H' Ø'	is there info ?
	BE	TMDT#003	n: goback
	CLC	SMF3ØUON, =H' Ø'	is there info ?
	BE	TMDT#003	n: goback
	LR	R3, R2	
	A	R3, SMF3ØUOF	
L	R1, SMF3ØTEP-SMF3ØURA(, R3)	Get excp count	
CVD	R1, DEXCP	Convert to decimal	

TMDT#003	DS	ØH	
	A	R2, SMF30COF	Add offset
	USING	SMF30CAS, R2	Let assembler know
	LA	R4, PRTLINE	Point to prtline
	CLI	STEP, ISJOB	is this job information ?
	BE	TMDT#005	y:
	LA	R4, HEDRSTEP(, R4)	bump past step header
	LA	R3, STEPTIME	indicate steptime dd
	B	TMDT#010	
TMDT#005	DS	ØH	
	LA	R4, HEDRJOB(, R4)	bump past job header
	LA	R3, JOBTIME	indicate jobtime
TMDT#010	DS	ØH	
	USING	TIMEDSCT, R4	Let assembler know
	MVC	TIMEEXCP, EDMASKN	move edit mask to excp
	ICM	R1, 15, SMF30CPT	Get tcb cpu time
	BAS	R14, CNVT#000	convert time
	MVC	TIMECPU, WORKTIME+2	write cpu time to output
	ICM	R1, 15, SMF30CPS	Get tcb srb time
	BAS	R14, CNVT#000	convert time
	MVC	TIMESRB, WORKTIME+2	write srb time to output
	ED	TIMEEXCP, DEXCP+3	Edit excps
	CLI	STEP, ISSTEP	Is this step info ?
	BE	TMDT#015	
	MVC	TIMEWLM, JOBWLMC	
	MVC	TIMESRVN, JOBWLVMS	
TMDT#015	DS	ØH	
	LR	R1, R3	point to dd offset
	BAS	R14, PRTF#000	write this record
TMDT#999	PR		goback
	DROP	R4, R2	
STDT#000	BAKR	R14, Ø	write storage details
	BAS	R14, HEDR#000	write header part of output
	USING	SMFRCD30, R2	Let assembler know
	CLC	SMF30ROF, =F' Ø'	is there info ?
	BE	STDT#999	n: goback
	CLC	SMF30RLN, =H' Ø'	is there info ?
	BE	STDT#999	n: goback
	CLC	SMF30RON, =H' Ø'	is there info ?
	BE	STDT#999	n: goback
	A	R2, SMF30ROF	Add offset
	USING	SMF30SAP, R2	Let assembler know
	LA	R4, PRTLINE	point to prtline
	CLI	STEP, ISJOB	is this job info ?
	BE	STDT#999	y: goback
	LA	R4, HEDRSTEP(, R4)	bump step header info
	LA	R3, STEPSTOR	indicate stepstor
	USING	STORDSCT, R4	Let assembler know
	MVC	STORUB16, EDMASKN	Move edit mask to UStor<16M

MVC	STORUA16, EDMASKN	Move edit mask to UStor>16M
MVC	STORSB16, EDMASKN	Move edit mask to SStor<16M
MVC	STORSA16, EDMASKN	Move edit mask to SStor>16M
ICM	R14, B'1111', SMF30URB	Get Ustor<16M
SRL	R14, 10	Convert to KB
CVD	R14, DWORD	Convert to decimal
ED	STORUB16, DWORD+3	Edit to output area
ICM	R14, B'1111', SMF30EUR	Get Ustor>16M
SRL	R14, 10	Convert to KB
CVD	R14, DWORD	Convert to decimal
ED	STORUA16, DWORD+3	Edit to output area
ICM	R14, B'1111', SMF30ARB	Get Sstor<16M
SRL	R14, 10	Convert to KB
CVD	R14, DWORD	Convert to decimal
ED	STORSB16, DWORD+3	Edit to output area
ICM	R14, B'1111', SMF30EAR	Get Sstor>16M
SRL	R14, 10	Convert to KB
CVD	R14, DWORD	Convert to decimal
ED	STORSA16, DWORD+3	Edit to output area
LR	R1, R3	Point to dd offset
BAS	R14, PRTF#000	write this output rec
STDT#999	PR	goback
	DROP R4, R2	
CNVT#000	BAKR R14, Ø	Convert time to HH:MM:SS
	SLR RØ, RØ	R1-> time in 100th of sec
	D RØ, =F'360000'	Get hours
	CVD R1, DWORD	Convert to decimal
	SRP DWORD, 4, Ø	00000000HH0000C
	ZAP WTIME, DWORD	Save it for later
	LR R1, RØ	Get reminder
	SLR RØ, RØ	
	D RØ, =F'6000'	Get minutes
	CVD R1, DWORD	Convert to decimal
	SRP DWORD, 2, Ø	000000000000MM00C
	AP WTIME, DWORD	Add to saved time
	LR R1, RØ	Get reminder
	SLR RØ, RØ	
	D RØ, =F'100'	Get seconds
	CVD R1, DWORD	Convert to decimal
	AP WTIME, DWORD	Add to saved time
	MVC WORKTIME, EDMASKT	Move edit mask for time
	ED WORKTIME, WTIME+4	edit time
CNVT#999	PR	goback
*	Convert date from ØnYYDDDF to CCYY-MM-DD format	
CNVD#000	BAKR R14, Ø	
	ZAP DWORD, Ø(4, R1)	R1 points to input date
	OI DWORD+L' DWORD-1, X'ØF'	Make pack!
	AP DWORD, =P'1900000'	Add 1900000
	ZAP WDATE, DWORD	Save for later

	SRP	DWORD, 64-3, Ø	Remove days
	CVB	R1, DWORD	Convert to binary
	SLR	RØ, RØ	
	D	RØ, =F' 4'	Divide by 4
	C	RØ, =F' Ø'	Any remainder ?
	BE	CNVD#ØØ5	y: is leap year
	L	R1, =A(NLEAPYR)	not a leap year
	B	CNVD#Ø1Ø	
CNVD#ØØ5	DS	ØH	
	L	R1, =A(LEAPYR)	is a leap year
CNVD#Ø1Ø	DS	ØH	
	ZAP	DWORD, WDATE	restore date
	SRP	WDATE, 64-3, Ø	remove days
	SRP	WDATE, 4, Ø	ØØØØØØØCCYYØØØØF
	SRP	DWORD, 12, Ø	DDDØØØØØØØØØØØØF
	SRP	DWORD, 64-12, Ø	ØØØØØØØØØØØØDDDF
	AP	WDATE, =P' 1ØØ'	Indicate 1st month
CNVD#Ø15	DS	ØH	
	CP	DWORD, Ø(2, R1)	is the days < table
	BNH	CNVD#Ø2Ø	n:
	AP	WDATE, =P' 1ØØ'	Add 1 month
	SP	DWORD, Ø(2, R1)	Subtract date in month
	LA	R1, 2(, R1)	Bump to next table entry
	B	CNVD#Ø15	loop back
CNVD#Ø2Ø	DS	ØH	
	AP	WDATE, DWORD	Add date
	MVC	WORKDATE, EDMASKD	Move mask for date
	ED	WORKDATE, WDATE+3	Edit date
CNVD#999	PR		goback
*		Literals	
Z	DC	CL1' *'	zero or more char
O	DC	CL1' ?'	exactly 1 char
LEAPYR	DC	P' 31, 29, 31, 3Ø, 31, 3Ø, 31, 31, 31, 30, 31, 31, 30, 31, 31'	Leap table
NLEAPYR	DC	P' 31, 28, 31, 3Ø, 31, 3Ø, 31, 31, 31, 30, 31, 31, 30, 31, 31'	
EDMASKN	DC	X' 4Ø2Ø2Ø2Ø2Ø2Ø2Ø2Ø212Ø'	Edit mask for number
EDMASKT	DC	X' 4Ø212Ø2Ø7A2Ø2Ø7A2Ø2Ø'	Edit mask for time
EDMASKD	DC	X' 4Ø212Ø2Ø2Ø2Ø6Ø2Ø2Ø6Ø2Ø2Ø'	Edit mask for date
BLANK	DC	C' '	Blanks
	DS	ØF	DD table
DDINFO	DC	C' SYSIN ' , C' I' , C' O' , A(8Ø) , A(SMFP#Ø1Ø) , A(SMFP#Ø15)	
	DC	A(Ø)	
SYSIN	EQU	Ø	
DDINFOL	EQU	*-DDINFO	
	DC	C' SMFIN ' , C' I' , C' M' , A(6ØØØØ) , A(SMFP#Ø3Ø) , A(SMFP#Ø35)	
	DC	A(Ø)	
SMFIN	EQU	1	
	DC	C' STEPSU ' , C' O' , C' O' , A(133) , A(Ø) , A(Ø)	
	DC	A(SUSHDR)	
STEPSU	EQU	2	

```

        DC      C' JOBSU      ' , C' 0' , C' 0' , A(133) , A(Ø) , A(Ø)
        DC      A(SUJHDR)
JOBSU    EQU    3
        DC      C' STEPTIME' , C' 0' , C' 0' , A(133) , A(Ø) , A(Ø)
        DC      A(TMSHDR)
STEPTIME EQU    4
        DC      C' JOBTIME ' , C' 0' , C' 0' , A(133) , A(Ø) , A(Ø)
        DC      A(TMJHDR)
JOBTIME  EQU    5
        DC      C' STEPSTOR' , C' 0' , C' 0' , A(133) , A(Ø) , A(Ø)
        DC      A(STSHDR)
STEPSTOR EQU    6
DDINFOS  EQU    (*-DDINFO)/DDINFOL
        DC      X' FFFFFFFF'
MODDCBO  DCB    DCBE=MODDCBE, MACRF=(PM) , DSORG=PS, LRECL=133 output DCB
MODDCBOL EQU    *-MODDCBO
MODDCBI  DCB    DCBE=MODDCBE, DDNAME=Ø, MACRF=(GM) , DSORG=PS Input DCB
MODDCBIL EQU    *-MODDCBI
MODDCBE  DCBE   RMODE31=BUFF, EODAD=Ø, SYNAD=Ø Model DCBE
MODDCBEL EQU    *-MODDCBE
MODOPEN   OPEN   (,) , MF=L, MODE=31 Model OPEN
MODOPENL  EQU    *-MODOPEN
MODCLOS   CLOSE  (,) , MF=L, MODE=31 Model CLOSE
MODCLOSL EQU    *-MODCLOS
*
LTORG
SMF30PG CSECT
SUSHDR   DS     ØF          Header information
        DC      C' JobName  JobNum  Start Date St Time End D'
        DC      C' ate     End Time StepName Program   Total SU '
        DC      C'       CPU SU     SRB SU     IO SU
SUJHDR   DS     ØF
        DC      C' JobName  JobNum  Start Date St Time End D'
        DC      C' ate     End Time Total SU     CPU SU     SRB'
        DC      C' SU       IO SU WLM Cls  Serv Name
TMSHDR   DS     ØF
        DC      C' JobName  JobNum  Start Date St Time End D'
        DC      C' ate     End Time StepName Program   TCB CPU   SR'
        DC      C' B CPU     EXCP
TMJHDR   DS     ØF
        DC      C' JobName  JobNum  Start Date St Time End D'
        DC      C' ate     End Time TCB CPU   SRB CPU     EXCP '
        DC      C' WLM Cls  Serv Name
STSHDR   DS     ØF
        DC      C' JobName  JobNum  Start Date St Time End D'
        DC      C' ate     End Time StepName Program   USR<16(K)'
        DC      C'   USR>16(K) SYS<16(K)  SYS>16(K)
$DYN     DSECT   Dynamic area
SAVEAREA DS     18F          Save area

```

FILEDATA	DS	F	Pointer to file information
MASKLIST	DS	F	Pointer to mask list
RETCODE	DS	F	Return code
DWORD	DS	D	work
DEXCP	DS	D	work excp count
WDATE	DS	D	work date
WTIME	DS	D	work time
WORKTIME	DS	CL(L' EDMASKT)	work area
WORKDATE	DS	CL(L' EDMASKD)	work area
JOBWLMC	DS	CL08	wlm class
JOBWLMS	DS	CL08	wlm service name
STRTDTE	DS	CL10	Start date
ENDDATE	DS	CL10	End date
STRTTME	DS	CL08	Start Time
ENDTIME	DS	CL08	End time
JOBNAME	DS	CL08	Jobname
JOBNUM	DS	CL08	Job Number
STEPNAME	DS	CL08	Step Name
PROGRAM	DS	CL08	Program name
STEP	DS	X	Step indicator
ISSTEP	EQU	X' 00'	... Is step information
ISJOB	EQU	X' 01'	... Is job information
INCL	DS	X	Include indicator
INCLALL	EQU	X' 01'	... Include all
PRTLINE	DS	CL133	Print line
		ASAXWC MF=(L, MYLIST)	
WORKAREA	DS	CL256	WorkArea
\$DYNL	EQU	*-\$DYN	Length of Dyn area
FILEINFO	DSECT		DSECT for file area
FILEOPEN	DS	F	Open area addr
FILECLOS	DS	F	Close area addr
FILEDCB	DS	F	DCB addr
FILEDCBE	DS	F	DCBE addr
FILEREC	DS	F	rec addr
FILERECL	DS	F	rec length
FILEHEDR	DS	F	header addr
FILEIND	DS	X	indicator
	DS	CL3	filler
FILESIZ	EQU	*-FILEINFO	length of file area
HEDR	DSECT		Header information for rec
	DS	CL1	
HEDRNJME	DS	CL8	Jobname
	DS	CL1	
HEDRNJNUM	DS	CL8	Jobnumber
	DS	CL1	
HEDRSTDT	DS	CL10	Start date
	DS	CL1	
HEDRSTTM	DS	CL08	Start Time
	DS	CL1	

HEDRENDT	DS	CL10	End date
	DS	CL1	
HEDRENTM	DS	CL08	End time
	DS	CL1	
HEDRJOB	EQU	*-HEDR	Length of job header
HEDRSTNM	DS	CL08	Step Name
	DS	CL1	
HEDRPGNM	DS	CL08	Program Name
	DS	CL1	
HEDRSTEP	EQU	*-HEDR	Length of Step Header
HEDRL	EQU	*-HEDR	Total length of header
SUDSECT	DSECT		Dsect for SU details
SUTOT	DS	CL10	Total SU units
	DS	CL1	
SUCPU	DS	CL10	CPU SU units
	DS	CL1	
SUSRB	DS	CL10	SRB SU units
	DS	CL1	
SUI O	DS	CL10	I/O su units
	DS	CL1	
SUWLM	DS	CL08	WLM class name
	DS	CL1	
SUSRVNM	DS	CL08	WLM serv class
SUL	EQU	*-SUDSECT	Length
TIMEDSCT	DSECT		Dsect for TIME details
TIMECPU	DS	CL08	TCB CPU time
	DS	CL1	
TIMESRB	DS	CL08	TCP SRB time
	DS	CL1	
TIMEEXCP	DS	CL10	Total EXCP
	DS	CL1	
TIMEWLM	DS	CL08	WLM class name
	DS	CL1	
TIMESRVN	DS	CL08	WLM serv class
TIMEL	EQU	*-TIMEDSCT	Length
STORDSCT	DSECT		Dsect for Storage section
STORUB16	DS	CL10	User stor < 16M
	DS	CL1	
STORUA16	DS	CL10	User stor > 16M
	DS	CL1	
STORSB16	DS	CL10	Sys stor < 16M
	DS	CL1	
STORSA16	DS	CL10	Sys Stor > 16M
STORL	EQU	*-STORDSCT	Length
MASKDSCT	DSECT		Dsect for Mask details
MASKNEXT	DS	F	Ptr to Next mask info elem
MASKNAML	DS	F	Length of Mask
MASKNAME	DS	CL8	Name of Mask
MASKL	EQU	*-MASKDSCT	Length
DCBD	DSORG=PS, DEVD=DA		

I HADCBE
I FASMF 30
END SMF30PG

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Analysing data-in-virtual statistics

INTRODUCTION

Since its introduction a long time ago (with MVS/XA!), when it received some attention, DIV (Data-In-Virtual) seems to have fallen into oblivion. The main reason for that is the fact that DIV is somewhat difficult to use because the Assembly-language primitive functions one must use are not readily available in high-level languages. However, DIV, which is a set of primitive functions, enables an application program to load and manage substantial amounts of data into memory from a VSAM Linear DataSet (LDS). The LDS itself can grow to 4GB and the program can map up to (almost) 2GB of it at a time in central memory . Applications can create, read, and update data without the I/O buffer, blocksize, and record considerations that the traditional GET and PUT types of access method require. An application written for data-in-virtual views its permanent storage data as a seamless body of data without internal record boundaries. Among the applications that can be considered for a data-in-virtual implementation are applications that process large arrays, VSAM relative record applications, and BDAM fixed-length record applications. The potential benefits may be realized eventually as DIV merges with hiperspaces (a related concept) and as subsystems, languages, and application packages exploit the DIV benefits. For example, DIV is used by DFSMS when I/O to SMS control datasets is needed.

The data-in-virtual services process the application data in 4096-byte (4KB) units on 4KB boundaries called blocks. The application data resides in what is called a data-in-virtual object, a data object, or simply an object. The data-in-virtual object is a continuous string of uninterrupted data. When one writes an application using the techniques of data-in-virtual, the I/O takes place only for the data referenced and saved. Only the referenced pages of such an object are brought into virtual storage. Bytes of the mapped pages can be accessed and changed in normal program execution without regard to the need for updating. On request, or at the time the connection to the DIV object is terminated, only the changed pages are written back to the linear dataset. If one runs an application using conventional access methods, and then runs it again using data-in-virtual, one will notice a difference in performance. This difference depends on both the size of the dataset and its access pattern. To gain the right to view or update the object, an application must use the ACCESS service. ACCESS is similar to the OPEN macro of VSAM. It has a mode parameter of READ or UPDATE, and it gives your application the right to read or update the object. If the application has finished processing the object, it uses UNACCESS to relinquish access to the object.

Before using the DIV macro to process a linear dataset object (or a hiperspace object), one must create the dataset (or the hiperspace). *OS/390 MVS Programming:Authorized Assembler Services Guide* (GC28-1763) explains how to use DIV macro functions. The ‘how to’ reference for hiperspaces is the *Extended Addressability Guide*, (GC28-1468). Also, it is worth consulting *An Introduction to Data-in-Virtual* (GG66-0259), which may be a bit old but provides a few Assembler, Fortran, and PL/I examples.

COLLECTING DIV DATA

When enabled by the SMFPRMxx TYPE parameter, SMF creates record type 41, which provides resource usage information regarding data-in-virtual. There are two subtypes of this SMF record: subtype 1 is an ACCESS record – the ACCESS data

section is written when a DIV object is accessed; subtype 2 is an UNACCESS record – the counts for the I/O activity section are accumulated by data-in-virtual while the object is in use and are reported at the time of the UNACCESS request. The subtype 2 record is written whenever a data-in-virtual object is unaccessed.

A detailed description of layout of SMF type 41 record and its subtypes can be obtained from the *MVS System Management Facilities (SMF)* (SA22-7630-03) manual. You can also find the subtype descriptions in the macro ITVSMF41 in SYS1.MACLIB.

Based on record descriptions obtained from above mentioned manual a simple DIV report writer was written.

CODE

The code is a two part stream. In the first part (COPY412), selected SMF records (selection being defined by INCLUDEs condition) are copied from an SMF dataset to a VB file, which is the input file for the second step.

In the second step (DIV412), the captured records are formatted by invoking REXX EXEC (DIV412), and a report produced. The report shows the users performing access/unaccess of an object, along with the timestamp of when the object was accessed/unaccessed, the size of the object, and its read/write/re-read count data. Elapsed time during which the object was read or updated is calculated from two TOD timestamps.

JCL

```
//DIVJOB JOB ACCT#,
//           MSGLEVEL=(1, 1),
//           MSGCLASS=R,
//           NOTIFY=&SYSUID
//COPY412 EXEC PGM=ICETOOL
//TOOLMSG DD SYSOUT=*
//DFSMMSG DD SYSOUT=*
//RAWSMF DD DISP=SHR, DISP=HL q. SMFDUMPW
//SMF412 DD DSN=your.copied.by.sort.to.VB.smf.dataset,
//           SPACE=(CYL,(1)), UNIT=SYSDA,DISP=(NEW,PASS),
//           DCB=(RECFM=VB,LRECL=32756,BLKSIZE=32760)
```

```

//TOOLIN DD *
  COPY FROM(RAWSMF) TO(SMF412) USING(SMFI)
//SMFI CNTL DD *
  OPTION SPANIN=RC4, VLSHRT
  INCLUDE COND=(6, 1, BI , EQ, 41, AND, 23, 2, BI , EQ, 2)
/*
//DIV412 EXEC PGM=IKJEFT01, REGI ON=0M, DYNAMNBR=50, PARM=' %DIV412'
//SYSEXEC DD DISP=SHR, DSN=your. rexx. library
//SMF DD DISP=(SHR, PASS), DSN=your. copied. by. sort. to. VB. smf. dataset,
//DIV41 DD DSN=your. div. report. dataset,
//          SPACE=(CYL, (1, 1)), UNIT=SYSDA,
//          DISP=(NEW, KEEP), DCB=(RECFM=FB, LRECL=150)
//SYSPRINT DD SYSOUT=*
//SYSTSPRT DD SYSOUT=*
//SYSTSIN DD DUMMY
/*

```

DIV412 EXEC

```

/* REXX EXEC to read and format SMF record 41.2*/
ADDRESS TSO
/*-----*/
/* Print report header and labels */
/*-----*/
Out.1 = left(' ', 30, ' '),
         ||center('Data in Virtual Report ', 22, ),
         ||left(' ', 15, ' ')
Out.2 = left(' ', 20, ' '),
         ||center('Report produced on', 18, ),
         ||left(' ', 1, ' ')||left(date(), 11),
         ||left(' ', 1, ' ')||left('at ', 3, ' '),
         ||left(time(), 10)
Out.3 =
Out.4 = left('SMF date', 11) left('SMF time', 9),
        left('SID', 4)      left('DFP lvl.', 7),
        left('Job name', 9) left('DD name', 9),
        left('Access Time', 15) left('Unaccess Time', 15),
        left('A. Size', 6)   left('U. Size', 7),
        left('Mode', 5)     left('Elapsed. sec.', 10),
        left('Block', 5)    left('Block', 6),
        left('Block', 6)    left('Read', 5),
        left('Write', 5)
Out.5 = left(' ', 118)      left('read', 5),
        left('write', 6)    left('rread', 6),
        left('I/O', 5)      left('I/O', 3)
Out.6 = LEFT('-', 149, '-')
"EXECIO * DISKW DIV41 (STEM Out.)"
' EXECIO * DISKR SMF ( STEM x. FINIS'

```

```

    do i = 1 to x.0
/*-----*/
/*          Header/Self-defining Section           */
/*-----*/
smftype = c2d(SUBSTR(x.i, 2, 1))           /* SMF record type */
smfstype = c2d(SUBSTR(x.i, 19, 2))          /* Record subtype */
IF smftype = '41' Then
Do
  smfdate = SUBSTR(c2x(SUBSTR(x.i, 7, 4)), 3, 5) /* Unpack SMF date */
  smftime = smf(c2d(SUBSTR(x.i, 3, 4)))           /* Decode SMF time */
  sysid   = SUBSTR(x.i, 11, 4)                      /* System identification */
  trp     = c2d(SUBSTR(x.i, 21, 2))                /* number of triplets */
  opd     = c2d(SUBSTR(x.i, 25, 4))                /* offset to product section */
  lpd     = c2d(SUBSTR(x.i, 29, 2))                /* length of product section */
  npd     = c2d(SUBSTR(x.i, 31, 2))                /* number of product sections */
  od1    = c2d(SUBSTR(x.i, 33, 4))                /* offset to access section */
  ld1    = c2d(SUBSTR(x.i, 37, 2))                /* length of access section */
  nd1    = c2d(SUBSTR(x.i, 39, 2))                /* number of access sections */
  od2    = c2d(SUBSTR(x.i, 41, 4))                /* offset to unaccess section */
  ld2    = c2d(SUBSTR(x.i, 45, 2))                /* length of unaccess section */
  nd2    = c2d(SUBSTR(x.i, 47, 2))                /* number of unaccess sections */
  od3    = c2d(SUBSTR(x.i, 49, 4))                /* offset to i/o activity */
  ld3    = c2d(SUBSTR(x.i, 53, 2))                /* length of i/o activity */
  nd3    = c2d(SUBSTR(x.i, 55, 2))                /* number of i/o activity */
/*-----*/
/*          Product Section                     */
/*-----*/
IF opd <> 0 AND npd <> 0 Then do
  opd=opd-3
  dfplvl = SUBSTR(x.i, opd, 8)                  /* product level */
  prod   = SUBSTR(x.i, opd+8, 16)                /* component name */
end
/*-----*/
/*          Object ACCESS Data Section        */
/*-----*/
IF od1 <> 0 and nd1 <> 0 Then do
  od1=od1-3
  dda = SUBSTR(x.i, od1, 8)                      /* object ddname */
  aza = c2d(SUBSTR(x.i, od1+8, 4))              /* object size */
  ata= SUBSTR(ct(c2x(SUBSTR(x.i, od1+12, 4))), 11, 15) /* TOD */
  tya = c2d(SUBSTR(x.i, od1+16, 1))              /* object type */
  ama = c2d(SUBSTR(x.i, od1+17, 1))              /* access mode */
  jbn = SUBSTR(x.i, od1+18, 8)                  /* jobname/started task */
  SELECT
    when ama=1 then mode='Read'
    when ama=2 then mode='Update'
  END
end
/*-----*/
/*          Object UNACCESS Data Section      */
/*-----*/

```

```

IF od2 <> 0 and nd2 <> and then do
  od2=od2-3
  uzu = c2d(SUBSTR(x.i , od2, 4))           /* object size*/
  utu = SUBSTR(ct(c2x(SUBSTR(x.i , od2+4, 4))), 11, 15) /* TOD*/
end
/*-----*/
/*          Object I/O Activity Section */
/*-----*/
IF od3 <> 0 and nd3 <> 0 then do
  od3=od3-3
  brd = c2d(SUBSTR(x.i , od3, 4))           /* tot.no. of reads*/
  bwr = c2d(SUBSTR(x.i , od3+4, 4))           /* tot.no. of writes*/
  brr = c2d(SUBSTR(x.i , od3+8, 4))           /* tot.no. of re-reads*/
  inc = c2d(SUBSTR(x.i , od3+12, 4))          /*tot.no.of i/o for read*/
  ouc = c2d(SUBSTR(x.i , od3+16, 4))          /*tot.no.of i/o for write*/
end
timedi f=diff(utu,ata)      /* how long the object was accessed*/
/*-----*/
/* formatting and printing a DIV entry */
/*-----*/
divout = left(Date('N', smfdate, 'J'), 11) left(smftime, 9),
  left(sysid, 4) left(dfplvl, 8) left(jbn, 8),
  left(dda, 8) right(ata, 15) left(utu, 17),
  left(aza, 6) left(uzu, 6) left(mode, 7),
  right(timedi f, 8),
  right(brd, 4) right(bwr, 5) right(brr, 5),
  right(inc, 5) right(ouc, 5)
PUSH divout
  "EXECIO 1 DISKW DIV41"
end
end
exit

```

SMF PROCEDURE

```

/* REXX - convert a SMF time */
arg time
  time1 = time % 100
  hh = time1 % 3600
  hh = RIGHT("0"||hh, 2)
  mm = (time1 % 60) - (hh * 60)
  mm = RIGHT("0"||mm, 2)
  ss = time1 - (hh * 3600) - (mm * 60)
  ss = RIGHT("0"||ss, 2)
  otme = hh||":"||mm||":"||ss           /* Compose SMF time*/
return otme

```

CT PROCEDURE

```
/*
 * TOD timestamp is a 16-byte EBCDIC representat */
/* The BLSUXTOD proc is described in "z/OS */
/* V1R3 MVS IPCS Customization" */
*/
arg todtime
If todtime    <> '0000000000000000' Then
Do
  TOD_Value = X2C(todtime)
  Returned_Date = '-----'
  address LINKPGM "BLSUXTOD TOD_Value Returned_Date"
End
Else
  Returned_Date = ''
Return Returned_Date
```

DIF PROCEDURE

```
/*
 * Dif: REXX subroutine to find the
 * difference between two timestamps, in this
 * case in seconds
 */
arg time2,time1
parse var time2 h2 ':' m2 ':' s2 ':' t2
parse var time1 h1 ':' m1 ':' s1 ':' t1
tot2 = h2*3600 + m2*60 + s2
tot1 = h1*3600 + m1*60 + s1
es = tot2 -tot1
if es<0 then es=es+86400
eh=es%3600
es=es//3600
ex=es
em=es%60
es=es//60
/* et=right(eh,2,0)' : 'right(em,2,0)' : 'right(es,2,0) */
return ex
```

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Waiting for datasets

We have a group of FTP datasets ‘pushed’ to us – the first part of the dataset name is known, but the trailing qualifier is variable. The following program is started by our automation package, and is the name of a volume to look on and some further variable (time-based) information that is then used for a ‘cvaffilt’ call, which passes back the names (if any) of datasets found that match our pattern. We then extract the dataset names and write them to a dataset that is passed to the next step (obviously this could be replaced by a WTO/WTOS, a command being issued, or any other relevant action). The program is written to look for three datasets, and will wait for five minutes for them to arrive.

WAIT4FTP PROGRAM

```
TITLE 'WAIT4FTP - WAIT FOR INCOMING FTP DATASETS'
*****
* WAIT4FTP: SCAN A VOLUME, USING A DSN PATTERN, TO SEE HOW MANY
*           DATASETS ARE PRESENT. WAIT FOR UP TO 5 MINUTES TO SEE
*           IF THREE HAVE ARRIVED. IF NOT, EXIT WITH RC.
*
*           FOR FTP TRANSFER OF DATASETS.
*
* PARMs:   TWO PARMs, BOTH MANDATORY, SEPARATED BY A COMMA -
*           1/ A VOLID (TO LOOK ON FOR THE DATASETS)
*           2/ 4 DIGITS, TO BE INSERTED INTO THE DATASET MASK. THIS
*               MASK WILL MAKE UP A DATASET NAME, IN THE FORM:
*
*           FTPXFER. DATA. FILE. PFTP. DNNNNNTY. **
*
*           WHERE 'NNNN' ARE THE 4 DIGITS PASSED TO US.
*
* OUTPUT:  DATASET CONTAINING THE NAMES OF THE DATASETS LOCATED,
*           WRITTEN TO DDNAME 'DSNAMES'.
*
*****
PRINT NOGEN
*****
* HOUSEKEEPING...
*****
WAIT4FTP CSECT
```

```

WAI T4FTP AMODE 31
WAI T4FTP RMODE 24
    BAKR  R14, Ø           SAVE CALLER DATA ON STACK
    LR    R12, R15         GET ENTRY POINT
    LA    R11, 2048(R12)   LOAD SECOND BASE
    LA    R11, 2048(R11)
USING WAI T4FTP, R12, R11   ADDRESSABILITY
    L    R9, Ø(R1)          GET A(PARMS)
    LTR   R9, R9            ANY?
    BZ    BADPARMS         NO...
    CLC   Ø(2, R9), =H' 11' 11-BYTE PARMS PASSED?
    BNE   BADPARMS         NO...
    CLI   8(R9), C', '     COMMA IN BETWEEN?
    BNE   BADPARMS         NO...
    MVC   VOLID(6), 2(R9)  SAVE VOLID...
    MVC   DSNSER(4), 9(R9) ... AND SERIAL
    MVC   WT01+23(6), VOLID MOVE VOLID TO WTO
    MVC   WT01+35(L' DSNAME), DSNAME MOVE MASK TO WTO
WT01    WTO   ' WAIT4FTP- SCAN XXXXXX FOR: .....X
                ..... ', ROUTCDE=11
*****
* INITIALIZE BUFFER LIST HEADER (BFLH) AND ELEMENTS (BFLE)...
*****
INITBFL DS  ØH
*
    LA    R2, BFLHDEF      ADDRESS WORKAREA
    LA    R3, BFLSIZE       LOAD LENGTH TO CLEAR
    XR   R5, R5            Ø PADDING BYTE
    MVCL R2, R4            SET BUFFER LIST AREA TO Ø'S
*
    LA    R1, BFLHDEF      TEMP ADDRESSABILITY TO BFLH
    USING BFLMAP, R1
    MVI   BFLHNOE, BUFFNUM SET NUMBER OF BUFFER ELEMENTS
    OI    BFLHFL, BFLHDSCB ID AS DSCB BUFFER ELEMENT LIST
    LA    R2, BFLHDEF+BFLHLEN R2 -> 1ST BUFFER LIST ELEMENT
    USING BFLE, R2          TEMP ADDRESSABILITY
    LA    R3, DSCBDEF      R3 -> 1ST DSCB BUFFER
    LA    R4, BUFFNUM       R4 = NO. OF ELEMENTS & BUFFERS
BFLEINIT DS  ØH
*
    OI    BFLEFL, BFLECHR REQUEST CCHHR ON RETURN
    MVI   BFLELTH, DSCBSIZE SET BUFF LEN TO FULL DSCB SIZE
    ST    R3, BFLEBUF       SET A(DSCB BUFFER)
    LA    R2, BFLELN(R2)   R2 -> NEXT BUFFER LIST ELEMENT
    LA    R3, DSCBSIZE(R3) R3 -> NEXT DSCB BUFFER
    BCT   R4, BFLEINIT     LOOP THROUGH ALL ELEMENTS
    DROP  R1, R2            DROP TEMP ADDRESSABILITY
*****
* INITIALIZE THE FILTER CRITERIA LIST (FCL) HEADER AND ELEMENT...
*****
XC    FCLDEF(FCLSIZE), FCLDEF   SET FCL AREA TO Ø

```

LA	R1, FCLDEF	R1 -> FCL HEADER
USING	FCLMAP, R1	TEMP ADDRESSABILITY TO FCL
MVC	FCLID, =C' FCL '	SET EYECATCHER 'FCL'
MVC	FCLCOUNT, =H' 1'	SET NUMBER OF FCL ELEMENTS = 1
OI	FCL1FLAG, FCL1EQF1	RETURN ONLY FORMAT1 DSCBS
LA	R2, FCLHDEND	R2 -> 1ST (ONLY) FCL ELEMENT
USING	FCLDSN, R2	TEMP ADDRESSABILITY
MVC	FCLDSNLG, DSNAMELN	SET LENGTH OF DSN PATTERN
LA	R3, DSNAME	GET DSN PATTERN ADDRESS
ST	R3, FC LDSNA	... SAVE IN FCL
DROP	R1, R2	

* SCAN THROUGH THE UCBS, LOOKING FOR OUR VOLUME.		*

SCANVOLS	DS ØH	
	USING UCB0B, R4	ADDRESSABILITY TO UCB
	LA R4, UCBAREA	LOCATE UCB WORKAREA
	XC UCBWORK, UCBWORK	+INITIALIZE UCBSCAN WORKAREA
UCBLOOP	DS ØH	
*		
	UCBSCAN COPY,	X
	WORKAREA=UCBWORK,	X
	UCBAREA=UCBAREA,	X
	DCEAREA=None,	X
	DCELEN=Ø,	X
	VOLSER=VOLID,	X
	DEVN=Ø,	X
	DYNAMIC=YES,	X
	RANGE=ALL,	X
	NONBASE=NO,	X
	DEVCLASS=DASD,	X
	DEVCID=Ø,	X
	IOCTOKEN=None,	X
	LINKAGE=SYSTEM,	X
	PLISTVER=MAX	X
*		
	LTR R15, R15	GOT UCB OK?
	BZ GETREAL	YES.. GET THE REAL UCB ADDRESS
	C R15, =F' 4'	END OF UCBS?
	BE NOTFND	YES.. SAY WE COULDN'T FIND VOL
	B BADCALL	NO... ERROR

* NOW WE HAVE OUR VOLUME - GET THE *REAL* UCB ADDRESS BY CALLING THE *		*
* "UCBLOOK" MACRO... .		*

GETREAL	DS ØH	
*		
	MODESET KEY=ZERO, MODE=SUP	SUPERVISOR STATE FOR "UCBLOOK"
*		
UCBLOOK	UCBLOOK VOLSER=UCBVOLI ,	USE VOLID FROM UCB COPY X

	UCBPTR=UCBADDR, LOC=ANY, NOPIN, RANGE=ALL, DEVCLASS=DASD, DYNAMIC=YES	SAVE UCB ADDRESS IN HERE UCB CAN BE ANYWHERE DON'T PIN UCB 3&4 DIGIT UCBS MAKE SURE ITS A DISK! CHECK DYNAMIC UCBS	X X X X X
*	ST R15, SAVER15	SAVE RC	
*	MODESET KEY=NZERO, MODE=PROB	BACK TO NORMAL	
*	L R15, SAVER15 LTR R15, R15 BNZ LOOKERR	RELOAD RC OK? NO... PANIC...	

*	NOW WE HAVE THE VOLUME WE WANT - ISSUE A CVAFFILT 'ACCESS=READ' *		
*	REQUEST TO SEE IF THERE ARE ANY RELEVANT FTP DATASETS, AND, IF SO, *		
*	HOW MANY... *		

CVPLLOOP	LA R7, 5 DS ØH	WAIT UP TO 5 TIMES	
*			
CVPL	CVAFFILT ACCESS=READ, FCL=FCLDEF, BUFLIST=BFLHDEF, UCB=UCBADDR		X X X
*			
WT02	C R15, F4 BE MOREBUFF LTR R15, R15 BNZ CVAFERR LA R3, FCLDEF USING FCLMAP, R3 LH R2, FCLDSCBR CVD R2, DWORD UNPK UNPKFLD(5), DWORD+5(3) OI UNPKFLD+4, X' FØ' MVC WTO2+49(1), UNPKFLD+4	WAS RC=4? YES.. NEED MORE BUFFERS RC OTHER THAN Ø? YES... DODGY R3 -> FCL HEADER TEMP ADDRESSABILITY TO FCL GET DSCB COUNT MAKE IT DECIMAL UNPACK IT SET CORRECT SIGN	
WT03	WTO 'WAIT4FTP- NUMBER OF FTP DATASETS LOCATED=?', ROUTCDE=11 CLC FCLDSCBR, =H' 3' BE GOTEM	GOT ALL 3? YES.. CARRY ON	
WT04	WTO 'WAIT4FTP- NOT ALL DATASETS HAVE ARRIVED - WAITING...', X ROUTCDE=11 STIMER WAIT, BINTVL=BIN6Ø BCT R7, CVPLLOOP	ROUTCDE=11 HAVE ANOTHER GO	
GOTEM	MVC RETC, F16 B RETURN DS ØH	SET RC=16	

```

OPEN   (DSNAMES, OUTPUT)      OPEN OUTPUT DCB
LA    R3, DSCBDEF          POINT TO FIRST DSCB
DSNLOOP DS  ØH
MVC   OUTREC(44), Ø(R3)     MOVE DSNAME TO OUTPUT RECORD
PUT   DSNAMES, OUTREC       WRITE IT OUT
LA    R3, DSCBSIZE(R3)      BUMP TO NEXT ONE
BCT   R2, DSNLOOP          GET THE REST
CLOSE  DSNAMES
RETURN DS  ØH
L     R15, RETC            LOAD RETURN CODE
PR   ,                      RESTORE CALLER DATA, RETURN
*****
* DIDN' T FIND OUR VOLUME... *
*****
NOTFND DS  ØH
MVC   WT05+26(6), VOLID      MOVE VOLID TO WTO
WT05   WTO  'WAIT4FTP- VOLUME "XXXXXX" NOT LOCATED...', ROUTCDE=11
MVC   RETC, F8
B    RETURN
*****
* INSUFFICIENT BUFFERS TO HOLD DSCBS... *
*****
MOREBUFF DS  ØH
WTO  'WAIT4FTP- TOO MANY DSCBS LOCATED FOR THE # OF BUFFERS DX
DEFINED...', ROUTCDE=11
MVC   RETC, F12
B    RETURN
*****
* BAD RETURN CODE FROM 'UCBSCAN'... *
*****
BADCALL DS  ØH
LR   R10, R15                SAVE RETC
WTO  'WAIT4FTP- BAD CALL TO "UCBSCAN"...', ROUTCDE=11
LR   R15, R10                RELOAD RETC
DS   F                        BANG
*****
* NO/INVALID PARMS PASSED... *
*****
BADPARMS DS  ØH
WTO  'WAIT4FTP- NO/INVALID PARMS PASSED...', ROUTCDE=11
ABEND 111, DUMP
*****
* BAD RETURN CODE FROM 'UCBLOOK'... *
*****
LOOKERR DS  ØH
LR   R10, R15                SAVE RETC
ST   RØ, SAVERØ              SAVE REASON
WTO  'WAIT4FTP- BAD CALL TO "UCBLOOK"...', ROUTCDE=11
L    RØ, SAVERØ              RELOAD REASON
LR   R15, R10                RELOAD RETC

```

ABEND 123

BANG

```
*****  
* ERROR IN 'CVAFFILT' MACRO... *  
*****  
CVAFERR DS ØH  
LA R1, CVPL+4  
USING CVPLMAP, R1  
STH R15, HWORD  
BAL R9, CONVR15  
MVC CVAFR15(2), UNPKFLD  
MVI HWORD, X' ØØ'  
MVC HWORD+1(1), CVSTAT  
BAL R9, CONVR15  
MVC CVAFSTAT(2), UNPKFLD  
MVC WT06+18(58), CVAFMSG  
WT06 WTO 'WAIT4FTP- XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX', ROUTCDE=11  
ABEND 555 BANG...  
DROP R1  
*****
```

```
* SUBROUTINE: CONVERT RETURN CODES, ETC INTO PRINTABLE HEX... *  
*****  
CONVR15 DS ØH  
UNPK UNPKFLD(3), HWORD+1(2) UNPACK RETC + 1 BYTE  
TR UNPKFLD(2), TRTAB-24Ø XLATE TO PRINTABLE HEX  
BR R9 RETURN FROM ROUTINE  
EJECT  
*-----*
```

```
*-----*  
* LTORG LITERAL POOL  
*  
DSNAME DC C' FTPXFER. DATA. FILE. PFTP. DNNNNNTY. **'  
DSNSER EQU DSNAME+24, 4  
DSNAMELN DC AL1(*-DSNAME)  
VOLID DC CL6' ??????  
UNPKFLD DS XL5  
TRTAB DC C' Ø123456789ABCDEF'  
HWORD DC H' Ø'  
DWORD DC D' Ø'  
F4 DC F' 4'  
F8 DC F' 8'  
F12 DC F' 12'  
F16 DC F' 16'  
RETC DC F' Ø'  
UCBADDR DC F' Ø'  
SAVER15 DC F' Ø'  
SAVERØ DC F' Ø'  
SWITCH DC C' N'  
BIN6Ø DC F' 1000' 6Ø SEC INTERVAL  
CVAFMSG DS ØCL58
```

```

        DC      CL35' ERROR IN "CVAFFILT" MACRO - RC = X"
CVAFR15  DC      CL2' '
        DC      CL13' " CVSTAT = X"
CVAFSTAT DC      CL2' '
        DC      CL6' ''
*-----*
* UCB SCAN ROUTINE PARAMETERS... *
*-----*
        DS      0F
UCBAREA  DS      XL48          HOLDS UCB COMMON & DEV SEGS
UCBWORK   DS      XL100         UCBSCAN WORKAREA
*-----*
* DCB FOR OUTPUT... *
*-----*
DSNAMES   DCB    DDNAME=DSNAMES,           X
            LRECL=80,           X
            RECFM=FB,           X
            DSORG=PS,           X
            MACRF=PM
OUTREC    DC     CL80' '
*-----*
* SPACE ALLOCATION FOR CVPL, FCL, BFL, AND DSCB BUFFERS... *
*-----*
FCLDEF    DS     (FCLHDLLEN+FCLDSNEL)X  FCL HEADER AND 1 FCL ELEMENT
FCLSIZE   EQU    *-FCLDEF
*=====DEFINE A CVAF BUFFER LIST WITH 'N' BUFFER LIST ELEMENTS=====
BFLHDEF   DS     (BFLHLEN)X      BUFFER LIST HEADER (BFLH)
BFLEDEF   DS     (BUFFNUM*BFLELN)X  'N' BUFFER LIST ELEMENTS (BFLE'S)
BFLSIZE   EQU    *-BFLHDEF
*=====DEFINE 'N' FULL DSCB BUFFERS=====
DSCBDEF   DS     (BUFFNUM*DSCBSIZE)X
*-----*
* REGISTERS EQUATES, ETC... *
*-----*
BUFFNUM   EQU    3             3 BUFFER LIST ELEMENTS AND BUFFERS
*
YREGS
*
CVPLMAP   I CVAFPL CVPLFSA=YES
FCLMAP    I CVFCL
BFLMAP    I CVAFBFL
        I EFUCBOB ,
        PRINT NOGEN
        CVT   DSECT=YES
DSCBMAP   DSECT
        I ECSDL1 (1)           FORMAT1 DSCB MAPPING TO BEF BUFFSIZE
DSCBSIZE EQU    *-I ECSDL1
*
END           , END OF PROGRAM

```

SAMPLE JCL TO RUN WAIT4FTP

```
//JOBCARD  
//S1      EXEC PGM=WAIT4FTP, PARM='FTP001, 1312' <-- VOLID, TIMESTAMP  
//STEPLIB  DD DISP=SHR, DSN=TEST. LINLIB  
//SYSPRINT DD SYSOUT=*  
//DSNAMES  DD DSN=&&DSNAMES, DISP=(MOD, PASS), UNIT=SYSDA, SPACE=(TRK, 1)
```

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Data conversion

This utility takes a COBOL copybook and converts it into a REXX include file. If you compile REXX programs, the REXX compiler parses the program looking for special INCLUDE instructions and then copies that code into the source before finishing the compile – much like a pre-processor.

The utility was written to perform data conversion. A project leader in the group decided that we REXX users had to use the same names for our variables as the COBOL people. The COBOL people were getting their copybooks from a data-mapping group, but that left us REXXers out in the cold. Initially, some of us started to manually convert COBOL copybooks to a REXX-like format, but with frequent changes in the data-mapping. Well, you can see that it had problems.

So I wrote the attached utility. With it we did the majority of the conversion using REXX rather than COBOL (much to the chagrin of the aforementioned project leader), and we came in a full five months ahead of schedule.

I have recently found another use for it, and once again it's a lifesaver.

COPYBOOK REXX

```
/* ----- REXX -----  
- Title: Convert COBOL Copy Books To REXX Include Format -  
- ======  
- Program Function:  
-  
- Converts a COBOL copybook to something REXX can work with.  
-  
- -----  
- Feel free to modify this program as you see fit.  
- -----  
-  
- Syntax:  
-  
- COPYBOOK MEMBER SOURCEPDS DESTPDS  
-  
- Where:  
- COPYBOOK -> This program.  
- MEMBER -> Name of COBOL copybook member to be  
- converted to REXX.  
- SOURCEPDS -> PDS containing COBOL copybook members.  
- DESTPDS -> PDS that will contain the converted  
- copybook in REXX format.  
-  
-----*/  
ConvertCobolCopybooks:  
arg member source dest .  
call Initialize  
call PerformCopyBookConversion  
call Finalize  
return  
/* -----*/  
Initialize:  
k = 0; offset = 1; true = 1; false = 0 ; fieldlen = 0  
level.0 = ''; variable.0 = ''; offset.0 = ''; stackptr = 0  
/* Initialize stacks */  
levelq.0 = ''; variableq.0 = ''; offsetq.0 = ''; queuetop = 0  
/* Initializing queues */  
lastlevel = 0  
/* Just what it says */  
i = 0  
return  
/* -----*/  
PerformCopyBookConversion:  
call ReadCopybook  
do while (i < copybook.0)  
i = i + 1  
copybook = PreProcessCopybook(copybook.i)
```

```

parse var copybook level variable pic field remainder
if ((level < lastlevel) | (pic = '' & level = Level StackTop())))
then
    do while (level <= Level StackTop())
        call PopStacks
    end
    lastlevel = level
select
    when (left(copybook, 1) = "*") then
        nop
    when (pic = 'REDEFINES') then
        do
            offset = RedefinesProcessing(field)
            parse var copybook level variable redword redvar
            pic field remainder
                if (field ~= '') then
                    call ProcessCobol Level
                else
                    call PushStacks
                end
            when level = 88 then
                call Level 88Processing
            when pic = '' then
                call PushStacks
            otherwise
                call ProcessCobol Level
            end
        end
    end
    do while (stackptr ~= 0)
        call PopStacks
    end
return
/* -----*/
PreProcessCopybook:
    arg copybook
    if (left(copybook, 1) ~= " ") | (left(copybook, 1) ~= "*") then
        parse var copybook 1 junk 7 copybook 73 .
    copybook = strip(copybook)
select
    when (copybook = '') then
        copybook = '*'
    when (left(copybook, 1) ~= '*') then
        copybook = BuildCopybookLine(copybook)
    otherwise nop
end
return copybook;
/* -----*/
BuildCopybookLine:
    arg line

```

```

    line = strip(line)
    do while ((substr(line,length(line),1) != '.') & (left(line,1)
    != "") & (i < copybook.0))
        i = i+1; line =line' 'strip(copybook.i)
    end
    line = substr(line,1,length(line)-1) /* gets rid of period */
return line
/* -----*/
Level StackTop:
return level.stackptr
/* -----*/
Variable StackTop:
return variable.stackptr
/* -----*/
PushStacks:
    stackptr = stackptr + 1
    variable.stackptr = variable
    level.stackptr = level
    offset.stackptr = offset
return
/* -----*/
PopStacks:
    newoffset = offset.stackptr
    newlevel = level.stackptr
    newvariable = translate(variable.stackptr,'_','-')
    stackptr = stackptr - 1
    k = k + 1
    outrec.k = " parse var record" left(newoffset,6)
left(newvariable, 40) offset
    x = Enqueue(newlevel,newvariable,newoffset)
return
/* -----*/
Level 88Processing:
    variable = translate(variable,'_','-')
    field = TranslateField(field)
    k = k + 1
    outrec.k = " /* 88-Level */ ||variable|| =
(||lastVariable||"|"||field||");
    if (strip(remainder) != "") then
        do
            string = (pos(" ",remainder) != 0);
            if (string) then
                do while (remainder != "")
                    parse var remainder "" "val" "" remainder
                    outrec.k = outrec.k|| |
(||lastVariable||"|"||val||");
                end
            else
                do while (remainder != "")

```

```

        parse var remainder val remainder
        outrec.k = outrec.k || " |
("||lastVariable| "="||val||")";
    end
end
return
/* -----
TranslateField: procedure expose fieldlen
arg field
select
    when (field = 'SPACES') | (field = 'SPACE') then
        field = "" "copies(" ",fieldlen)"" "
    when (field = 'ZERO') | (field = 'ZEROS') then
        field = Ø
    when (field = 'LOW-VALUES') then
        field = copies('00'x,fieldlen)
    when (field = 'HIGH-VALUES') then
        field = copies('FF'x,fieldlen)
    otherwise nop
end
return field
/* -----
ProcessCobol Level :
value =
computational = (pos("COMP-3", remainder) -= Ø)
binary = (pos("COMP ", remainder) -= Ø)
if pos("VALUE", remainder) -= Ø then
    parse var remainder "VALUE" value
fieldlen = DetermineFieldLength(field)
newoffset = fieldlen + offset
k = k + 1
variable = translate(variable,'_','')
lastvariable = variable /* storage bucket in case of 88 levels*/
outrec.k = " parse var record" left(offset,6) left(variable,40)
newoffset
x = Enqueue(level,variable,offset)
offset = newoffset
if value -= '' then
    do
        k = k + 1
        outrec.k = variable" = "translateField(value);
    end
return
/* -----
RedefinesProcessing: procedure expose variableno. offsetq. queueptr
arg redefinesvariable
queueptr = 1
redefinesvariable = translate(redefinesvariable,'_','')
do while redefinesvariable -= variableno.queueptr & queueptr <

```

```

queue top
    queueptr = queueptr + 1
end
return offsetq.queueptr
/* -----*/
Enqueue: procedure expose levelq.variableq.offsetq.queue top
    parse arg level, variable, offset
    queue top = queue top + 1
    levelq.queue top = level
    variableq.queue top = variable
    offsetq.queue top = offset
return rc
/* -----*/
DetermineFieldLength: procedure expose computational binary
    arg picmap
    parse var picmap decimal 'V' fraction
    fieldlen = 0
    if decimal ~= '' then
        do
            parse var decimal pre('len')
            if len = '' then
                do
                    if substr(pre, 1, 1) = 'S' then
                        len = length(pre) - 1
                    else
                        len = length(pre)
                end
            end
            fieldlen = fieldlen + len
        end
    if fraction ~= '' then
        do
            parse var fraction pre('len')
            if len = '' then
                len = length(pre)
            fieldlen = fieldlen + len
        end
    end
    select
        when computational then
            fieldlen = trunc((fieldlen + 1)/2)
        when binary then
            select
                when fieldlen <= 4 then fieldlen = 2
                when fieldlen <= 9 then fieldlen = 4
                when fieldlen <= 16 then fieldlen = 6
                otherwise fieldlen = 8
            end
        otherwise nop
    end
return fieldlen
/* -----*/

```

```

ReadCopybook:
    source = source' (' member')';
    if sysdsn(" "source""") != "OK" then
        exit 8
    else
        x = ReadDataset(source, "COPYBOOK")
    return
/* -----*/
ReadDataset:
    parse arg file,stemvar .
    address tso "alloc dd(INDD) da(' "file""') shr"
    "execio * diskr INDD (stem "stemvar". finis"
    address tso "free dd(INDD)"
return rc
/* -----*/
Finalize:
    convertedcopybook = dest"("member")"
    if sysdsn(" "dest""") != 'OK' then
        do
            address tso "alloc dd(TEMPDD) da(' "dest""') new reuse tr dir(20)",
                "sp(30 30) lrecl(255) recfm(V B) dsorg(P0) blksze(0)"
            address tso "free dd(TEMPDD)"
        end
    address tso "alloc dd(OUTDD) da(' "convertedcopybook""') shr"
    "execio "k" diskw OUTDD (stem outrec. finis"
    address tso "free dd(OUTDD)"
return
/* =====*/

```

SAMPLE COBOL COPYBOOK

This is an example of a COBOL copybook:

```

*****
* CABT0025 - CUI JURISDICTION TABLE RECORD *
*          GTE HISTORY SECTION                  TCID *
*  DATE    INITIAL   RVL    DESCRIPTION      DPSR *
*  07/01/87 KAN      R1V01L01  CABS ENHANCEMENT *
*  12/28/88 SAC      R2V01L01  DATE SENSITIVITY CABD0070 *
*  12/28/88 SAC      R2V01L01  FGB MPB INDICATOR CABD0119 *
*  06/05/90 REW      R2V01L01  AT&T Ø EMR1101 CABD0273 *
*  04/28/91 SAC      R2V06L01  DB800 INDICATORS CABD0306 *
*  04/28/91 SAC      R2V06L01  CELL NATL THRESHLD CABD0307 *
*  03/12/93 LB       R1V24L01  ADD VALUE OF 'L'  CBSD0838 *
*                           TO TRAFFIC-TYPE- *
*                           BILLABLE-INDICATOR *
*****
03  T0025-TABLE-IDENT          PIC X(05).

```

```

88 T0025-TABLE-ID           VALUE 'T0025' .
03 T0025-TABLE-DATA.
05 T0025-KEY.
 10 T0025-PARTIAL-KEY.
    15 T0025-FROM-NPA-NXX.
      20 T0025-FROM-NPA          PIC X(03).
      20 T0025-FROM-NXX         PIC X(03).
    15 T0025-TO-NPA            PIC X(03).
 10 T0025-EFFECTIVE-START   PIC 9(6).
 10 T0025-EFFECTIVE-END    PIC 9(6).

05 T0025-FUNCTION.
 10 T0025-TABLE-LEVEL        PIC X(1).
 10 T0025-JURISDICTION     PIC 9.
 10 T0025-STATE-CODE        PIC X(02).
 10 T0025-BILLING-LOCATION-CODE. 00002400
    15 T0025-OPERATING-GROUP  PIC X(01).
    15 T0025-CONSOL-COMPANY  PIC X(01).
    15 T0025-COMPANY          PIC X(01).
    15 T0025-AREA-DIV         PIC X(01).
    15 T0025-DIV-DIST        PIC X(02).
    15 T0025-PLANT-CODE      PIC X(04).
 10 T0025-BAN-STATE-CODE    PIC X(01).
 10 T0025-ORIG-MMU          PIC X(01).
 10 T0025-ORIG-RATE-CNTR   PIC 9(03).
 10 T0025-OLATA-CODE        PIC 9(03).
 10 T0025-BILLING-RAO       PIC X(03).
 10 T0025-ASSOC-BELL-RAO   PIC X(03).
 10 T0025-RCC-THRESHOLD-SEC PIC 9(02).
 10 T0025-IND-EQUAL-ACCESS  PIC X(01).
    88 T0025-EQUAL-ACCESS     VALUE 'Y'.
    88 T0025-NON-EQUAL-ACCESS VALUE 'N'.
 10 T0025-FGB-MPB-IND       PIC X(01).
 10 T0025-FGCD-MPB-IND     PIC X(01).
 10 T0025-CONV-OPH-INTRA   PIC S9V9(02) COMP-3.
 10 T0025-CONV-OPH-INTER   PIC S9V9(02) COMP-3.
 10 T0025-CONV-DIAL        PIC S9V9(02) COMP-3.
 10 T0025-RECORD-POINT-DDD  PIC S9(06)  COMP-3.
 10 T0025-RECORD-POINT-OPH  PIC S9(06)  COMP-3.
 10 T0025-RECORD-POINT-TERM PIC S9(06)  COMP-3.

10 T0025-ATT-SOURCE-OF-DATA.
 15 T0025-ATT-OPH-ZERO-PLUS  PIC X(01).
    88 T0025-ATT-TYPE-1-TOLL   VALUE 'T' 'L' 'C'
    88 T0025-ATT-TYPE-1-UMS   VALUE 'U'.
    88 T0025-ATT-TYPE-1-EMR   VALUE 'D'.
 15 T0025-ATT-OPH-ZERO-MINUS PIC X(01).
    88 T0025-ATT-TYPE-2-TOLL   VALUE 'T' 'L' 'C'
    88 T0025-ATT-TYPE-2-UMS   VALUE 'U'.
    88 T0025-ATT-TYPE-2-EMR   VALUE 'D'.
 15 T0025-ATT-MTS-ORIG      PIC X(01).
    88 T0025-ATT-TYPE-3-TOLL   VALUE 'T' 'L'.

```

```

88 T0025-ATT-TYPE-3-UMS      VALUE 'U'.
15 T0025-ATT-700-ORIG        PIC X(01).
88 T0025-ATT-TYPE-4-TOLL    VALUE 'T' 'L'.
88 T0025-ATT-TYPE-4-UMS    VALUE 'U'.
15 T0025-ATT-800-ORIG        PIC X(01).
88 T0025-ATT-TYPE-5-TOLL    VALUE 'T' 'L'.
88 T0025-ATT-TYPE-5-UMS    VALUE 'U'.
15 T0025-ATT-DB800-ORIG     PIC X(01).
88 T0025-ATT-DB800-TOLL    VALUE 'T' 'L'.
88 T0025-ATT-DB800-UMS    VALUE 'U'.
88 T0025-ATT-DB800-EMR    VALUE 'E'.
15 T0025-ATT-900-ORIG        PIC X(01).
88 T0025-ATT-TYPE-6-TOLL    VALUE 'T' 'L'.
88 T0025-ATT-TYPE-6-UMS    VALUE 'U'.
15 T0025-ATT-MTS-TERM       PIC X(01).
88 T0025-ATT-TYPE-7-UMS    VALUE 'U'.

10 T0025-NONATT-SOURCE-OF-DATA.
15 T0025-NONATT-OPH-ZERO-PLUS  PIC X(01).
88 T0025-NONATT-TYPE-1-TOLL   VALUE 'T' 'L'.
88 T0025-NONATT-TYPE-1-UMS   VALUE 'U'.
15 T0025-NONATT-OPH-ZERO-MINUS PIC X(01).
88 T0025-NONATT-TYPE-2-TOLL   VALUE 'T' 'L'.
88 T0025-NONATT-TYPE-2-UMS   VALUE 'U'.
15 T0025-NONATT-MTS-ORIG     PIC X(01).
88 T0025-NONATT-TYPE-3-UMS   VALUE 'U'.
15 T0025-NONATT-700-ORIG     PIC X(01).
88 T0025-NONATT-TYPE-4-UMS   VALUE 'U'.
15 T0025-NONATT-800-ORIG     PIC X(01).
88 T0025-NONATT-TYPE-5-TOLL   VALUE 'T' 'L'.
88 T0025-NONATT-TYPE-5-UMS   VALUE 'U'.
15 T0025-NONATT-DB800-ORIG   PIC X(01).
88 T0025-NONATT-DB800-TOLL   VALUE 'T' 'L'.
88 T0025-NONATT-DB800-UMS   VALUE 'U'.
88 T0025-NONATT-DB800-EMR   VALUE 'E'.
15 T0025-NONATT-900-ORIG     PIC X(01).
88 T0025-NONATT-TYPE-6-TOLL   VALUE 'T' 'L'.
88 T0025-NONATT-TYPE-6-UMS   VALUE 'U'.
15 T0025-NONATT-MTS-TERM     PIC X(01).
88 T0025-NONATT-TYPE-7-UMS   VALUE 'U'.

10 T0025-CI C000-DB800-ORIG   PIC X(01).
88 T0025-CI C000-DB800-TOLL   VALUE 'T' 'L'.
88 T0025-CI C000-DB800-UMS   VALUE 'U'.
88 T0025-CI C000-DB800-EMR   VALUE 'E'.

```

SAMPLE COPYBOOK REXX

Here is the same sample after it has been run through the utility and converted:

parse var record 1	T0025_TABLE_IDENT	6
/* 88-Level */ T0025_TABLE_ID = (T0025_TABLE_IDENT='T0025')		
parse var record 6	T0025_FROM_NPA	9
parse var record 9	T0025_FROM_NXX	12
parse var record 6	T0025_FROM_NPA_NXX	12
parse var record 12	T0025_TO_NPA	15
parse var record 6	T0025_PARTIAL_KEY	15
parse var record 15	T0025_EFFECTIVE_START	21
parse var record 21	T0025_EFFECTIVE_END	27
parse var record 6	T0025_KEY	27
parse var record 27	T0025_TABLE_LEVEL	28
parse var record 28	T0025_JURISDICTION	29
parse var record 29	T0025_STATE_CODE	31
parse var record 31	T0025_OPERATING_GROUP	32
parse var record 32	T0025_CONSOL_COMPANY	33
parse var record 33	T0025_COMPANY	34
parse var record 34	T0025_AREA_DIV	35
parse var record 35	T0025_DIV_DIST	37
parse var record 37	T0025_PLANT_CODE	41
parse var record 31	T0025_BILLING_LOCATION_CODE	41
parse var record 41	T0025_BAN_STATE_CODE	42
parse var record 42	T0025_ORIG_MMU	43
parse var record 43	T0025_ORIG_RATE_CNTR	46
parse var record 46	T0025_OLATA_CODE	49
parse var record 49	T0025_BILLING_RAO	52
parse var record 52	T0025_ASSOC_BELL_RAO	55
parse var record 55	T0025_RCC_THRESHOLD_SEC	57
parse var record 57	T0025_IND_EQUAL_ACCESS	58
/* 88-Level */ T0025_EQUAL_ACCESS = (T0025_IND_EQUAL_ACCESS='Y')		
/* 88-Level */ T0025_NON_EQUAL_ACCESS =		
(T0025_IND_EQUAL_ACCESS='N')		
parse var record 58	T0025_FGB_MPB_IND	59
parse var record 59	T0025_FGCD_MPB_IND	60
parse var record 60	T0025_CONV_OPH_INTRA	62
parse var record 62	T0025_CONV_OPH_INTER	64
parse var record 64	T0025_CONV_DIAL	66
parse var record 66	T0025_RECORD_POINT_DDD	69
parse var record 69	T0025_RECORD_POINT_OPH	72
parse var record 72	T0025_RECORD_POINT_TERM	75
parse var record 75	T0025_ATT_OPH_ZERO_PLUS	76
/* 88-Level */ T0025_ATT_TYPE_1_TOLL =		
(T0025_ATT_OPH_ZERO_PLUS='T') (T0025_ATT_OPH_ZERO_PLUS='L')		
(T0025_ATT_OPH_ZERO_PLUS='C')		
/* 88-Level */ T0025_ATT_TYPE_1_UMS =		
(T0025_ATT_OPH_ZERO_PLUS='U')		
/* 88-Level */ T0025_ATT_TYPE_1_EMR =		
(T0025_ATT_OPH_ZERO_PLUS='D')		
parse var record 76	T0025_ATT_OPH_ZERO_MINUS	77
/* 88-Level */ T0025_ATT_TYPE_2_TOLL =		
(T0025_ATT_OPH_ZERO_MINUS='T') (T0025_ATT_OPH_ZERO_MINUS='L')		

```

(T0025_ATT_OPH_ZERO_MINUS='C')
    /* 88-Level */ T0025_ATT_TYPE_2_UMS =
(T0025_ATT_OPH_ZERO_MINUS='U')
    /* 88-Level */ T0025_ATT_TYPE_2_EMR =
(T0025_ATT_OPH_ZERO_MINUS='D')
    parse var record 77      T0025_ATT_MTS_ORIG          78
        /* 88-Level */ T0025_ATT_TYPE_3_TOLL = (T0025_ATT_MTS_ORIG='T')
| (T0025_ATT_MTS_ORIG='L')
    /* 88-Level */ T0025_ATT_TYPE_3_UMS = (T0025_ATT_MTS_ORIG='U')
    parse var record 78      T0025_ATT_700_ORIG          79
        /* 88-Level */ T0025_ATT_TYPE_4_TOLL = (T0025_ATT_700_ORIG='T')
| (T0025_ATT_700_ORIG='L')
    /* 88-Level */ T0025_ATT_TYPE_4_UMS = (T0025_ATT_700_ORIG='U')
    parse var record 79      T0025_ATT_800_ORIG          80
        /* 88-Level */ T0025_ATT_TYPE_5_TOLL = (T0025_ATT_800_ORIG='T')
| (T0025_ATT_800_ORIG='L')
    /* 88-Level */ T0025_ATT_TYPE_5_UMS = (T0025_ATT_800_ORIG='U')
    parse var record 80      T0025_ATT_DB800_ORIG         81
        /* 88-Level */ T0025_ATT_DB800_TOLL = (T0025_ATT_DB800_ORIG='T')
| (T0025_ATT_DB800_ORIG='L')
    /* 88-Level */ T0025_ATT_DB800_UMS = (T0025_ATT_DB800_ORIG='U')
    /* 88-Level */ T0025_ATT_DB800_EMR = (T0025_ATT_DB800_ORIG='E')
    parse var record 81      T0025_ATT_900_ORIG          82
        /* 88-Level */ T0025_ATT_TYPE_6_TOLL = (T0025_ATT_900_ORIG='T')
| (T0025_ATT_900_ORIG='L')
    /* 88-Level */ T0025_ATT_TYPE_6_UMS = (T0025_ATT_900_ORIG='U')
    parse var record 82      T0025_ATT_MTS_TERM          83
        /* 88-Level */ T0025_ATT_TYPE_7_UMS = (T0025_ATT_MTS_TERM='U')
    parse var record 75      T0025_ATT_SOURCE_OF_DATA     83
    parse var record 83      T0025_NONATT_OPH_ZERO_PLUS   84
        /* 88-Level */ T0025_NONATT_TYPE_1_TOLL =
(T0025_NONATT_OPH_ZERO_PLUS='T') | (T0025_NONATT_OPH_ZERO_PLUS='L')
        /* 88-Level */ T0025_NONATT_TYPE_1_UMS =
(T0025_NONATT_OPH_ZERO_PLUS='U')
    parse var record 84      T0025_NONATT_OPH_ZERO_MINUS 85
        /* 88-Level */ T0025_NONATT_TYPE_2_TOLL =
(T0025_NONATT_OPH_ZERO_MINUS='T') | (T0025_NONATT_OPH_ZERO_MINUS='L')
        /* 88-Level */ T0025_NONATT_TYPE_2_UMS =
(T0025_NONATT_OPH_ZERO_MINUS='U')
    parse var record 85      T0025_NONATT_MTS_ORIG         86
        /* 88-Level */ T0025_NONATT_TYPE_3_UMS =
(T0025_NONATT_MTS_ORIG='U')
    parse var record 86      T0025_NONATT_700_ORIG         87
        /* 88-Level */ T0025_NONATT_TYPE_4_UMS =
(T0025_NONATT_700_ORIG='U')
    parse var record 87      T0025_NONATT_800_ORIG         88
        /* 88-Level */ T0025_NONATT_TYPE_5_TOLL =
(T0025_NONATT_800_ORIG='T') | (T0025_NONATT_800_ORIG='L')
        /* 88-Level */ T0025_NONATT_TYPE_5_UMS =
(T0025_NONATT_800_ORIG='U')

```

parse var record 88 T0025_NONATT_DB800_ORIG	89
/* 88-Level */ T0025_NONATT_DB800_TOLL =	
(T0025_NONATT_DB800_ORIG='T') (T0025_NONATT_DB800_ORIG='L')	
/* 88-Level */ T0025_NONATT_DB800_UMS =	
(T0025_NONATT_DB800_ORIG='U')	
/* 88-Level */ T0025_NONATT_DB800_EMR =	
(T0025_NONATT_DB800_ORIG='E')	
parse var record 89 T0025_NONATT_900_ORIG	90
/* 88-Level */ T0025_NONATT_TYPE_6_TOLL =	
(T0025_NONATT_900_ORIG='T') (T0025_NONATT_900_ORIG='L')	
/* 88-Level */ T0025_NONATT_TYPE_6_UMS =	
(T0025_NONATT_900_ORIG='U')	
parse var record 90 T0025_NONATT_MTS_TERM	91
/* 88-Level */ T0025_NONATT_TYPE_7_UMS =	
(T0025_NONATT_MTS_TERM='U')	
parse var record 83 T0025_NONATT_SOURCE_OF_DATA	91
parse var record 91 T0025_CI_C000_DB800_ORIG	92
/* 88-Level */ T0025_CI_C000_DB800_TOLL =	
(T0025_CI_C000_DB800_ORIG='T') (T0025_CI_C000_DB800_ORIG='L')	
/* 88-Level */ T0025_CI_C000_DB800_UMS =	
(T0025_CI_C000_DB800_ORIG='U')	
/* 88-Level */ T0025_CI_C000_DB800_EMR =	
(T0025_CI_C000_DB800_ORIG='E')	
parse var record 27 T0025_FUNCTION	92
parse var record 6 T0025_TABLE_DATA	92

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Space abend reporter

INTRODUCTION

Since the dawn of computers, every now and then you get given the ‘expert’ advice: “If you prevent application abends, you will improve throughput”. No matter how easy it may sound, doing it for an entire installation is not such a trivial job. The fact is that exponential growth in data storage resulting from data-intensive applications, e-business, and regulatory requirements to retain data have driven the need for more efficient systems to manage

storage resources. It is also a fact that few businesses today can afford processing delays or application failures and reruns.

However, application failures and errors do occur. Although not explicitly responsible for failures in application programs, a storage administrator is the one who is called upon when failures occur, especially if there is even the slightest hint or suspicion that the error is related to storage management. On the other hand, the interpretation of cryptic and often hard to understand IBM operating system error codes and messages can be tedious and frustrating, especially for application programmers. Often the appropriate manuals are not readily available for the users to access. If the manuals are available, the error description and corrective procedure are often impossible to interpret – it is the storage administrator again who must interpret the error and find a cause, and provide the ‘patch’.

Amongst a large variety of possible errors the most common type is the inability to acquire a sufficient amount of DASD space. This usually produces a JCL error if primary space cannot be obtained while an x37 type of error (B37, D37, or E37 abend) is generated if secondary space is unavailable. These abends are simply telling us that the job was not able to obtain enough DASD space to continue. The most frequent response to it is to increase the amount of space requested. Calculating how much would be enough is not an easy task, however, and space requirements are often guessed at – and generally guessed high, to avoid any chance of x37 abends. This can only waste resources and increase volume fragmentation without eliminating x37 abends: a job that requests much more space than it will really use may not only run short of space itself, but may cause errors when the space it holds is unavailable for other jobs. It is often interesting to look at file usage statistics on a system, and one may find some files with up to 90% of their allocated space free. Therefore increasing the amount of space requested may make the situation worse rather than better.

In order to prevent these out-of-space errors, some installations scrutinize applications’ JCL parameters, which turns out to be

not only unreliable but also time-consuming and expensive in terms of human resources.

A more effective and efficient technique is to trap abends, and make available a variety of actions that would allow the job to complete. An important aspect of this automatic ‘only-as-needed’ approach is the assurance that overhead, fragmentation, and other side-effects that waste space are minimized. When a storage administrator develops a management structure and assigns constructs to DFSMS storage groups, the goal is to anticipate and avoid running out of allowable space and abending a job. Even when additional parameters are made available in DFSMS to handle space-related problems, the same issues remain: the complex implementation of global space recovery functions that are processed in an all-or-none fashion. High-use thresholds are established for each storage group in order to leave some percentage of space unused in the storage group to service an unusual space request, or some volumes/storage groups are placed in quiesced status, so that files are allocated only on these volumes/storage groups when a high threshold is exceeded.

Success in reducing the number of x37s is directly related to the number of volumes in the storage group – the more volumes each group has, the merrier. The use of dataclass is paramount too: a proper dataclass will ensure extra candidate volumes are allocated to the datasets, so even if the user does allocate insufficient space, DFSMS will prevent it from failing by allowing it to extend to additional volumes. However, if you abuse this multivolume concept, you will have too much waste space, because RLSE is not issued against multivolume datasets. Some users have found this to be a major stumbling block because of the JCL changes necessary to remove the traditional space parameters and to specify the proper dataclass.

Worth the attention of the storage administrator are the space constraint relief enhancements – they do two main things to address out-of-space abends:

- They relax the limit on the number of extents that can be used to satisfy allocations for new datasets or extends to a new volume.
- The amount of space requested can be reduced by a percentage if the initial attempt fails, thus increasing the likelihood of a successful allocation (and, of course, being made up for by subsequent secondary extents).

These are designed to reduce the likelihood of x37 abends, and experience so far shows that they work fairly well.

With z/OS 1.3, DFSMS gets even better at handling out-of-space (x37) abends:

- Extend storage groups – one can specify a second storage group name when defining a storage group. If a dataset wants to extend to another volume, but there is insufficient space in the primary group, the dataset can extend into this storage group. Primary allocations will not use this storage group.
- Overflow storage groups – similar to the above, except that they are used when DFSMS cannot initially allocate onto the primary storage group. This storage group becomes the dataset's primary storage group, ie second and other volumes must go into this storage group.

One can combine both features so that a storage group can be both an extend and overflow storage group. There is a lot of good information in the redbook *z/OS V1R3 DFSMS Technical Guide* (SG24-6569).

When it comes to dealing with out-of-space failures for VSAM datasets, things get a bit more complicated. The errors that occur when creating or extending VSAM datasets are believed to be the most cryptic in MVS, mainly consisting of just a return code and a reason code. Among these messages one of the most common is IDC3009I, issued when a catalog function fails. This particular message includes literally hundreds of possible return/reason codes, a few of which are encountered more frequently

than others. A return code 68 says that there is a space error in DADSM (no space is available on the user volume). When accompanied by reason code 20, 22, or 24, it tells us that the storage group did not contain sufficient space to allocate the dataset. More precisely, it means that an attempt to allocate a DFSMS-managed dataset, using the best-fit interface, allocated space on all supplied volumes, but the amount of space allocated was not the total amount required. The DADSM error codes are fully documented in the manual *DFPSMSdfp Diagnosis Reference* (LY27-9606).

COLLECTING X37 DATA

No matter how fine-tuned DFSMS might be, the x37 abends will not disappear, simply because DFSMS won't keep applications from abending – no such facility exists in DFSMS. Obviously, x37 errors can be reduced through the use of DFSMS facilities like extended format datasets or space constraint relief, extend storage groups, and/or overflow storage groups. These are good facilities that can greatly reduce the likelihood, but they don't stop a job from aborting when allocated space is insufficient.

In order to reduce the occurrence of x37 abends as well as to keep things under control a 'space abend' reporting program was written. Running it on a daily basis may help the storage administrator by identifying necessary pool size changes that can be made, and in discussion with the applications programmer on how the dataset got to be bigger than anticipated. This report could also be used to measure and predict how the storage policies impact the user community.

One of the best ways to report on space abends is through the SMF records. Before proceeding any further it might be helpful to understand the context in which certain SMF records, and the values within them, are written. In this particular case, one may be tempted to choose to process less useful SMF type 4 and type 30 records, and report the jobs that abended with B37, D37, and E37 abends.

When enabled by the proper SMFPRMxx TYPE parameter, SMF creates type 42, subtype 9 records, which are written each time an x37 failure occurs for non-VSAM files, while record type 64 (VSAM component/cluster status) was used to report out-of-space abends pertaining to a VSAM dataset. A detailed description of the layout of SMF type 42 and 64 records can be obtained from the *MVS System Management Facilities (SMF)*, SA22-7630-03 manual. One can also find the type 42 subtype descriptions in macro IGWSMF in SYS1.MACLIB. For type 64 one should consult macro IDASMF64 in the same library.

CODE

Based on record descriptions obtained from the manual, a sample report writer was written.

The code is a two-part stream. In the first part (COPYSMF) selected SMF records (selection being defined by INCLUDE conditions) are copied from the SMF dataset to a file, which can be used for archived records.

In the second part (ABD37), the captured records are formatted by invoking EXEC (AB37) and two reports are produced.

The X37 report shows information such as jobname, dataset name, volume serial number, number of dataset extents, DFSMS-related information, as well as date and time. The second report (under DDn V37) is for VSAM out-of space conditions and it provides information such as jobname, component name, number of extents, etc.

COPYSMF

```
//DEL      EXEC PGM=IDCAMS
//SYSPRINT DD SYSOUT=X
//SYSIN    DD *
      DELETE ui d. ABEND. TEST
      DELETE ui d. X377. DATA
      DELETE ui d. V377. DATA
      SET MAXCC=0
```

```

//COPYSMF EXEC PGM=ICETOOL
//TOOLMSG DD SYSOUT=*
//DFSMMSG DD SYSOUT=*
//RAWSMF DD DSN=your.smf.dataset,DISP=SHR
//SMFX37 DD DSN=uid.ABEND.TEST,
//          SPACE=(CYL,(1)),UNIT=SYSDA,
//          DISP=(NEW,CATLG,KEEP),
//          DCB=(RECFM=VB,LRECL=32756,BLKSIZE=32760)
//TOOLIN DD *
   COPY FROM(RAWSMF) TO(SMFX37) USING(SMFI)
//SMFI CNTL DD *
   OPTION SPANNING=RC4,VLSHRT
   INCLUDE COND=(6,1,BI,EQ,42,AND,23,2,BI,EQ,9,      * Get SMF type 42.9
                 OR,(6,1,BI,EQ,64,AND,43,1,BI,EQ,X'20')) * also copy SMF 64
/*
//ABD37 EXEC PGM=IKJEFT01,REGION=0M,DYNAMNBR=50,PARM='%AB37'
//SYSEXEC DD DISP=SHR,DSN=your.rexx.library
//SMF DD DISP=SHR,DSN=uid.ABEND4.TEST
//X37 DD DSN=uid.X377.DATA,           <-- X37 report list
//          SPACE=(CYL,(1,1)),UNIT=SYSDA,
//          DISP=(NEW,KEEP),DCB=(RECFM=FB,LRECL=145)
//V37 DD DSN=uid.V377.DATA,           <-- VSAM report list
//          SPACE=(CYL,(1,1)),UNIT=SYSDA,
//          DISP=(NEW,KEEP),DCB=(RECFM=FB,LRECL=95)
//SYSPRINT DD SYSOUT=*
//SYSTSPRT DD SYSOUT=*
//SYSTSIN DD DUMMY
/*

```

AB37 EXEC

```

/* REXX EXEC to read and format space abend records */
/* trace ?r */

/*
/* Print Space Abend report header and labels
*/
Out.1 = left(' ',30,' '),
        ||center('Space Abend Report ',22,),
        ||left(' ',15,' ')
Out.2 = left(' ',20,' '),
        ||center('Report produced on',18,),
        ||left(' ',1,' ')||left(date(),11),
        ||left(' ',1,' ')||left('at ',3,' '),
        ||left(time(),10)
Out.3 =
Out.4 = left('Date',11)    left('Time',9),
        left('Job',6)      left('Step',4),
        left('X37',4)      left('Data Set',41),

```

```

    right('Disp', 4)      right('Dsorg', 6),
    left('Volume', 7)     left('Ext.', 4),
    left('Trks', 4)       left('Stor.Cl.', 12),
    left('Mngt.Cl.', 10)  left('Data Cl.', 10)
Out.5 = LEFT('-', 145, '-')

"EXECIO * DISKW X37 (STEM Out.)"
/*-----*/
/* Print VSAM 'out-of space' header and labels */
/*-----*/
lin.1 = left(' ', 30, ' ')
    ||center('VSAM out-of-space Report ', 30, ),
    ||left(' ', 15, ' ')
lin.2 = left(' ', 20, ' ')
    ||center('Report produced on', 18, ),
    ||left(' ', 1, ' ')||left(date(), 11),
    ||left(' ', 1, ' ')||left('at ', 3, ' '),
    ||left(time(), 10)
lin.3 =
lin.4 = left('Date', 11) left('Time', 9),
        left('Job', 8)   left('Component name', 45),
        left('Trks', 4)   right('Ext', 4) left('Volume', 6)
lin.5 = LEFT('-', 145, '-')

"EXECIO * DISKW V37 (STEM Lin.)"

/*-----*/
/* Header for SMF records */
/*-----*/
' EXECIO * DISKR SMF ( STEM x. FINIS'
do i = 1 to x.Ø

smftype = c2d(SUBSTR(x. i , 2, 1))          /* SMF record type */
smftime = smf(c2d(SUBSTR(x. i , 3, 4)))      /* Decode SMF time */
smfdate = SUBSTR(c2x(SUBSTR(x. i , 7, 4)), 3, 5) /* Unpack SMF date */

IF smftype = '64' Then call SMF64
IF smftype = '42' Then
Do
    sysid   = SUBSTR(x. i , 11, 4)           /* System identification */
    sywid   = SUBSTR(x. i , 15, 4)           /* Subsystem id */
    smfstype = c2d(SUBSTR(x. i , 19, 2))      /* Record subtype */

/*-----*/
/* Product Section */
/*-----*/

ops = c2d(SUBSTR(x. i , 25, 4))      /*Offset to product section */
lps = c2d(SUBSTR(x. i , 29, 4))      /*Length to product section */

```

```

nps = c2d(SUBSTR(x.i, 31, 4))      /*Number to product sections*/
IF ops <> 0 AND lps <> 0 Then do
  ops=ops -3
  pdl = SUBSTR(x.i, ops, 8)          /* Product level */
  pdn = SUBSTR(x.i, ops+8, 10)        /* Product name */
  psv = c2d(SUBSTR(x.i, ops+18, 1))  /* Subtype version number */
  SELECT
    when psv='0' then vhead ='No vol. header section'
    when psv='1' then vhead ='Vol. header exists'
  END
end

/*
/*-----*/
/*      SMF42 subtype 9 header section (B37/D37/E37 abend)      */
/*-----*/
abo = c2d(SUBSTR(x.i, 33, 4))      /*Offset to X37 abend section */
abl = c2d(SUBSTR(x.i, 37, 2))        /*Length of X37 abend section */
abn = c2d(SUBSTR(x.i, 39, 2))        /*Number of X37 abend sections*/
smo = c2d(SUBSTR(x.i, 41, 4))        /*Offset to SMS data section */
sml = c2d(SUBSTR(x.i, 45, 2))        /*Length of SMS data section */
smn = c2d(SUBSTR(x.i, 47, 2))        /*Number of SMS data sections */

/*
/*-----*/
/*      B37/D37/E37 Abend Data Section (part 1)                */
/*-----*/
IF abo <> 0 AND abl <> 0 Then do
  abo=abo -3
  sys1 = SUBSTR(x.i, abo, 4)          /* System ID */
  job = SUBSTR(x.i, abo+4, 8)          /* Job name */
  time = smf(c2d(SUBSTR(x.i, abo+12, 4))) /*Start time */
  date = SUBSTR(c2x(SUBSTR(x.i, abo+16, 4)), 3, 5) /*Start date */
  uid = SUBSTR(x.i, abo+20, 8)          /* User id. */
  step = c2d(SUBSTR(x.i, abo+28, 1))  /* Step no. */

  a37b = SUBSTR(x.i, abo+29, 1)        /* X37 flags */
  SELECT
    when a37b='10000000' b then abend ='B37'
    when a37b='01000000' b then abend ='D37'
    when a37b='00100000' b then abend ='E37'
  END

  dsorg= c2x(SUBSTR(x.i, abo+34, 2))  /*Dataset organization */
  SELECT
    when dsorg='4000' then Dorg='PS'
    when dsorg='2000' then Dorg='DA'
    when dsorg='0200' then Dorg='PO'
    otherwise dorg='PO'
  END

```

```

di sp = SUBSTR(x. i , abo+36, 1)                                /* Dataset di sp*/
SELECT
    when di sp ='11000001'b then di s ='Temp'
    when di sp ='11000000'b then di s ='New'
    when di sp ='10000000'b then di s ='Mod'
    when di sp ='01000000'b then di s ='01 d'
    when di sp ='01001000'b then di s ='Shr'
END

dsn  = SUBSTR(x. i , abo+37, 44)                                /* Dataset name*/
vol  = SUBSTR(x. i , abo+81, 6)        /*Vol ser of current volume*/
ucb  = c2d(SUBSTR(x. i , abo+87, 4))      /*UCB of current volume*/

ext  = c2d(SUBSTR(x. i , abo+91, 1)) /*No. of extents for dataset*/
                                         /* on current volume*/

trk  = c2d(SUBSTR(x. i , abo+92, 4)) /*Total tracks allocated for*/
                                         /*dataset on current volume*/

sal  = c2x(SUBSTR(x. i , abo+96, 4)) /*Secondary allocation amount*/
bl   = c2d(SUBSTR(x. i , abo+100, 3)) /*Average data block length*/
                                         /* if specified*/

/*-----*/
/*  SMS Data Secton
/*-----*/
IF smo > 0 Then do
    smo=smo -3
    mclas = SUBSTR(x. i , smo, 10)      /* Management class name*/
    sclas = SUBSTR(x. i , smo+30, 12)    /* Storage class name*/
    dclas = SUBSTR(x. i , smo+60, 10)    /* Data class name*/
end
else do
    mclas = ''
    sclas = 'Non SMS file'
    dclas = ''
end
end

x37out = right(Date('N', date, 'J'), 11) left(smftime, 9),
           left(job, 8) right(step, 2) left(abend, 4),
           left(dsn, 39) right(dis, 6) center(dorg, 6),
           left(vol, 8) right(ext, 3) right(trk, 4),
           right(sclas, 12) right(mclas, 10) right(dclas, 10)

PUSH x37out
"EXECIO 1 DISKW X37"

```

```

    end
end
/*-----*/
/* Print abend legend */
/*-----*/
Leg. 1 = LEFT('-', 145, '-')
Leg. 2 = ' '
Leg. 3 = left('X37 Abend Legend: ', 20)
Leg. 4 = ' '
Leg. 5 = left('B37: All 16 extents used or', 60)
Leg. 6 = left(' ', 5, ' '),
         ||left('No more space was available on the volume or', 80)
Leg. 7 = left(' ', 4, ' '),
         left('Unable to update the VTOC (full)', 78)
Leg. 8 = left('D37: Used all the primary space,', 33),
         ||left('and no secondary space was requested', 40)
Leg. 9 = left('E37: Used all space available on', 32),
         ||left(' ', 1, ' '),
         ||left('the current volume,', 31)
Leg. 10= left(' ', 5, ' '),
         ||left('and no more volumes were available', 40)
Leg. 11 = ' '
Leg. 12 = left('For details see the following messages', 70)
Leg. 13 = left('B37: IEC030I; D37: IEC031I; E37: IEC032I', 80)
"EXECIO * DISKW X37 (STEM Leg.)"
exit

/*-----*/
/* Selected SMF64 variables */
/*-----*/
SMF64:
jbb = SUBSTR(x.i, 15, 8)                      /* Job name*/
rin = SUBSTR(x.i, 39, 1)                         /*Recording indicators*/
SELECT
when rin='10000000' b then sit='Component closed';
when rin='01000000' b then sit='Vol switched';
when rin='00100000' b then sit='No space avail';
when rin='00010000' b then sit='Cat or CRA rec';
when rin='00001000' b then sit='Closed type=t';
when rin='00000100' b then sit='Abend process';
when rin='00000010' b then sit='VVDS or ICF';
otherwise sit='Reserved'
END
dnm = SUBSTR(x.i, 85, 44)                      /* Dataset name*/
ntr = c2d(SUBSTR(x.i, 129, 2))      /*Number of tracks required*/
esl = c2d(SUBSTR(x.i, 135, 2))      /*Extent segment length*/

/*-----*/
/* Extent table - one per volume online at time record written */

```

```

/*-----*/
noext=esl/26                                /*No. of extent segment*/
do extno=0 to noext - 1
  incr = (140 + (extno*26)) - 3
  vol.extno = SUBSTR(x.i, incr+8, 6) /*Vol ser of curr. volume*/
  volume    = SUBSTR(x.i, incr+8, 6) /*Vol ser of last volume*/
  cuu =c2d(substr(x.i, incr+14, 2))           /* device number*/
  ind =substr(x.i, incr+16, 2)                  /* Spindle id.*/
  uty =c2d(substr(x.i, incr+18, 4))           /* Unit type - ucbtype*/
end

v37out = right(Date('N', smfdate, 'J'), 11) left(smftime, 9),
           left(jbb, 8) left(dnm, 44),
           right(ntr, 4) right(noext, 5) left(volume, 6)

PUSH v37out
"EXECIO 1 DISK W V37"
return

SMF: procedure
/* REXX - convert a SMF time */
arg time
  time1 = time % 100
  hh    = time1 % 3600
  hh    = RIGHT("0"||hh, 2)
  mm    = (time1 % 60) - (hh * 60)
  mm    = RIGHT("0"||mm, 2)
  ss    = time1 - (hh * 3600) - (mm * 60)
  ss    = RIGHT("0"||ss, 2)
  otme = hh||":"||mm||":"||ss                /* Compose SMF time*/
  return otme

```

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As a free service to subscribers and to remove the need to rekey, the code from individual articles of *MVS Update* can be accessed on our Web site. You will be asked to enter a word from the printed issue.

Extending the life of your mainframe with in-memory table management

In a perfect world, mainframe computing is all about blazing performance, unparalleled scalability, and 99.999% availability. In reality, today's influx of process/transaction-intensive applications, such as CRM, knowledge management, data mining, Internet-based applications, etc, are taxing mainframe systems. Furthermore, because of budget cutbacks, most companies today are either reluctant or unable to spend extra money on mainframe hardware upgrades that are required to support the increased processing that these large and complex systems necessitate. For many, the answer to this dichotomy is found by making effective use of tables in software applications, coupled with the implementation of an in-memory table management system. For over 20 years, in-memory table management technology has helped Fortune 1000 companies to improve significantly the speed and performance of their mainframe applications, maximize transaction volumes, and support more end-users, without consuming additional hardware resources.

Most of you are aware that using tables in application development means easier and less costly application maintenance because table data, such as a product price list for example, is easily updated without incurring the costs of re-testing the application – and the risk of introducing new errors is virtually eliminated.

Although tables can appear to be a programmer's best friend, often their processing efficiency is compromised because of the I/O required, especially if the tables are frequently-accessed.

By using an in-memory table management system that loads tables into memory and allows these tables to be shared among applications and regions, application performance is significantly increased while I/O and CPU usage is significantly decreased. This is but one advantage of using an in-memory data management system. There are many others.

An in-memory data management system assures optimal use of existing hardware, and can even go so far as to support the postponement of costly hardware upgrades by making the applications more efficient – the more efficient the application, the less CPU usage is required. In addition, a table management system manages the tables and the memory, unburdening the development and maintenance staff to focus on other tasks.

For those unfamiliar with a table management system, it is similar to a database management system (DBMS), except that tables are loaded into memory. Since data manipulation is performed directly in memory, data access speed is greatly enhanced and the extensive I/O and CPU overhead of disk is avoided. The result is a radical increase in the performance of mainframe applications while extending the useful life of the mainframe system.

Other data management schemes do not provide holistic access to large volumes of structured data. They typically provide only sequential access to small amounts of data – normally thought of as a record, such as customer records from a CRM database or an employee record from an HR application. However, this type of record-level access is not optimal *if*:

- The data contained in the record needs to be accessed more than once – and possibly repeatedly – during an application run, which means 24x7 for several consecutive weeks for online applications.
- The data needs to be accessed by multiple mainframe applications at the same time, possibly from different regions.
- The data is transient, and only exists for a short period of time, like stock-price quote data.

Data falling into the above categories is handled effectively using a table-driven application design approach coupled with an in-memory table management system.

LOOKING OVER TABLES

Most people accept that a table is an iterative data structure consisting of a series of homogeneous data holders known as rows. Most also realize that a table is a logical structure, independent of its physical representation in memory or on a disk drive.

In the case of programming and database scenarios, the contents of tables can be thought of as consisting of two parts: keys (or indexes) and the data accessed using those keys. When dealing with most catalogue type applications, for example, a part number is used as the key required to access the detailed product description. In many cases, the key occurs within a specific column within the table. See column 1 in the mail-order catalogue example below:

<i>Item #</i>	<i>Product</i>	<i>Price</i>	<i>Availability</i>	<i>Link</i>
AP1183	Adobe Photoshop 6.5	549.00	Y	www.adobe.com ...
AP1266	Adobe Acrobat Full 5.4	499.00	Y	www.adobe.com ...
AP1270	Adobe InDesign 2.0	699.00	N	www.adobe.com ...
AP4881	AutoCAD LT 2002	619.95	N	www.autodesk.com ...
AP4895	AutoDesk Symbol 2000	94.95	Y	www.autodesk.com ...

The single-key-in-one-column approach may not be flexible or powerful enough in more complex situations. The application may want to use *multiple* keys – to search for items that are below \$500 *and* are in stock for example – or even use a retrieval mechanism that is not key-centric. In some instances, various applications – and even modules within them – may want to access the data using dissimilar criteria.

Tables, given their inherent logical nature, may be implemented in a number of different ways. Some popular approaches for implementing tables are:

- As structured, and possibly even indexed, files, like VSAM on disk.
- With a relational database using a database management system (DBMS), such as DB2 or Oracle.

- As a programming language array within a specific application, just for use by that application.
- As application-independent data structures resident in memory, external to any specific application programs, that are created, accessed, and maintained using a table management system.

Using memory-based data structures external to applications, which is true in-memory table management, is the only approach that delivers blazing performance, and simplified addition of new data extraction, data manipulation, and data correlation functions to existing applications.

COMBINING DBMSS AND IN-MEMORY TABLE MANAGEMENT

DBMSs and almost all other data-access technologies are designed to work with files, file structures or databases. They operate on the premise that, any time an update is performed, the DBMS will journal that update, reconcile all necessary changes and then make the updated record available to the next user as quickly as possible. Most enterprise data is processed this way. However, this approach is extremely I/O intensive and saps up valuable CPU cycles.

With in-memory table management, these same data tables can be loaded into memory and be readily available. This permits access to that data much faster than DBMS, and ensures that all requisite data manipulations can be performed directly within memory. For heavily-accessed data, application run-times and on-line responses can improve by as much as 90%.

EXTENDING LEGACY TO THE WEB

In-memory table management also helps to support personalization, and is proving particularly useful for extending legacy applications to the Web for real-time, high-volume e-commerce or e-business. It can help accommodate the resulting increase in traffic volumes within existing mainframe CPU, LPAR, and storage constraints, reducing or even eliminating the

need for hardware reallocation or additional hardware resources.

Three specific features make table management technology ideally suited for implementing fine-tuned personalization into corporate portal and e-business applications: speed, flexibility, and the capability to develop rules-based applications. All three support a successful delivery of customized user experiences to individual users. For instance, banking institutions that have adopted this technology reportedly credit memory-based table management as the enabling technology for providing personalized statements via the Web. These banks also use table management technology to produce consolidated and personalized master account statements for their customer base. To achieve this, the table-management-centric statement engine application applies 15 to 20 dynamic business rules per customer account in the process of handling millions of transactions per day.

In-memory table management is also suited for implementing incisive rules-based security policies for Web sites and corporate portals, and can be used for interactive, high-volume currency conversion applications as well as sales tax calculation applications. One of Ireland's largest banking corporations uses an in-memory table management system to handle all of its interactive currency conversion requirements, including the standardization to the euro currency in 2002. The bank's currency conversion application falls into the repetitive data-access category discussed earlier. Real-time sales tax calculation for point-of-sale applications and online mail-order applications also fall into this repetitive processing category. So do many security-related functions pertaining to corporate portals, whether they are for rules-based security policy enforcement, such as controlling access rights, or data encryption and decryption functions, like those related to SSL security.

BOTTOM LINE

Application design that takes advantage of tables can facilitate a host of requirements, from adding new functionality to existing

applications without requiring any modification or rewrites, to extending mainframe applications to the Web. Combining this approach with an in-memory table management system enables companies to speed-up process-intensive applications, support increased user volumes, optimize system efficiency and performance, facilitate and accelerate application development, and ultimately reduce costs. All in all, in-memory table management is a 20-year-old software methodology that is still in profitable use today, granting a continued lease on life to mission-critical mainframe applications.

*Robert Koblovsky
Vice President of Sales and Marketing
Data Kinetics Ltd (USA)*

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MVS news

BMC Software has announced three new MAINVIEW Storage Resource Manager (SRM) solutions – MAINVIEW SRM Allocation, MAINVIEW SRM Reporting, and MAINVIEW SRM Automation.

These products help customers to manage their mainframe storage environment by reducing storage costs and space shortages, providing automation capabilities that maximize staff and CPU time and minimize wasted space and schedule interruptions.

With these new products, BMC has automated allocation features to reduce the incidence of space-related processing problems and automated actions for the prevention and recovery from abends, ensuring higher levels of availability.

For further information contact:
BMC Software, 2101 City West Blvd,
Houston, TX 77042, USA.
Tel: (713) 918 8800.
URL: http://www.bmc.com/products/proddocview/0,2832,19052_19429_6115596_9804,00.html.

* * *

Serena Software has announced Version 4.1 of its Application Performance Manager, StarTool.

StarTool APM is a performance measurement and analysis system that helps to resolve OS/390 and z/OS job performance issues. Users can focus their activities on tuning specific areas of an application to boost productivity.

Version 4.1 provides improved, real-time monitoring capabilities with extended support for a wider range of subsystems including CICS, IMS, and DB2.

In addition, Version 4.1 delivers better reporting by providing customers with more in-depth information on an application to help quickly solve performance problems and eliminate the need for further analysis.

For further information contact:
Serena Software, 2755 Campus Drive, 3rd Floor, San Mateo, California 94403, USA.
Tel: (650) 522 6600.
URL: http://www.serena.com/product/aa_st_apm_ov.html.

* * *

iWay Software has announced improvements in its data integration capabilities, which include new cluster management, remote installation and management, tighter integration with SQL Server, and new data adapters.

The Remote Installation/Management Console is used for remote installations, which reduces the manual tasks involved in an OS/390, z/OS server installation. Server installations are simplified because the installation files can be transferred and configured from a local laptop or a remote computer, eliminating the need for certain types of technical and operational support.

Intelligent adapters have been added and enhanced to support DB2 Version 8 and IMS Version 8. Additionally, iWay's Adapter for CICS now supports RRS (Resource Recovery Services) and integrated two-phase commit.

For further information contact:
iWay Software, Two Penn Plaza, New York, NY 10121-2898, USA.
Tel: (212) 330 1700.
URL: <http://www.iwaysoftware.com/products/index.html>.



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